

ON DIFFERENCE-SPLITTING AND THE EQUAL WEIGHT VIEW

ABSTRACT. Dawid, DeGroot and Mortera showed, a quarter century ago, that any agent who regards a fellow agent as a peer—in particular, defers to the fellow agent’s prior credences in the same way that she defers to her own—and updates by split-the-difference is prone (on pain of triviality) to diachronic incoherence. On the other hand one may show that there are special scenarios in which Bayesian updating approximates difference splitting, so it remains an important question whether it remains a viable (approximate) response to “generic” peer update. We critique arguments by two teams of philosophers (Fitelson & Jehle and Nissan-Rozen & Spectre) against this updating scheme, then suggest an alternative “Equal Weight” response to cases of peer disagreement.

1. ON AN OLD TRIVIALITY RESULT FOR SPLIT-THE-DIFFERENCE

According to the so-called *Equal Weight View*, “When you count an advisor as an epistemic peer, you should give her conclusions the same weight as your own” (Elga 2007). Some philosophers have taken “splitting the difference” (i.e., adoption of the arithmetic mean) between competing peer credences to be constitutive of the Equal Weight View. Kelly (2010), e.g., writes:

...if the agnostic gives credence .5 to the proposition that God exists while the atheist gives credence .1 to the same proposition, the import of The Equal Weight View is clear: upon learning of the other’s opinion, each should give credence .3 to the proposition that God exists.

The popularity of Equal Weight difference splitting persists, despite the fact that it was shown, a quarter century ago in Dawid, DeGroot and Mortera (1995), to entail probabilistic incoherence. Indeed, it has been disputed that this coherence result applies at all to the situation that interests us, namely that of two peers having identical evidence but different priors. Bradley (2018), for example, rehearses a version of the result, but goes on to write:

...this study leaves open the question of whether linear averaging is the appropriate response to situations in which you find yourself in disagreement with peers who hold the same information as you and are as good at judging its significance. In the philosophical literature, the view that one should respond to such disagreements by taking an equal-weighted average of your opinions has been hotly debated. But nothing presented here militates either for or against this view.

In this section, we make the case that considerations very much along the lines of Dawid et. al. (1995) do “militate” (decisively) against split-the-difference in cases where two peers¹ hold the same information but do not have almost surely identical prior credences. We begin by supposing that i and j are agents and P is a proposition-valued random variable, so that i 's initial credence x in P and j 's initial credence y in P are also random variables. We take P 's distribution to be supported on propositions P_0 for which i and j are peers.

Because P is a random variable, there is some ambiguity concerning “credence in P ” we need to address. Suppose that a card, c , is drawn from a standard deck and P is either *Card c is a face card* or *Card c is an Ace* (with equal probabilities). In this case, P is in fact *Card c is an Ace*, but i doesn't know this. In such a case there are two readings of “ i 's credence in P ”. On the first reading, it refers to i 's credence in *Card c is an Ace*, namely $\frac{1}{13}$. On the second reading it refers to i 's credence in the proposition *either P is “Card c is a face card” and Card c is a face card or P is “Card c is an Ace” and Card c is an Ace*, namely $\frac{2}{13}$. We'll write $Cr_i(P)$ when we intend the first reading and $Cr_i(T(P))$ when we intend the second. (It may help to think of $T(P)$ as “ P is true”, read *de dicto*.) Notice that $Cr_i(P) = x$ regardless of whether i knows the value of P , whereas it will typically be the case that $Cr_i(T(P)) = x$ only after i learns P or at least $Cr_i(P)$.

When we say that i is “diachronically coherent almost surely” or “Reflection² obeying almost surely”, we mean that if i 's credence in P_0 is x_0 then

$$E_y(u(x, y)|P = P_0) = x_0,$$

where $u(x, y)$ denotes the posterior credence in P adopted by i upon learning the values of P , x and y . (Or in $T(P)$ upon learning just x and y .) In particular, for almost every x_0 in the essential range of x one has

$$E_y(u(x, y)|x = x_0) = x_0.$$

That is, if she were to learn her own initial credence $x = x_0$ (without learning P) then she would both come to have credence x_0 in $T(P)$ and expected posterior (posterior to learning y , that is) credence x_0 in $T(P)$.

¹Splitting the difference with a non-peer can, by contrast, be coherent; if my credence in P is one-half and I believe that your credence in whichever of P , $\neg P$ is true is 1 with probability .75 and zero with probability .25 then I surely don't consider you a peer (it seems that I think you are more sensitive to which of P , $\neg P$ is true, yet wildly overconfident), but should intend to split the difference with you when you tell me your credence.

²“Reflection” was coined by van Fraassen (1984). Roughly, an agent satisfies it when her current credence in a proposition P is equal to the expectation of her credence in P at a future time t , where t is typically an almost-surely future, possibly random time satisfying certain technical criteria (a so-called “stopping time”—see Schervish et. al. 2004). In the current application t is the time immediately after j 's credence in P is revealed.

