Epistemic Modality, Mind, and Mathematics

Hasen Khudairi

June 20, 2017
Abstract

This book concerns the foundations of epistemic modality. I examine the nature of epistemic modality, when the modal operator is interpreted as concerning both apriority and conceivability, as well as states of knowledge and belief. The book demonstrates how epistemic modality relates to the computational theory of mind; metaphysical modality; the types of mathematical modality; to the epistemic status of undecidable propositions and abstraction principles in the philosophy of mathematics; to the modal profile of rational intuition; and to the types of intention, when the latter is interpreted as a modal mental state. Chapter 2 argues for a novel type of expressivism based on the duality between the categories of coalgebras and algebras, and argues that the duality permits of the reconciliation between modal cognitivism and modal expressivism. Chapter 3 provides an abstraction principle for epistemic intensions. Chapter 4 advances a two-dimensional truthmaker semantics, and provides three novel interpretations of the framework along with the epistemic and metasemantic. Chapter 5 applies the modal μ-calculus to account for the iteration of epistemic states, by contrast to availing of modal axiom 4 (i.e. the KK principle). Chapter 6 advances a solution to the Julius Caesar problem based on Fine’s "criterial" identity conditions which incorporate conditions on essentiality and grounding. Chapter 7 provides a ground-theoretic regimentation of the proposals in the metaphysics of consciousness and examines its bearing on the two-dimensional conceivability argument against physicalism. The epistemic two-dimensional truthmaker semantics developed in chapter 4 is availed of in order for epistemic states to be a guide to metaphysical states in the hyperintensional setting. Chapter 8 examines the modal commitments of abstractionism, in particular necessitism, and epistemic modality and the epistemology of abstraction. Chapter 9 examines the modal profile of Ω-logic in set theory. Chapter 10 examines the interaction between epistemic two-dimensional se-
mantics and absolute decidability. Chapter 11 avails of modal coalgebraic automata to interpret the defining properties of indefinite extensibility, and avails of epistemic two-dimensional semantics in order to account for the interaction of the interpretational and metaphysical modalities thereof. The hyperintensional, epistemic two-dimensional truthmaker semantics developed in chapter 4 is applied in chapters 8, 10, and 11. Chapter 12 provides a modal logic for rational intuition. Chapter 13 examines modal responses to the alethic paradoxes. Chapter 14 examines, finally, the modal semantics for the different types of intention and the relation of the latter to evidential decision theory.
Acknowledgements

From 2014 to 2017, I was a Ph.D. Student at the Arché Philosophical Research Centre for Logic, Language, Metaphysics and Epistemology at the University of St Andrews. St Andrews is an ideal place to live and work. At Arché, I was supported by a St Leonard’s College (SASP) Research Scholarship, for which I record my gratitude.

This book is a revised version of my dissertation. The dissertation was written between the foregoing years, and revised in the years that followed.

For comments which lead to revisions to individual chapters, I am grateful to Aaron Cotnoir, Josh Dever, Peter Milne, and Gabriel Uzquiano.

For productive conversations at Arché, I am grateful to Mark Bowker, Sarah Broadie, Federico Faroldi, Katherine Hawley, Patrik Hummel, Ryo Ito, Bruno Jacinto, Li Kang, Kris Kersa, Martin Lipman, Poppy Mankowitz, Matthew McKeever, Daniel Nolan, Laurie Paul, Andrew Peet, Stephen Read, Justin Snedegar, Mark Thakkar, Jens Timmermann, Michael Traynor, Brian Weatherson, and Erik Wielenberg. For her administrative assistance, I am grateful to Lynn Hynd.

From 2005 to 2008, I was an honors undergraduate in philosophy at Johns Hopkins University. For their encouragement and example, I am grateful to Michael Williams, Hent de Vries, Meredith Williams, Dean Moyar, and Maura Tumulty.

For his friendship and for visits at the beginning and end of the time this dissertation was written, I am grateful to Armand Leblois. For their unwavering support of my endeavors over the years, I am grateful to my parents.

The book is dedicated to Alison Bowen.

Table of Contents

1. Methodological Forward (p. 11)

   Part I: A Framework for Epistemic Modality

2. Modal Cognitivism and Modal Expressivism (p. 24)
   2.1 Introduction
   2.2 The Hybrid Proposal
      2.2.1 Epistemic Modal Algebra
      2.2.1.1 Epistemic Truthmaker Semantics
      2.2.2 Modal Coalgebraic Automata
      2.3 Material Adequacy
      2.4 Expressivist Semantics for Epistemic Possibility
      2.5 Modal Expressivism and the Philosophy of Mathematics
      2.6 Concluding Remarks

3. Cognitivism about Epistemic Modality (p. 44)
   3.1 Introduction
   3.2 An Abstraction Principle for Epistemic Intensions
   3.3 Examples in Philosophy and Cognitive Science
   3.4 Objections and Replies
   3.5 Concluding Remarks

4. Two-Dimensional Truthmaker Semantics (p. 57)
   4.1. Introduction
4.2 Two-Dimensional Truthmaker Semantics
4.2.1 Intensional Semantics
4.2.2 Truthmaker Semantics
4.2.3 Two-Dimensional Truthmaker Semantics
4.3 New Interpretations
4.3.1 Fundamental and Derivative Truths
4.3.2 Decision Theory
4.3.3 Intentional Action
4.4 Concluding Remarks

5. Non-Transitive Self-Knowledge: Luminosity via Modal $\mu$-Automata (p. 73)

Part II: Conceivability and Metaphysical Possibility

6. Conceivability, Haecceities, and Essence (p. 81)
6.1 Introduction
6.2 Super-Rigidity
6.3 Two Dogmas of Semantic Rationalism
6.3.1 The First Dogma
6.3.2 The Second Dogma
6.3.2.1 The Julius Caesar Problem
6.3.3 Mereological Parthood
6.3.4 Summary
6.4 Determinacy and Consistency
6.5 Concluding Remarks