Here is a natural necessary condition on peerhood.

PN: If i is diachronically coherent almost surely and regards j as a peer then for almost every y_0 in the essential range of y one has $E_x(u(x, y)|y = y_0) = y_0$.

PN cashes out the intuition that if i regards j as a peer then she has the same confidence in j 's initial credences that she has in her own. As observed above, if i were to learn that, and only that, $x = x_0$, then she would come to have credence x_0 in $T(P)$. So if she regards j as a peer and were to learn that, and only that, $y = y_0$, then she ought, similarly, to come to have credence y_0 in $T(P)$. But if she learns the value of x , say $x = x_0$, after learning that $y = y_0$, her posterior credence in $T(P)$ will be $u(x_0, y_0)$. By Reflection, then, her current (i.e. after learning $y = y_0$ but before learning $x = x_0$) credence, namely y_0 , should be the expectation of this posterior, that is $E_x(u(x, y)|y = y_0)$.

We can now establish incoherence (on pain of triviality) of split-the-difference (and more “general linear pooling”) for Equal Weighters.

Theorem 1. If i is diachronically coherent almost surely, regards j as a peer and updates by linear pooling (i.e. $u(x, y) = wx + (1 - w)y$, where $0 < w < 1$) then $Prob(x \neq y) = 0$.

Proof. For *reductio*, assume that $Prob(x \neq y) > 0$. The following is routine:

Lemma 1. If X is a random variable and $E(X) = k$ then

$$E(X - k|X > k)Prob(X > k) = E(k - X|X < k)Prob(X < k).$$

Since j is diachronically coherent almost surely, for almost every x_0 in the essential range of x one has

$$E_y((u(x, y)|x = x_0) = E_y(wx + (1 - w)y|x = x_0) = x_0.$$

But obviously $E_y(x|x = x_0) = x_0$, so in fact $E_y(y|x = x_0) = x_0$. By Lemma 1,

$$\begin{aligned} & E_y(y - x|x = x_0 \wedge y > x)Prob(y > x|x = x_0) \\ &= E_y(x - y|x = x_0 \wedge y < x)Prob(y < x|x = x_0). \end{aligned}$$

Multiplying both sides of this equation by x_0 , we can move occurrences of x_0 inside the expectations (since these are integrals in the variable y). So for a.e. x_0 ,

$$\begin{aligned} & E_y(x(y - x)|x = x_0 \wedge y > x)Prob(y > x|x = x_0) \\ &= E_y(x(x - y)|x = x_0 \wedge y < x)Prob(y < x|x = x_0). \end{aligned}$$

Integrating over x_0 , we get

$$E(x(y-x)|y > x)Prob(y > x) = E(x(x-y)|y < x)Prob(y < x). \quad (1)$$

Since i regards j as a peer, for almost every y_0 in the essential range of y one has $E_x(wx + (1-w)y|y = y_0) = y_0$ (by **PN**). But obviously $E_x(y|y = y_0) = y_0$, so in fact $E_x(x|y = y_0) = y_0$. By Lemma 1, it follows that

$$\begin{aligned} & E_x(x-y|y = y_0 \wedge x > y)Prob(x > y|y = y_0) \\ &= E_x(y-x|y = y_0 \wedge x < y)Prob(x < y|y = y_0). \end{aligned}$$

Multiplying by y_0 and employing $x > y$ or $y > x$ where applicable,

$$\begin{aligned} & E_x(x(x-y)|y = y_0 \wedge x > y)Prob(x > y|y = y_0) \\ & \geq E_x(y(x-y)|y = y_0 \wedge x > y)Prob(x > y|y = y_0) \\ & = E_x(y(y-x)|y = y_0 \wedge x < y)Prob(x < y|y = y_0) \\ & \geq E_x(x(y-x)|y = y_0 \wedge x < y)Prob(x < y|y = y_0), \end{aligned}$$

for a.e. y_0 , with strict inequality wherever $Prob(x > y|y = y_0)$ is positive. Integrating over y_0 , we therefore get

$$E(x(y-x)|x > y)Prob(x > y) > E(x(x-y)|x < y)Prob(x < y).$$

Multiplying both sides by -1 , swapping sides and rearranging some comparisons, we get

$$E(x(y-x)|y > x)Prob(y > x) > E(x(x-y)|y < x)Prob(y < x),$$

contradicting (1). qed

These considerations look to kill the difference splitting implementation of the Equal Weight View. On the other hand, coherent peer update schemes may *approximate* difference splitting. There is, moreover, a simple way of constructing such schemes. Namely, by considering the parallel case in which the agents had the same original priors, but have since acquired different evidence. Though this isn't the case we are interested in, the existence of these scenarios limits the methods by which one can argue. In particular, the existence of different-evidence scenarios in which approximate difference splitting is mandated prevents one from arguing against difference splitting in the same evidence case on incoherence grounds; for some joint distributions of the agents' priors (those arising in the different evidence case), it *is* coherent.

We sketch such a scenario. Suppose a point x is chosen uniformly at random on the unit interval. A standard Brownian motion Z is initiated at x and evolves

until it exits the interval. P is the event that it exits to the right, i.e. at 1. Neither i nor j know the value x . Suppose next that two independent standard Brownian motions, Z_i and Z_j , are initiated at x and stopped at time $t = 10^{-24}$. i is told the value $x_i = Z_i(t)$ and j is told the value $x_j = Z_j(t)$. Since the standard deviation, 10^{-12} , of x_i is so small, and since the expectation of x_i is x , the expectation x'_i of x conditional on x_i ($= i$'s probability for P conditional on x_i) will, with high probability, be extremely close to (distance much less than 10^{-12}) x_i . Similarly, j 's credence in P will be, with high probability, extremely close to x_j .

On the other hand, the expected value x' of x conditional on x_i and x_j will with high probability be extremely close to (distance much less than $|x_i - x_j|$) the midpoint of x_i and x_j ; so when i and j share their credences, they will, with high probability, adopt posterior credence $x' = u(x'_i, x'_j)$ in P extremely close to (relative to $|x'_i - x'_j|$) the midpoint of their shared credences x'_i and x'_j .

To reiterate, though this is not a same evidence scenario, the update rule $u(\cdot, \cdot)$ that falls out of it will be coherent in any same-evidence scenario in which the joint distribution of the peers agents' priors x'_i and x'_j is the same. So to argue against peer update schemes that approximate difference splitting, one would have either to propose norms directly constraining such joint distributions, or propose indirect constraints. We examine two approaches of the latter sort presently.

2. ON AN OVERRESTRICTIVE PEERHOOD CONSTRAINT OF NISSAN-ROZEN AND SPECTRE

Ittay Nissan-Rozen and Levi Spectre (2017) present an original argument against difference splitting as an implementation of the Equal Weight View. It fails, as we shall demonstrate. It begins with a novel proposed constraint on peerhood:

Our main contribution takes the form of a pragmatic constraint on the notion of peerhood: if an agent, j , is your peer, then assuming that j is sympathetic—she wants you to gain as much as possible—you should be willing—in exchange for a certain payoff—to let her decide for you whether to accept a bet with positive expected utility. If you are not willing to accept this exchange even for a sure payoff, you do not seriously regard j as your peer.

Nissan-Rozen and Spectre now prove the following theorem, in which P is a proposition for which i has an initial credence, and i is committed to updating via linear pooling (with weight w) upon learning j 's initial credence.

Theorem 2. (Nissan-Rozen and Spectre 2017) Let i be an agent for whom j is a fully rational and sympathetic peer. For any credence function of i that assigns a non-trivial probability value to the possibility that j 's degree of belief in P is different from i 's degree of belief in P , and for any $0 < w < 1$, there always exists

a bet with positive expected utility such that i (if she updates by linear pooling with weights $w, 1 - w$) will be willing to pay a positive amount of utility in order to avoid passing the choice of whether to accept the bet (on i 's behalf) to j .

The bet guaranteed by Theorem 2 violates Nissan-Rozen and Spectre's pragmatic constraint. If the constraint is viable, then, linear pooling Equal Weighters do not regard their fellow agents as peers, which implies in particular that difference splitting cannot be a viable implementation of the Equal Weight View.³

Nissan-Rozen and Spectre claim that their constraint (in conjunction with Theorem 2) "makes room for the development of a new Conciliatory view that calls for varying weights" (Nissan-Rozen and Spectre 2017). We interpret "variable weight" to imply an update rule for which $u(x, y)$ lies strictly between x and y , if $x \neq y$, and $u(x, x) = x$.⁴ But on this understanding, Theorem 2 doesn't, in fact, make room for such views; to what extent the argument attaches a deficient notion of peerhood to split-the-difference, plausible variable weight Conciliatory views are collateral damage.

To see why, let $u(y)$ be i 's posterior credence under such a scheme when i learns $y = Cr_j(P)$. Since we are assuming a "variable weight" rule the function $u(y)$ satisfies $u(y) = x$ for $x = y$, with $u(y)$ strictly between x and y otherwise. Since i regards j as a peer, meanwhile, we can assume that $u(y)$ is strictly increasing. Finally, since i is coherent, she should obey Reflection; in particular her initial credence $x = Cr_i(P)$ is the expectation of her posterior, i.e. $x = E_y(u(y))$.

Under these assumptions, one can always find a bet that i will pay a positive amount to avoid passing to j whenever y isn't, by i 's lights, equal to x almost

³Nissan-Rozen and Spectre also prove that there will be a bet with positive expected utility such that a difference splitting i will be inclined to pay a positive amount of utility in order to pass the bet to j . This violates an apparently endorsed (if only implicitly) variant of their constraint whereby you should be willing—in exchange for a certain payoff—to decide for yourself whether to accept a bet with positive expected utility in a case where you are otherwise obliged to pass it to j . The details aren't precisely the same, but this variant overgeneralizes as well, and so cannot be used to resuscitate the Nissan-Rozen/Spectre argument. In any event we set this aside, as they don't formally invoke (or even formulate) the variant in question.