Part III: Epistemic Modality and the Philosophy of Mathematics

8. Epistemic Modality, Necessitism, and Abstractionism (p. 112)
8.1 Introduction
8.2 The Abstractionist Foundations of Mathematics
8.3 Abstraction and Necessitism
8.3.1 Hale and Wright’s Arguments against Necessitism
8.3.2 Hale on the Necessary Being of Purely General Properties and Objects
8.3.2.1 Objections
8.3.3 Cardinality and Intensionality
8.4 Epistemic Modality, Metaphysical Modality, and Epistemic Utility and Entitlement
8.4.1 Epistemic Two-dimensional Truthmaker Semantics
8.5 Concluding Remarks

9. $\Omega$-Logicism: Automata, Neo-logicism, and Set-theoretic Realism (p. 136)
9.1. Introduction
9.2 Definitions
9.2.1 Axioms
9.2.2 Large Cardinals
9.2.3 $\Omega$-Logic
9.3 Discussion
9.3.1 Neo-Logicism
9.3.2 Set-theoretic Realism
9.4 Concluding Remarks

10. Epistemic Modality and Absolute Decidability (p. 152)
10.1 Introduction
10.2 Mathematical Modality
10.2.1 Metaphysical Mathematical Modality
10.2.2 Epistemic Mathematical Modality
10.2.3 Interaction
10.2.4 Modal Axioms
10.3 Departures from Precedent
10.4 Knowledge of Absolute Decidability
10.5 Concluding Remarks

11. Grothendieck Universes, and Indefinite Extensibility (p. 165)
11.1 Introduction
11.2 Indefinite Extensibility in Set Theory: Modal and Extensional Approaches
11.3 Grothendieck Universes
11.4 Modal Coalgebraic Automata and Indefinite Extensibility
11.5 Concluding Remarks

12. A Modal Logic for Gödelian Intuition (p. 180)
12.1 Introduction
12.2 Rational Intuition as Cognitive Phenomenology
12.3 Modalized Rational Intuition and Conceptual Elucidation
12.4 Concluding Remarks

13. An Epistemicist Solution to Curry’s Paradox (p. 192)
13.1 Introduction
13.2 Scharp’s Replacement Theory
13.2.1 Properties of Ascending and Descending Truth
13.2.2 Scharp’s Theory: ADT
13.2.3 Semantics for ADT
13.3 New Extensions of ADT
13.3.1 First Extension: The Preface Paradox
13.3.2 Second Extension: Absolute Generality
13.3.3 Third Extension: Probabilistic Self-reference
13.3.4 Fourth Extension: The Sorites Paradox
13.4 Issues for ADT
13.4.1 Issue 1: Revenge Paradoxes
13.4.2 Issue 2: Validity
13.4.3 Issue 3: Hybrid Principles and Compositionality
13.4.4 Issue 4: ADT and Indeterminacy
13.4.5 Issue 5: Descending Truth, Ascending Truth, and Objectivity
13.4.6 Issue 6: Paradox, Sense, and Signification

13.5 Epistemicism and Alethic Paradox
13.6 Concluding Remarks

14. Epistemic Modality, Intention, and Decision Theory (p. 212)

14.1 Introduction
14.2 The Modes of Intention
14.2.1 Intention-in-Action
14.2.2 Intention-with-which
14.2.3 Intention-for-the-Future
14.3 Intention in Decision Theory
14.4 Concluding Remarks

Bibliography (p. 221)
Chapter 1

Methodological Forward

This book concerns the foundations of epistemic modality. The work aims to advance our present understanding of the defining contours of epistemic modal space. I endeavor, then, to develop the theory of epistemic modality, by accounting for its interaction with metaphysical modality; the types of mathematical modality; the epistemic status of undecidable propositions and abstraction principles in the philosophy of mathematics; the modal profile of rational propositional intuition; and the types of intention, when the latter is interpreted as a modal mental state. In each chapter, I examine the philosophical significance of the foregoing, by demonstrating its import to a number of previously intransigent philosophical issues.

In Section 1, I provide a summary of each of the chapters. In Section 2, I examine the limits of competing proposals in the literature, and outline the need for a new approach.

1.1 Chapter Summary

In Chapter 2, I provide a mathematically tractable background against which to model both modal cognitivism and modal expressivism. I argue that epistemic modal algebras, endowed with a hyperintensional truthtmaker semantics, comprise a materially adequate fragment of the language of thought. I demonstrate, then, how modal expressivism can be regimented by modal coalgebraic automata, to which the above epistemic modal algebras are dual. I examine, in particular, the virtues unique to the modal expressivist approach here proffered in the setting of the foundations of mathematics, by
contrast to competing approaches based upon both the inferentialist approach to concept-individuation and the codification of speech acts via intensional semantics.

In Chapter 3, I aim to vindicate the thesis that cognitive computational properties are abstract objects implemented in physical systems. I avail of the equivalence relations countenanced in Homotopy Type Theory, in order to specify an abstraction principle for epistemic intensions. The homotopic abstraction principle for epistemic intensions provides an epistemic conduit into our knowledge of intensions as abstract objects. I examine, then, how intensional functions in Epistemic Modal Algebra are deployed as core models in the philosophy of mind, Bayesian perceptual psychology, and the program of natural language semantics in linguistics, and I argue that this provides abductive support for the truth of homotopic abstraction. Epistemic modality can thus be shown to be both a compelling and a materially adequate candidate for the fundamental structure of mental representational states, comprising a fragment of the language of thought.