⁴Some authors (e.g. Easwaran et. al. 2016) advocate for *synergy*, which implies that in case $y = x \neq \frac{1}{2}$, i 's posterior distribution should be more extreme than the common initial credence. Synergy is appropriate to the more common case where disagreeing peers have different evidence and the same priors. Suppose for example that i and j have common prior distribution that is uniform on $[0, 1]$ for the bias of a coin and are each allowed to toss the coin once, privately. If they reconvene and simultaneously announce credence of $\frac{1}{3}$ in the next toss of the coin landing *heads*, they will update not to $\frac{1}{3}$ but to $\frac{1}{4}$ (Laplace rule of succession). That is because their disclosures effectively allow for a pooling of evidence. Something like this is going on, for example, when so-called *meta-analyses* obtain "statistically significant" results (i.e. sufficiently extreme p values) by pooling studies that individually were unable to derive such results. In the same-evidence case we are interested in, however, the practice is plainly unjustified, indeed a bit like concluding that, because a certain balloon looks orange to everyone in the room, it must therefore be red.

surely. For in such a case i must, by non-triviality and Reflection, assign positive probability to the event $y < u(y) < x$. Let y_0 be the essential infimum of y .⁵ Choose k with $y_0 < k < u(y_0) < x$. Bet 1 pays 1 if P is true and pays 0 if P is false, if accepted; one receives a sure k if Bet 1 is rejected. Since i 's posterior (after learning y , that is) credence in P is almost surely greater than k , acceptance of Bet 1 has positive expected utility for i . She will be willing, moreover, to pay any amount less than $(u(y_0) - k)Prob(y < k) > 0$ to avoid having this bet passed to j , since j would reject it whenever $y < k$.

We think what Nissan-Rozen and Spectre had in mind was that one should deem a peer as being no worse (in expectation) than oneself when it comes to accepting or rejecting a bet of the form given (1 if the proposition is true and 0 if it is false, if accepted; a certain amount c if rejected), *prior to learning one's own initial credence x in the proposition in question*. Once one learns the value of x , that might change. If x is very close to c , the agent will recognize that the expected relative utility of her choice is small (non-existent, when $x = c$), and she may want to pass the bet to j . In at least some other cases (cases in which x and c are not close, typically), she will be inclined to want to field the bet herself.

The proposed constraint is therefore implausible—it would, if valid, rule out too much. That is to say, it can't be a requirement of peerhood that for every such offer one should think that one's peer has expected return not less than one's own. We conclude that Nissan-Rozen and Spectre's argument fails.

3. ON A WOULD-BE DESIDERATUM OF FITELSON AND JEHLÉ

Fitelson and Jehle (2009) attempt to discredit difference splitting *simpliciter* (i.e. their argument does not invoke peerhood) on the grounds that it fails to commute with conditionalization. Such an argument, it's probably worth mentioning, cannot counsel against difference splitting for two cell partitions, for the simple reason that if one conditionalizes on a non-trivial event from a two cell partition, the resulting space is trivial and there is only one candidate credence function over it. So Theorem 1 is more general, even if this alternative argument has merit.

The argument in its current form has serious problems, however, owing to the fact that Fitelson and Jehle believed the matter to be much simpler than it is. Indeed, they regarded it as transparent enough to relegate to a footnote:

⁵That is, $Prob(y < y_0) = 0$, but $Prob(y > y_0 + \epsilon) > 0$ for every $\epsilon > 0$.

Some Bayesian defenders of EWW require that (ideally) the result of an EWW update should be equivalent to a (classical) conditionalization, which conditionalizes “on whatever you...have learned about the circumstances of the disagreement” (Elga 2007, 490). If that’s right, then [commutativity] will follow from the definition of (classical) Bayesian conditionalization, since pairs of (classical) conditionalizations must commute. (Fitelson and Jehle 2009, footnote 12.)

That unfortunately isn’t right. What’s true is that if i conditions on A and then conditions on B , she should arrive at the same posterior as if she were to have first conditioned on B and then on A . But that’s not what’s going on here.

Suppose for example that i ’s original prior on (A, B, C) is $(\frac{1}{2}, \frac{1}{4}, \frac{1}{4})$ and j ’s is $(\frac{1}{4}, \frac{1}{2}, \frac{1}{4})$. It is certainly true that i should arrive at the same posterior if she first learned j ’s credence function then learned $\neg C$ as she should if she first learned $\neg C$ and then learned j ’s original prior. And these are just the propositions i will learn in a case where the agents first perform a peer update and then condition on $\neg C$. It is not, however, what i will learn if the agents first condition on $\neg C$ and then perform a peer update. In the latter case, i will learn only that j ’s original prior was of the form $(x, 2x, 1 - 3x)$ for some $0 < x \leq \frac{1}{3}$. There is no reason to think, then, that her posterior here must be the same.