In Chapter 4, I endeavor to establish foundations for the interaction between hyperintensional semantics and two-dimensional indexing. I examine the significance of the semantics, by developing three, novel interpretations of the framework. The first interpretation provides a characterization of the distinction between fundamental and derivative truths. The interaction between the hyperintensional truthmaker semantics and modal ontology is further examined. The second interpretation demonstrates how the elements of decision theory are definable within the semantics, and provides a novel account of the interaction between probability measures and hyperintensional grounds. The third interpretation concerns the contents of the types of intentional action, and the semantics is shown to resolve a puzzle concerning the role of intention in action. Two-dimensional truthmaker semantics can be interpreted epistemically and metasemantically, as well, and epistemic two-dimensional truthmaker semantics is examined in the chapter and in chapters 7 and 8, as well as appealed to in chapters 10-11.

In Chapter 5, I provide a novel account of iterated epistemic states. I argue that states of epistemic determinacy might be secured by countenancing self-knowledge on the model of fixed points in monadic second-order modal logic, i.e. the modal $\mu$-calculus. Despite the epistemic indeterminacy witnessed by the invalidation of modal axiom 4 in the sorites paradox – i.e. the KK principle: $\Box \phi \rightarrow \Box \Box \phi$ – an epistemic interpretation of a $\mu$-automaton permits fixed points to entrain a principled means by which to account for
necessary conditions on self-knowledge.

In Chapter 6, I aim to redress the contention that epistemic possibility cannot be a guide to the principles of modal metaphysics. I argue that the interaction between the two-dimensional intensional framework and the mereological parthood relation enables epistemic possibilities to target the haecceitistic properties of individuals. I specify, then, a two-dimensional intensional formula encoding the relation between the epistemic possibility of haecceity comprehension and its metaphysical possibility. I examine the Julius Caesar problem as a test case. I then generalize the approach to essential properties. I conclude by addressing objections from the indeterminacy of ontological principles relative to the space of epistemic possibilities, and from the consistency of epistemic modal space.

In Chapter 7, I argue that Chalmers’s (1996; 2010) two-dimensional conceivability argument against the derivation of phenomenal truths from physical truths risks being obviated by a hyperintensional regimentation of the ontology of consciousness. The regimentation demonstrates how ontological dependencies between truths about consciousness and about physics cannot be witnessed by epistemic constraints, when the latter are recorded by the conceivability – i.e., the epistemic possibility – thereof. Generalizations and other aspects of the philosophical significance of the hyperintensional regimentation are further examined.

In Chapter 8, I aim to provide modal foundations for mathematical platonism. I examine Hale and Wright’s (2009) objections to the merits and need, in the defense of mathematical platonism and its epistemology, of the thesis of Necessitism. In response to Hale and Wright’s objections to the role of epistemic and metaphysical modalities in providing justification for both the truth of abstraction principles and the success of mathematical predicate reference, I examine the Necessitist commitments of the abundant conception of properties endorsed by Hale and Wright and examined in Hale (2013a,b); examine cardinality issues which arise depending on whether Necessitism is accepted at first- and higher-order; and demonstrate how a two-dimensional semantic approach to the epistemology of mathematics, augmented with Necessitism, is consistent with Hale and Wright’s notion of there being epistemic entitlement rationally to trust that abstraction principles are true. A choice point that I flag is that between availing of intensional or hyperintensional semantics. The hyperintensional semantics approach that I favor is an epistemic two-dimensional truthmaker semantics, for which I define a model. Epistemic and metaphysical states and possibilities may thus be shown to
play a constitutive role in vindicating the reality of mathematical objects and truth, and in explaining our possible knowledge thereof.

In Chapter 9, I examine the philosophical significance of $\Omega$-logic in Zermelo-Fraenkel set theory with choice (ZFC). The duality between coalgebra and algebra permits Boolean-valued algebraic models of ZFC to be interpreted as coalgebras. The modal profile of $\Omega$-logical validity can then be countenanced within a coalgebraic logic, and $\Omega$-logical validity can be defined via deterministic automata. I argue that the philosophical significance of the foregoing is two-fold. First, because the epistemic and modal profiles of $\Omega$-logical validity correspond to those of second-order logical consequence, $\Omega$-logical validity is genuinely logical, and thus vindicates a neo-logicist conception of mathematical truth in the set-theoretic multiverse. Second, the foregoing provides a modal-computational account of the interpretation of mathematical vocabulary, adducing in favor of a realist conception of the cumulative hierarchy of sets.

In Chapter 10, I aim to contribute to the analysis of the nature of mathematical modality, and to the applications of the latter to unrestricted quantification and absolute decidability. Rather than countenancing the interpretational type of mathematical modality as a primitive, I argue that the interpretational type of mathematical modality is a species of epistemic modality. I argue, then, that the framework of two-dimensional semantics ought to be applied to the mathematical setting. The framework permits of a formally precise account of the priority and relation between epistemic mathematical modality and metaphysical mathematical modality. The discrepancy between the modal systems governing the parameters in the two-dimensional intensional setting provides an explanation of the difference between the metaphysical possibility of absolute decidability and our knowledge thereof. I mention again the choice point between intensional and hyperintensional semantics, where the latter is defined in detail in chapters 4, 7, and 8.

In Chapter 11, I endeavor to define the concept of indefinite extensibility in the setting of category theory. I argue that the generative property of indefinite extensibility in the category-theoretic setting is identifiable with the Kripke functors of modal coalgebraic automata, where set-coalgebras model Grothendieck Universes and the functors are further inter-definable with the elementary embeddings of large cardinal axioms. The Kripke functors are argued to account for both reinterpretations of quantifier domains as well as the ontological expansion effected by the elementary embeddings in the category of sets. The interaction between the interpretational and metaphysical
modalities of indefinite extensibility is defined via the epistemic interpretation of two-dimensional semantics. By characterizing the modal profile of \( \Omega \)-logical validity, and thus the generic invariance of mathematical truth, modal coalgebraic automata are further capable of capturing the notion of definiteness, in order to yield a non-circular definition of indefinite extensibility.