In fact, one can easily construct coherent peer update rules that fail to commute with conditionalization: rules that approximate split-the-difference, for example, will violate this commutativity. The two-cell scheme presented at the end of Section 1 can be adapted to three cells to this end, as we now show.

Consider an equilateral triangle and a point generated uniformly at random in its interior, having barycentric coordinates (x, y, z) (x, y and z denote the distances from the point to the sides of the triangle; we assume $x + y + z = 1$). A two dimensional Brownian motion will be initiated at this point. When it hits a side (i.e. when one of the coordinates becomes zero), the Brownian motion will become 1-dimensional on that side until it terminates at a vertex. Let P_X be the event that the motion terminates at the vertex X having barycentric coordinates $(1, 0, 0)$; P_Y and P_Z are similarly defined.

Neither i nor j knows the initial point (x, y, z) . However, they each learn the identity of a nearby point—points (x_i, y_i, z_i) and (x_j, y_j, z_j) respectively—chosen from independent bivariate normal distributions having mean at the point with barycentric coordinates (x, y, z) and common, extremely small known variance. Assuming the agents to be rational, their resulting credences in (P_X, P_Y, P_Z) will be $(x'_i, y'_i, z'_i) \approx (x_i, y_i, z_i)$ and $(x'_j, y'_j, z'_j) \approx (x_j, y_j, z_j)$. Upon sharing these credences, they will each come to have posterior credence $(x', y', z') \approx \frac{1}{2}((x'_i, y'_i, z'_i) +$

(x'_j, y'_j, z'_j)) in (P_X, P_Y, P_Z) . The error in these approximations will be small compared to the distance between (x'_i, y'_i, z'_i) and (x'_j, y'_j, z'_j) with very high probability.

In particular, the error will always be small in those (extremely) rare cases where the points (x'_i, y'_i, z'_i) and (x'_j, y'_j, z'_j) are far from each other (and not too close to the edges). For example, when $(x'_i, y'_i, z'_i) = (.02, .2, .78)$ and $(x'_j, y'_j, z'_j) = (.8, .08, .12)$, peer update will result in a credence $\approx (.41, .14, .45)$. If one then conditions on $\neg P_Z$ one will obtain posterior $\approx (\frac{41}{55}, \frac{14}{55}, 0)$. On the other hand if i and j first condition on $\neg P_Z$ they will come to have credences $(\frac{1}{11}, \frac{10}{11}, 0)$ and $(\frac{10}{11}, \frac{1}{11}, 0)$, respectively. If now they perform a peer update, preservation of zero considerations and symmetry imply a posterior of $(\frac{1}{2}, \frac{1}{2}, 0) \not\approx (\frac{41}{55}, \frac{14}{55}, 0)$. So commutativity of conditionalization and peer update simply doesn't follow from naive Bayesian (i.e. coherence) considerations alone.

Fitelson and Jehle did go on to say (as a hedge, perhaps): “But even if we don't think of EWV-rules as equivalent to some conditionalization, we think [commutativity with conditionalization] should remain a desideratum for EWV-updates. We don't have the space to defend this claim here.” It's of course a pity that they do not defend the claim, as it certainly requires defense.

Any such defense would have to begin, we believe, with an attempt to explain away examples such as the foregoing one in which something near to difference splitting is rationally mandated. Note that the example favors difference splitting because the Euclidean midpoint of the segment connecting the ordered pairs whose barycentric coordinates correspond to the agents' priors minimizes the sum of the absolute deviations of the approximating bivariate normals (and so is near to the expectation of their mean). One would have to say, then, why the Euclidean metric is the wrong one to be working with in the generic situation in which two agents have identical evidence but different priors.

On the other hand, perhaps one would not have to say much here, for there is absolutely no reason to think that the Euclidean metric *would* be an appropriate metric in this context. When measuring the distance from a probability measure $x = (x_1, \dots, x_n)$ to another probability measure $y = (y_1, \dots, y_n)$, the information distance—so called *Kullback-Leibler divergence* $KL(x, y) = \sum_{i=1}^n x_i \log \frac{x_i}{y_i}$, is a far more likely default candidate. And, as we shall see below, when i and j 's common posterior is chosen so as to minimize the sum of these distances to their respective priors, the resulting update scheme *does* commute with conditionalization.

Alternatively, one can argue that commutativity is appropriate in cases where one doesn't have any reason to suspect it would fail. First one would argue that, in a case where i knows j 's prior credence function and knows that $\neg C$ (say) is the case, knowledge of her own current credence function “screens off” the significance of her prior credence in C . The example involving barycentric coordinates shows why one cannot make this assumption on the basis of coherence

considerations alone...the joint distribution of the two priors and the partition in question ($\{A, B, C\}$, say) may be such that i 's prior credence in C yields information about the relative likelihoods of A and B beyond that provided by j 's prior credence function and her own current credence function alone. That example relies heavily on the two agents' different evidence, however. One might argue that the same-evidence situation is different.

As an example, consider again the scenario in which i 's original prior on (A, B, C) is $(\frac{1}{4}, \frac{1}{4}, \frac{1}{2})$ and j 's is $(\frac{1}{4}, \frac{1}{2}, \frac{1}{4})$. Suppose that if the agents first perform a peer update, their common posterior will be (x, y, z) . (In this case $y = z$ follows from some seemingly innocuous permutation considerations, but we don't require this here.) If they next condition on $\neg C$, their common posterior will be $(\frac{x}{x+y}, \frac{y}{x+y})$. We want to say that if the agents instead: (1) condition on $\neg C$; then (2) perform a peer update, they will arrive at the same common posterior $(\frac{x}{x+y}, \frac{y}{x+y})$. It's clear that their credences after (1) will be $(\frac{1}{2}, \frac{1}{2}, 0)$ and $(\frac{1}{3}, \frac{2}{3}, 0)$ respectively.

Consider an alternate scenario in which the agents instead; (0) share their prior credences in C ; then (1) condition on $\neg C$; and finally (2) perform a peer update. Here the agents will definitely land at common posterior $(\frac{x}{x+y}, \frac{y}{x+y})$ after step (2), for in this case they would have acquired exactly the same information as in the original case where they first performed a peer update, then conditioned on $\neg C$. It is arguable, moreover, that j 's evidence in step (0), namely i 's prior credence in C , should in the absence of any reason for thinking the contrary be treated as neutral with respect to the relative likelihoods of A and B . Indeed, expert testimony as to probability of C (say) is considered *a* or *the* paradigm case in which so-called *Jeffrey conditionalization* (adopt the expert's credence in C as your own, preserve ratios of other partition cells; Jeffrey 1965) is appropriate. At the beginning of step (1) in the alternate scenario, learning your peer's credence in C is arguably no more nor less informative than learning your own credence in C at the conclusion of step (1), and so formally equivalent to taking expert (i.e. that of your future self) testimony as to the likelihood of C . If that's right, however, then the agents' credences after performing step (0) will have form (a, a, k) and $(b, 2b, l)$ respectively, so that after step (1) they will have credences $(\frac{1}{2}, \frac{1}{2}, 0)$ and $(\frac{1}{3}, \frac{2}{3}, 0)$ respectively, exactly as in the actual scenario of the previous paragraph. Since the final posteriors in the alternate scenario are $(\frac{x}{x+y}, \frac{y}{x+y})$, then, the final posteriors in the scenario from the previous paragraph will be $(\frac{x}{x+y}, \frac{y}{x+y})$ as well. That is, peer update commutes with conditionalization.

These considerations convince us that Fitelson and Jehle's desideratum is viable. Accordingly, we accept it; in the absence of known protocols to the contrary (in the same evidence case, in particular), peer update should commute with conditionalization.

4. PEER UPDATE AND RELATIVE ENTROPY MINIMIZATION

In this section we present an update scheme that we believe should in some sense be considered “standard”.⁶ The rule isn’t new. It is, for example, one member of a family of update rules considered in Easwaran et. al. (2016).⁷ On the other hand, we will argue its virtues in a couple of new ways. Among these virtues, one finds that it commutes with conditionalization, preserves independences and minimizes the sum of the Kullback-Leibler divergences from the priors to the common posterior. (We establish these properties in an appendix.)

According to this rule, if A and B are two cells of the partition under consideration and $r_i = \frac{P_i(A)}{P_i(B)}$, $r_j = \frac{P_j(A)}{P_j(B)}$ are ratios assigned by i and j to these cells’ respective probabilities, then the corresponding ratio arising from the common posterior ought to be the geometric mean of r_1 and r_2 . In the two-cell case $\{A, \neg A\}$, therefore, one updates to $(u, 1 - u)$ given priors $(x, 1 - x)$ and $(y, 1 - y)$, where

$$\frac{u}{1 - u} = \left(\frac{x}{1 - x} \frac{y}{1 - y} \right)^{1/2}. \quad (1)$$

As motivation for (1) we’ll provide two arguments (one heuristic) in the case where i ’s prior is $(1/5, 4/5)$ and j ’s is $(1/2, 1/2)$. That is, we’ll give plausible reasons why, in this case, the common posterior should be $(1/3, 2/3)$.

The first argument is simply the relative entropy one. The Kullback-Leibler divergence from $(u, 1 - u)$ to $(1/5, 4/5)$ is given by $u \log \frac{u}{1/5} + (1 - u) \log \frac{1-u}{4/5}$; the Kullback-Leibler divergence from $(u, 1 - u)$ to $(1/2, 1/2)$ is given by $u \log \frac{u}{1/2} + (1 - u) \log \frac{1-u}{1/2}$. The sum of these quantities is

$$H(u) = 2u \log u + 2(1 - u) \log(1 - u) - u \log 1/5 - (1 - u) \log 4/5 - \log 1/2.$$

The minimum of H occurs where

$$H'(u) = 2 \log \frac{u}{1 - u} + \log 4 = 0.$$

A quick calculation gives $(u, 1 - u) = (\frac{1}{3}, \frac{2}{3})$, in agreement with (1).