In Chapter 12, I aim to provide a modal logic for rational intuition. Similarly to treatments of the property of knowledge in epistemic logic, I argue that rational intuition can be codified by a modal operator governed by the axioms of a dynamic provability logic, which embeds GL within the modal \( \mu \)-calculus. Via correspondence results between modal logic and the bisimulation-invariant fragment of second-order logic, a precise translation can then be provided between the notion of 'intuition-of', i.e., the cognitive phenomenal properties of thoughts, and the modal operators regimenting the notion of 'intuition-that'. I argue that intuition-that can further be shown to entrain conceptual elucidation, by way of figuring as a dynamic-interpretational modality which induces the reinterpretation of both domains of quantification and the intensions of mathematical concepts that are formalizable in monadic first- and second-order formal languages.

In Chapter 13, I target a series of potential issues for a modal resolution to the alethic paradoxes. I aim, then, to provide a novel, epistemicist treatment to Curry’s Paradox. The epistemicist solution that I advance enables the retention of both classical logic and the traditional rules for the alethic predicate: truth-elimination and truth-introduction.

In Chapter 14, I argue that the types of intention can be modeled as a type of modal operator. I delineate the intensional-semantic profiles of the types of intention, and provide a precise account of how the types of intention are unified in virtue of both their operations in a single, encompassing, epistemic modal space, and their role in practical reasoning. I endeavor to provide reasons adducing against the proposal that the types of intention are reducible to the mental states of belief and desire, where the former state is codified by subjective probability measures and the latter is codified by a utility function. I argue, instead, that each of the types of intention – i.e., intention-in-action, intention-as-explanation, and intention-for-the-future – has as its aim the value of an outcome of the agent’s action, as derived by her partial beliefs and assignments of utility, and as codified by the value of expected utility in evidential decision theory.
1.2 The Need for a New Approach

The proposal that mental representations can be defined as possibilities relative to states of information dates at least back to Wittgenstein (1921/1974), although there are a number of precursors to the literature in the twentieth century. While novel, the limits of these incipient proposals consists in that they are laconic with regard to the explanatory foundations of the general approach.

Wittgenstein writes: 'A picture is a fact. / The fact that the elements of a picture are related to one another in a determinate way represents that things are related to one another in the same way. / Let us call this connexion of its elements the structure of the picture, and let us call the possibility of this structure the pictorial form of the picture. / Pictorial form is the possibility that things are related to one another in the same way as the elements of the picture. A logical picture of facts is a thought. / 'A state of affairs is thinkable': what this means is that we can picture it to ourselves. / The totality of true thoughts is a picture of the world. / A thought contains the possibility of the situation of which it is the thought. What is thinkable is possible too' (op. cit.: 2.141-2.151, 3-3.02). Wittgenstein notes, further, that 'The theory of knowledge is the philosophy of psychology' (4.1121), and inquires: 'Does not my study of sign-language correspond to the study of thought processes which philosophers held to be so essential to the philosophy of logic? Only they got entangled for the most part in unessential psychological investigations.'

1For an examination of epistemic logic in, e.g., the late medieval period, see Boh (1993). For the role of logical, rather than epistemic, modality in defining the modes of judgment, see Buridan (2001: 5.6), Kant (1787/1998: A74/B99-A76/B101), and Bolzano (1810/2004: 15-16). For the synthetic apriori determination of which of the possible predicates comprising a disjoint union ought to be applied to objects – i.e., transcendental logic – see Kant (op. cit: A53/B77-A57/B81; A571/B599-A574/B602). Anticipating Kripke (1980: 56), Husserl (1929/1999: §6) refers, in a section heading and the discussion therein, to transcendental logic as pertaining to conditions on the 'contingent apriori'.

For the role of possibilities in accounting for the nature of subjective probability measures, i.e., partial belief, see Bernoulli (1713/2006: 211), Wittgenstein (op. cit.: 4.464, 5.15-5.152), and Carnap (1945). Bernoulli (op. cit.) writes: 'Something is possible if it has even a very small part of certainty, impossible if it has none or infinitely little. Thus something that has 1/20 or 1/30 of certainty is possible'. For subjective interpretations of probability, see Pascal (1654/1959), Laplace (1774/1986), Boole (1854), Ramsey (1926/1960), de Finetti (1937/1964), and Koopman (1940). For the history of the development of the theory of subjective probability, see Daston (1988; 1994) and Joyce (2011).
tions, and there is an analogous danger for my method’ (op. cit.). Despite Wittgenstein’s reluctance to accept the bearing of cognitive psychology on thought, chapters 2 and 3 endeavor to argue that epistemic modality comprises a materially adequate fragment of the language of thought, i.e., the computational structure and semantic values of the mental representations countenanced in philosophy and cognitive science.

Modal analyses of the notions of apriority and of states of information broadly construed are further proffered in Russell (1919), Lewis (1923), and Peirce (1933).

Russell (op. cit.: 345-346) contrasts the possible truth-value of a propositional function (i.e., open formula) given an assignment of values to the variables therein with an epistemic – what he refers to as the ‘ordinary’ – interpretation of the modal according to which ‘when you say of a proposition that it is possible, you mean something like this: first of all it is implied that you do not know whether it is true or false, and I think it is implied; secondly, that it is one of a class of propositions, some of which are known to be true. When I say, e.g., ‘It is possible that it may rain to-morrow’ . . . We mean partly that we do not know whether it will rain or whether it will not, but also that we do know that that is the sort of proposition that is quite apt to be true, that it is a value of a propositional function of which we know some value to be true’ (op. cit.: 346).

Lewis (op. cit.: 172) defines the apriority of the laws of mathematical languages as consisting in their being ‘true in all possible worlds’.