⁶We don’t claim that it would be immune from (coherence-based, even) objections. In particular, it would likely be subject to the usual array of criticisms that have plagued entropy maximization solutions in other contexts (see, e.g., Seidenfeld 1986 or Friedman and Shimony 1971). On the other hand, such solutions arguably do have formal merit as approximations to ideal behavior in extreme or limiting cases (we urge caution here, but see Vasudevan 2018). At any rate it isn’t possible to criticize a method that one doesn’t know about; this scheme, whatever its faults, marks a clear advance on difference splitting, and should be disseminated.

⁷These authors, curiously, do not favor the member of the family we are interested in. This is because they advocate for synergy in updating; we took issue with this in footnote 4.

For the second argument, imagine that a fair coin will be tossed if and only if $\neg A$ obtains. If i and j expand the algebra to accommodate the toss then of course their expanded priors will be $(1/5, 2/5, 2/5)$ and $(1/2, 1/4, 1/4)$. Suppose now they were to condition on the disjunction of the first two cells (the event $A \vee (\neg A \wedge \text{heads})$, say) and then perform a peer update. Their credences after conditioning will be $(1/3, 2/3, 0)$ and $(2/3, 1/3, 0)$. Symmetry and preservation of zero considerations now indicate that their credences will be $(1/2, 1/2, 0)$ after the peer update.

We next assume that i and j 's peer updating commutes with conditionalization. (As we stated at the end of Section 3, we accept Fitelson and Jehle's proposed desideratum.) Thus they will also come to have credence function $(1/2, 1/2, 0)$ should they peer update then condition. After the initial peer update they will have credences of the form (x, x, y) . On the other hand, permutation considerations point to $x = y$. That implies that i 's posterior credence in A , when she considers the coin toss, is $1/3$. The final step is then that peer update should commute with marginalization onto the original sub-algebra.

Beware: i mustn't subscribe to the commutativity of peer updating and marginalization in general. What justifies it in this case (the proponent will say) is that the ratio of the sizes of the to-be-amalgamated subcells is uncontroversial. Indeed, in the case where i and j first marginalize, then update, i 's credences first evolve from $(1/5, 2/5, 2/5)$ to $(1/5, 4/5)$, whereupon she learns j 's post-marginalization credence function, namely $(\frac{1}{2}, \frac{1}{2})$. Since the coin is uncontroversially fair, however, this gives her knowledge of j 's pre-marginalization credence function as well, namely $(\frac{1}{2}, 1/4, 1/4)$. So she acquires the same information she would acquire if the agents were to perform a peer update first, then marginalize. Accordingly, the update and the marginalization ought to commute, *in this special case*.⁸

Unlike split-the-difference, the proposed scheme can be coherently implemented with a peer. For imagine a proposition-valued random variable P . Denote i 's initial credence in P by x and j 's initial credence in P by y . Suppose that i regards j as a peer and updates in agreement with (1). We assume, for simplicity, that i 's joint distribution for (x, y) is distributed on eight pairs, with weights as indicated in Table 1.

It is now easy to see that i is Reflection-obeying. For example, if P_0 is such that $x = 1/5$ then, upon learning that $P = P_0$, i 's posterior distribution for u will be $(2/5, 3/5)$ on $(0, 1/3)$. In particular, $E(u|P = P_0) = 1/5$. The remaining cases are similar, so i 's behavior under this model exhibits diachronic coherence.

⁸One may make a fruitful comparison to the "Infomin" solution to the Judy Benjamin problem (van Fraassen 1981) here. When Judy receives a message yielding information about the relative sizes she ought to assign the Red regions, this may (says Infomin) influence her credence in *Blue*—but not in a case where the message fails to alter Judy's relative credences in the Red regions (*a fortiori*, in a case where Judy knows this in advance).

TABLE 1

x	y	u	Prob
0	1/5	0	1/10
1/5	0	0	1/10
1/5	1/2	1/3	3/20
1/2	1/5	1/3	3/20
1/2	4/5	2/3	3/20
4/5	1/2	2/3	3/20
4/5	1	1	1/10
1	4/5	1	1/10

The model moreover represents a plausible implementation of the Equal Weight View. (Apart from employing the EWW-friendly (1), i 's joint distribution for (x, y) and update function $u(x, y)$ are symmetric in the variables x and y , implying that, from i 's perspective, her own credences and those of j are treated interchangeably.) Since, then, it is not the case that $x = y$ almost surely, we may conclude that Theorem 1 doesn't overgeneralize in the manner of Theorem 2.