Peirce (op. cit.: §65) writes:

'[L]et me say that I use the word information to mean a state of knowledge, which may range from total ignorance of everything except the meanings of words up to omniscience; and by informational I mean relative to such a state of knowledge. Thus, by ‘informationally possible,’ I mean possible so far as we, or the persons considered, know. Then, the informationally possible is that which in a given information is not perfectly known not to be true. The informationally necessary is that which is perfectly known to be true. The informationally contingent, which in the given information remains uncertain, that is, at once possible and unnecessary.’

The notion of epistemic modality was, finally, stipulated independently by Moore (c.1941-1942/1962) in his commonplace book. According to Moore,

---

2The remarks are anticipated in Wittgenstein [1979: 21/10/14, 5/11/14, 10/11/14, 12/11/14 (pp. 16, 24-29)].
'epistemic' possibilities include that 'Its possible that [for some individual, a: a] is [glad] right now [iff] [a] may be [glad]', where 'I know that he's not' contradicts the foregoing sentence (op. cit.: 187). Another instance of an epistemic possibility is advanced – 'It’s possible that I’m not sitting down right now’ – and analyzed as: 'It’s not certain that I am’ or 'I don’t know that I am” (184).

In the contemporary literature, there is a paucity of works devoted to the nature of epistemic modality and its relation to other modalities. Recent books and edited volumes which examine aspects of epistemic modality include Gendler and Hawthorne (2002); Yablo (2008); Gendler (2010; Egan and Weatherson (2011); and Chalmers (2012a). The present work is focused on the foundations and philosophical significance of the epistemic interpretation of modal logic and semantics. For the sake of completeness, a critical summary of the relevant literature is thus included below.

The Gendler and Hawthorne volume includes seminal contributions to the theory of the relationship between epistemic and metaphysical modality. By contrast, this book provides foundations for the nature of epistemic modality, when the modality concerns apriority and conceivability, as well as the logic of knowledge and belief; makes contributions to our understanding of the ontology of consciousness, by regimenting the ontology of consciousness using hyperintensional grounding operators; examines the nature and philosophical extensions of epistemic logic; and examines the relations between epistemic modality and the variety of other modalities (e.g., metaphysical and mathematical modalities and the types of intention in the setting of evidential decision theory).

The papers on modal epistemology in Yablo (2008) predominantly concern the relation between epistemic and metaphysical modalities, and, in particular, non-trivial conditions on modal error. Issues for the epistemic interpretation of two-dimensional intensional semantics are examined; e.g., the absence of conditions on ascertaining when an epistemic possibility is actual, and a dissociation in the case of recognitional concepts between conceptual necessity and apriority. The discussion is similar, in scope, to the discussions in the Gendler and Hawthorne volume. This book aims to redress the limits mentioned in the foregoing, and to proffer the positive proposals delineated above.

The Egan and Weatherson volume is comprised of papers which predominantly analyze epistemic modals in the setting of natural language semantics. Only three papers in the volume target epistemic possibilities as imaginable
or conceptual possibilities; those by Chalmers (‘The Nature of Epistemic Space’), Jackson (‘Possibilities for Representation and Credence’), and Yalcin (‘Non-factualism about Epistemic Modality’).

Chalmers’ paper examines some principles governing epistemic space and its interaction with metaphysical modality, as well as Kaplan’s paradox. As mentioned, the book endeavors, by contrast, further to explain the individuation-conditions on the terms and intensions defined in epistemic modal space; examines the interaction between epistemic modality and various other types of modality; and examines the role that epistemic logic plays in resolving the alethic paradoxes, as well as undecidable sentences in the philosophy of mathematics.

Jackson’s paper argues that conceptual possibilities and metaphysical possibilities ought to be defined within a single space, in order both to avoid cases in which a sentence is conceptually possible although metaphysically impossible and to secure the representational adequacy of conceptually possible terms. Chapter 4 adopts, by contrast, the modal dualist proposal to the effect that epistemic modality and metaphysical modality, as well as epistemic and metaphysical state spaces, occupy distinct spaces.

Yalcin’s paper argues that epistemic modal sentences in natural language semantics mirror the structure of the beliefs of speakers. Epistemic mental states are taken, then, to be expressive rather than representational, because the communication of epistemic modal and interrogative updates on an informational background shared by speakers is not truth-conditional. The present approach contrasts to the foregoing, by not taking the values of expressions in natural language semantics to be a guide to the nature of mental states (cf. Evans, 1982).

Chalmers (2012a) provides a book-length examination of the scrutability of truth, and the apriori entailment relations between different types of truths. The rigidity of intensions is availed of, in order to explain the relation between epistemic modality and metaphysical modality. By contrast, the book aims to examine novel philosophical extensions of epistemic two-dimensional semantics and the role of epistemic modality in the philosophy of mathematics and logic.

Gendler (2010) is a rare, empirically informed study of the limits of repre-

---

3The view that subject matters, broadly construed, have the form of an interrogative update on a set of worlds is anticipated by Lewis (1988/1998) and further defended by Yablo (2014 and Yalcin (2016). For further discussion of subject matters, see Chapter 4.
sentational capacities, when they target counterfactual assignments of values to variables in thought experiments – e.g., the conditions under which there might be resistance attending the states of imagining that fictional characters have variant value-theoretic properties – and when implicit biases and unconscious sub-doxtastic states affect the veridicality conditions of one’s beliefs. One crucial distinction between Gendler’s approach and the one pursued in this chapter, however, is that the former does not examine the interaction between epistemic modality and modal logic.