5. APPENDIX

Theorem 1. Let $\{A_1, A_2, \dots, A_n\}$ be a measurable partition. Suppose that i and j have priors

$$\mu_i = (a_1, a_2, \dots, a_n) \quad \text{and} \quad \mu_j = (b_1, b_2, \dots, b_n)$$

over (A_1, A_2, \dots, A_n) . Let $H(\mu) = KL(\mu, \mu_i) + KL(\mu, \mu_j)$ take on its minimum value (as μ ranges over probability measures on (A_1, A_2, \dots, A_n)) at $\mu_0 = (c_1, c_2, \dots, c_n)$. Then for any fixed indices $l, m, 1 \leq l \neq m \leq n$,

$$c_l^2 a_m b_m = c_m^2 a_l b_l. \quad (1)$$

In particular, if $a_m b_m \neq 0$ then $\frac{c_l}{c_m}$ is the geometric mean of $\frac{a_l}{a_m}$ and $\frac{b_l}{b_m}$.

Proof. Permuting indices if necessary, we may assume that $l = 1$ and $m = n$. Writing $0 \cdot \log 0 = 0$, $H(\mu)$ is continuous on a compact domain and so attains its minimum value at some $\mu_0 = (c_1, c_2, \dots, c_n)$. Plainly $c_1 = 0$ if $a_1 b_1 = 0$ and $c_n = 0$ if $a_n b_n = 0$; in either case, (1) follows.

We may therefore assume that $a_1 b_1 a_n b_n > 0$. Writing $\mu = (x_1, x_2, \dots, x_n)$,

$$H(\mu) = H(x_1, \dots, x_{n-1}) = \sum_{t=1}^n (2x_t \log x_t - x_t \log a_t b_t).$$

Since $x_n = 1 - x_1 - x_2 - \dots - x_{n-1}$, the first partial derivative of H is

$$H_{x_1}(x_1, \dots, x_{n-1}) = 2 \log \frac{x_1}{x_n} + \log \frac{a_n b_n}{a_1 b_1}.$$

Since H takes on its minimum value at one must have $H_{x_1}(\mu_0) = 0$. Therefore, $\log(\frac{c_1}{c_n})^2 = \log \frac{a_n b_n}{a_1 b_1}$. Clearing logarithms, we obtain (1). qed

Theorem 2. Let $\{A_1, A_2, \dots, A_n\}$ be a measurable partition. Let u be the update function having the property that

$$u((a_1, a_2, \dots, a_n), (b_1, b_2, \dots, b_n)) = (c_1, c_2, \dots, c_n),$$

where $c_l^2 a_m b_m = c_m^2 a_l b_l$, $1 \leq l \neq m \leq n$. Then u commutes with conditionalization on events in the algebra generated by A_1, A_2, \dots, A_n .

Proof. Let E be a union of partition cells. By rearranging indices if necessary we may assume that $E = A_1 \cup A_2 \cup \dots \cup A_k$ for some k . Suppose we first condition on E , then update. After conditioning on E , we get

$$\mu'_i = \left(\frac{a_1}{a_1 + a_2 + \dots + a_k}, \frac{a_2}{a_1 + a_2 + \dots + a_k}, \dots, \frac{a_k}{a_1 + a_2 + \dots + a_k}, 0, \dots, 0 \right)$$

and

$$\mu'_j = \left(\frac{b_1}{b_1 + b_2 + \dots + b_k}, \frac{b_2}{b_1 + b_2 + \dots + b_k}, \dots, \frac{b_k}{b_1 + b_2 + \dots + b_k}, 0, \dots, 0 \right).$$

Now $u(\mu'_i, \mu'_j) = (d_1, d_2, \dots, d_k, 0, \dots, 0)$, where

$$\begin{aligned} & d_l^2 \frac{a_m b_m}{(a_1 + a_2 + \dots + a_k)(b_1 + b_2 + \dots + b_k)} \\ &= d_m^2 \frac{a_l b_l}{(a_1 + a_2 + \dots + a_k)(b_1 + b_2 + \dots + b_k)}, \end{aligned}$$

which implies that $d_l^2 a_m b_m = d_m^2 a_l b_l$, $1 \leq l \neq m \leq k$.

Next we first update, then condition on E . We have $u(\mu_i, \mu_j) = (c_1, c_2, \dots, c_n)$, where $c_l^2 a_m b_m = c_m^2 a_l b_l$, $1 \leq l \neq m \leq n$. Conditioning next on E , we get to $(d_1, d_2, \dots, d_k, 0, \dots, 0)$, where $(c_1 + c_2 + \dots + c_k)d_m = c_m$, $1 \leq m \leq k$. Therefore

$$d_l^2 (c_1 + c_2 + \dots + c_k)^2 a_m b_m = d_m^2 (c_1 + c_2 + \dots + c_k)^2 a_l b_l,$$

so that $d_l^2 a_m b_m = d_m^2 a_l b_l$, $1 \leq l \neq m \leq k$. Since these equations clearly determine d_1, d_2, \dots, d_k subject to the constraint $d_1 + d_2 + \dots + d_k = 1$, we are done. qed

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