In the literature on modal epistemology, Hale (2013a) argues that modal knowledge ought to be pursued via the epistemology of essential definitions which specify conditions on sortal membership. Apriori knowledge of essence is explained in virtue of knowledge of the purely general terms – embedding no singular terms – which figure in the definitions. Thus – by being purely general – the essential properties and the objects falling in their extension have necessary being. Aposteriori knowledge of essential definitions can be pursued via theoretical identity statements, yet, because the terms figuring therein are not purely general, both the essential properties and the objects falling in their extension have contingent being. As mentioned, the book redefines the extant proposals in the ontology of consciousness using hyper-intensional grounding operators. The ground-theoretic interpretation of the ontology of consciousness, and an examination of the bearing of the latter for the relation between conceivability and metaphysical possibility is, as noted, examined in Chapter 7. Hale’s higher-order Necessitist proposal is examined in further detail, in Chapter 8.

Nichols (2006) features three essays on modal epistemology. Nichols’ ‘Imaginative Blocks and Impossibility’ examines introspection-based tasks in developmental psychology, in order to account for the interaction between imaginative exercises and counterfactual judgments. Hill’s ‘Modality, Modal Epistemology, and the Metaphysics of Consciousness’ examines the interaction between conceptual and metaphysical possibility, where conceptual possibilities are construed as Fregean thoughts, and the relation between conceivability and metaphysical possibility is then analyzed as the relation between Fregean thoughts (augmented by satisfaction-conditions such as conceptual coherence) and empirical propositions. Sorensen’s paper, ‘Meta-conceivability and Thought Experiments’, argues that meta-conceivable thought experiments are distinct from both conscious perceptual states and conceivable possibilities. [Sorensen (1999) argues that (thought) experiments track the consequences of reassignments of values to variables.] My approach dif-
fers from Hill’s by arguing in favor of both a possible worlds semantics as well as a hyperintensional, epistemic two-dimensional truthmaker semantics for thoughts, which is able to recover the virtues attending the Fregean model, as well as in accounting for the relations between epistemic modal algebras and various other interpretations of modality, including the mathematical (cf. Chapter 8 and 10 for further discussion). My approach is similar in methodology to Nichols’, although I endeavor to account to for the relation between conceivability and metaphysical possibility by availing of the epistemic interpretation of two-dimensional semantics. Finally, my approach is similar to Sorensen’s, in targeting both a formal semantic analysis of epistemic and related modalities, as well as the operators of knowledge and belief in the setting of epistemic logic.

Waxman (ms) endeavors to account for the interaction between the imagination and mathematics. Whereas I avail of conceivability as defined in epistemic two-dimensional semantics in Chapters 8 and 10 – which I refer to in the mathematical setting as epistemic mathematical modality – in order to account for how the epistemic possibility of abstraction principles and large cardinal axioms relates to their metaphysical possibility, Waxman’s aim is to account for how imagining a model of a mathematical theory entrains justification to believe its consistency (op. cit.). Unlike Waxman, epistemic mathematical modality is ideal, whereas imagination is, on his account, non-ideal (Waxman, op. cit.: 18; Chalmers, 2002), where ideal conceivable means true at the limit of apriori reflection unconstrained by finite limitations. Unlike Waxman, I believe, further, that imaginative contents are sensitive to hyperintensional subject-matters or topics (cf. Berto, 2018; Canavotto, Berto, and Giordani, 2020).

Finally, a class of views in the epistemology of modality can be characterized as being broadly empiricist. Stalnaker (2003) and Williamson (2007; 2013) refrain from countenancing the notion of epistemically possible worlds; and argue instead either that the imagination is identified with cognitive processes taking the form of counterfactual presupposition (Williamson, 2007); that one’s choice of the axioms governing modal logic should satisfy abductive criteria on theory choice (Williamson, 2013a); or that metaphysical modalities are properties of the actual world (Stalnaker, op. cit.). Vetter (2013) argues for a reduction of modal notions to actual dispositional properties, and Roca-Royes (2016) pursues a corresponding modal empiricist approach, according to which knowledge of the de re possibilities of objects consists in the extrapolation of properties from acquaintance with objects in one’s
surround to formally similar objects, related by reflexivity and symmetry. Generally, according to the foregoing approaches, the method of modal epistemology proceeds by discerning the modal truths – captured, e.g., by abductively preferred theorems in modal logic; conditional propositions; and dispositional and counterfactual properties – and then working backward to the exigent incompleteness of an individual’s epistemic states concerning such truths. By contrast, the approach advanced in this work both retains and provides explanatory foundations for epistemic modal space, and augments the examination by empirical research and an abductive methodology.

The foregoing texts either examine epistemic modality via natural language semantics; restrict their examination to the interaction between conceivable possibilities and metaphysical possibilities; eschew epistemic possibilities; provide a naturalistic approach to the analysis of epistemic modality, without drawing on formal methods; or provide a formal analysis of epistemic modality, without drawing on empirical results.

The book endeavors, by contrast, to examine the interaction between epistemic modality and the computational theory of mind; metaphysical modality; the types of mathematical modality; the modal profile of rational intuition; and the types of intention, when the latter is interpreted as a modal mental state.

The models developed here are of interest in their own right. However, this work is principally concerned with, and examines, their philosophical significance, as witnessed by the new distinctions and properties that they induce. Beyond conditions on theoretical creativity, both formal regimentation and empirical confirmation are the best methods available for truth-apt philosophical inquiry into both the space of epistemic modality and the multiple points of convergence between epistemically possible truth and the most general, fundamental structure of metaphysically possible worlds.
Part I: A Framework for Epistemic Modality
Chapter 2

Modal Cognitivism and Modal Expressivism

2.1 Introduction

This essay endeavors to reconcile two approaches to the modal foundations of thought: modal cognitivism and modal expressivism. The novel contribution of the paper is its argument for a reconciliation between the two positions, by providing a hybrid account in which both internal cognitive architecture, on the model of epistemic possibilities, as well as modal automata, are accommodated, while retaining what is supposed to be their unique and inconsistent roles.

Modal cognitivism is the proposal that the internal representations comprising the language of thought can be modeled via either a possible world or hyperintensional semantics. Modal expressivism has, in turn, been delineated in two ways. On the first approach, the presuppositions shared by a community of speakers have been modeled as possibilities (cf. Kratzer, 1979; Stalnaker, 1978, 1984). Speech acts have in turn been modeled as modal operators which update the common ground of possibilities, the semantic values of which are then defined relative to an array of intensional parameters (Stalnaker, op. cit.; Veltman, 1996; Yalcin, 2007). On the second approach, the content of concepts is supposed to be individuated via the ability to draw inferences, and the pragmatic abilities of individuals have been modeled as automata comprised of two transition functions. A counterfactual transition functional – encoding the recognition of distal properties – determines the
range of admissible values for another transition function encoding the individual’s actions (cf. Brandom, 2008). Inferential conditions constitutive of concept possession are then taken to have the same counterfactual form as the foregoing functions (Brandom, 2014), while truth-evaluable descriptions of the automata are specified in a metalanguage (Brandom, 2008). Both the modal approach to shared information and the speech acts which serve to update the latter, and the modal-inferential approach to concept-individuation – are thus consistent with mental states having semantic values or truth-conditional characterizations.

The notions of cognitivism and expressivism here targeted concern the role of internal – rather than external – factors in countenancing the nature of thought and information (cf. Fodor, 1975; Haugeland, 1978). Possible worlds or hyperintensional semantics is taken then to provide the most descriptively adequate means of countenancing the structure of the foregoing. Delineating cognitivism and expressivism by whether the positions avail of internal representations is thus orthogonal to the eponymous dispute between realists and antirealists with regard to whether mental states are truth-apt, i.e., have a representational function, rather than being non-representational and non-factive, even if real (cf. Dummett, 1959; Blackburn, 1984; Price, 2013). Whereas the type of modal cognitivism examined here assumes that thoughts and information take exclusively the form of internal representations, the target modal expressivist proposals assume that information states are exhaustively individuated by both linguistic behavior and conditions external to the cognitive architecture of agents.

So defined, the modal cognitivist and modal expressivist approaches have been assumed to be in constitutive opposition. While the cognitivist proposal avails of modal resources in order to model the internal representations comprising an abstract language of thought, the expressivist proposal targets informational properties which extend beyond the remit of internal cognitive architecture: both the form and the parameters relevant to determining the semantic values of linguistic utterances, where the informational common ground is taken to be reducible to possibilities; and the individuation of the contents of concepts on the basis of inferential behavior.

In this essay, I provide a background mathematical theory, in order to account for the reconciliation of the cognitivist and expressivist proposals. I avail, in particular, of the duality between Boolean-valued models of epistemic modal algebras and coalgebras; i.e., labeled transition systems defined
in the setting of category theory. The functors of coalgebras permit of flexible interpretations, such that they are able to characterize both modal logics as well as discrete-state automata. I argue that the correspondence between epistemic modal algebras and modal coalgebraic automata is sufficient then for the provision of a mathematically tractable, modal foundation for thought and action.

In Section 2, I provide the background mathematical theory, in order to account for the reconciliation of the cognitivist and expressivist proposals.

In Section 3, I provide reasons adducing in favor of modal cognitivism, and argue for the material adequacy of epistemic modal algebras as a fragment of the language of thought.

In Section 4, I outline an expressivist semantics for epistemic modals.

In Section 5, modal coalgebraic automata are argued, finally, to be preferred as models of modal expressivism, by contrast to the speech-act and inferentialist approaches, in virtue of the advantages accruing to the model in the philosophy of mathematics. The interest in modal coalgebraic automata consists, in particular, in the range of mathematical properties that can be recovered on the basis thereof. By contrast to the above competing approaches to modal expressivism, the functors of modal coalgebraic automata are able both to model and explain elementary embeddings in the category of sets; the intensions of mathematical terms; as well as the modal profile of Ω-logical consequence.

Section 6 provides concluding remarks.

2.2 The Hybrid Proposal

2.2.1 Epistemic Modal Algebra

An epistemic modal algebra is defined as $U = \langle A, 0, 1, \neg, \cap, \cup, l, m \rangle$, with A a set containing 0 and 1 (Bull and Segerberg, 2001: 28).

---

1For an algebraic characterization of dynamic-epistemic logic, see Kurz and Palmigiano (2013). Baltag (2003) develops a colagebraic semantics for dynamic-epistemic logic, where coalgebraic functors are intended to record the informational dynamics of single- and multi-agent systems. The current approach differs from the foregoing by examining the duality between static epistemic modal algebras and coalgebraic automata in a single-agent system.

2See Wittgenstein (2001: IV, 4-6, 11, 30-31), for a prescient expressivist approach to the modal profile of mathematical formulas.
\[ l_1 = 1, \]
\[ l(a \cap b) = l_a \cap l_b \]
\[ m_a = \neg l \neg a, \]
\[ m_0 = 0, \]
\[ m(a \cup b) = m_a \cup m_b, \text{ and} \]
\[ l_a = \neg m \neg a \text{ (op. cit.)}. \]

A valuation \( v \) on \( U \) is a function from propositional formulas to elements of the algebra, which satisfies the following conditions:

\[
\begin{align*}
v(\neg A) &= \neg v(A), \\
v(A \land B) &= v(A) \cap v(B), \\
v(A \lor B) &= v(A) \cup v(B), \\
v(\blacksquare A) &= l v(A), \text{ and} \\
v(\lozenge A) &= m v(A) \text{ (op. cit.)}. 
\end{align*}
\]

A frame \( F = \langle W, R \rangle \) consists of a set \( W \) and a binary relation \( R \) on \( W \) (op. cit.). \( R[w] \) denotes the set \( \{ v \in W \mid (w, v) \in R \} \). A valuation \( V \) on \( F \) is a function such that \( V(A, x) \in \{1, 0\} \) for each propositional formula \( A \) and \( x \in W \), satisfying the following conditions:

\[
\begin{align*}
V(\neg A, x) &= 1 \text{ iff } V(A, x) = 0, \\
V(A \land B, x) &= 1 \text{ iff } V(A, x) = 1 \text{ and } V(B, x) = 1, \\
V(A \lor B, x) &= 1 \text{ iff } V(A, x) = 1 \text{ or } V(B, x) = 1 \text{ (op. cit.)}. 
\end{align*}
\]

**Epistemic Truthmaker Semantics**

Chalmers endorses a principle of plenitude according to which "For all sentences \( s \), \( s \) is epistemically possible iff there exists a scenario [i.e. epistemically possible world - HK] such that \( w \) verifies \( s' \) (2011: 64), where \( w'[w] \)hen \( w \) verifies \( s \), we can say that \( s \) is true at \( w' \) (63). In this paper, I accept, instead, a hyperintensional truthmaker approach to epistemic possibility, defined by the notion of exact verification in a state space, where states are parts of whole worlds (Fine 2017a,b; Hawke and Özgün, forthcoming). According to truthmaker semantics for epistemic logic, a modalized state space model is a tuple \( \langle S, P, \leq, v \rangle \), where \( S \) is a non-empty set of states, i.e. parts of the elements in \( A \) in the foregoing epistemic modal algebra \( U \), \( P \) is the subspace of possible states where states \( s \) and \( t \) are compatible when \( s \sqcup t \in P, \leq \) is a partial order, and \( v: \text{Prop} \rightarrow (2^S \times 2^S) \) assigns a bilateral proposition \( \langle p^+, p^- \rangle \) to each atom \( p \in \text{Prop} \) with \( p^+ \) and \( p^- \) incompatible (Hawke and Özgün, forthcoming: 10-11). Exact verification \( (\vdash) \) and exact falsification
(-) are recursively defined as follows (Fine, 2017a: 19; Hawke and Özgün, forthcoming: 11):

\[ s \vdash p \text{ if } s \in \llbracket p \rrbracket^+ \]

(s verifies p, if s is a truthmaker for p i.e. if s is in p's extension);

\[ s \dashv p \text{ if } s \in \llbracket p \rrbracket^- \]

(s falsifies p, if s is a falsifier for p i.e. if s is in p's anti-extension);

\[ s \vdash \neg p \text{ if } s \vdash p \]

(s verifies not p, if s falsifies p);

\[ s \dashv \neg p \text{ if } s \dashv p \]

(s falsifies not p, if s verifies p);

\[ s \vdash p \land q \text{ if } \exists t,u, t \vdash p, u \vdash q, \text{ and } s = t \sqcup u \]

(s verifies p and q, if s is the fusion of states, t and u, t verifies p, and u verifies q);

\[ s \dashv p \land q \text{ if } s \dashv p \text{ or } s \dashv q \]

(s falsifies p and q, if s falsifies p or s falsifies q);

\[ s \vdash p \lor q \text{ if } s \vdash p \text{ or } s \vdash q \]

(s verifies p or q, if s verifies p or s verifies q);

\[ s \dashv p \lor q \text{ if } \exists t,u, t \vdash p, u \vdash q, \text{ and } s = t \sqcup u \]

(s falsifies p or q, if s is the state overlapping the states t and u, t falsifies p, and u falsifies q);

\[ s \text{ exactly verifies } p \text{ if and only if } s \vdash p \text{ if } s \in \llbracket p \rrbracket; \]

\[ s \text{ inexactly verifies } p \text{ if and only if } s \triangleright p \text{ if } \exists s' \subseteq S, s' \vdash p; \text{ and} \]

\[ s \text{ loosely verifies } p \text{ if and only if, } \forall t, s. s \sqcup t, s \sqcup t \vdash p \text{ (35-36);} \]

\[ s \vdash A \phi \text{ if and only if for all } t \in P \text{ there is a } t' \in P \text{ such that } t' \sqcup t \in P \text{ and } t' \vdash \phi; \]

\[ s \vdash A \phi \text{ if and only if there is a } t \in P \text{ such that for all } u \in P \text{ either } t \sqcup u \notin P \text{ or } u \vdash \phi, \text{ where } A \phi \text{ denotes the apriority of } \phi. \]

Epistemic (primary), subjunctive (secondary), and 2D hyperintensions can be defined as follows, where hyperintensions are functions from states to extensions, and intensions are functions from worlds to extensions:

- Epistemic Hyperintension:
  \[ \text{pri}(x) = \lambda s. [x]^{s,s}, \text{ with } s \text{ a state in the state space defined over the foregoing epistemic modal algebra, } U; \]

\[ \text{Note that the clauses for apriority here tie the notion to states of information, by contrast to the proposal in Edgington (2004: 6) according to which 'a priori knowledge is independent of the state of information'.} \]

28
\begin{itemize}
\item Subjunctive Hyperintension:
\[ \text{sec}_{v@}(x) = \lambda w. [x]^{v@,w}, \] with \( w \) a state in metaphysical state space \( W \);
\item 2D-Hyperintension:
\[ 2D(x) = \lambda s \lambda w [x]^{s,w} = 1. \]
\end{itemize}

An abstraction principle for epistemic hyperintensions can be defined as follows:
For all types, \( A, B \), there is a homotopy:
\[
H := [(f \sim g) \equiv \prod_{x:A}(f(x) = g(x))], \text{ where } \\
\prod_{f:A \rightarrow B}(f \sim f) \wedge (f \sim g \rightarrow g \sim f) \wedge (f \sim g \rightarrow g \sim h \rightarrow f \sim h)],
\]
such that, via Voevodsky’s (2006) Univalence Axiom, for all type families \( A, B : U \), there is a function:
\[ \text{idtoequiv} : (A =_{U} B) \rightarrow (A \simeq B), \]
which is itself an equivalence relation:
\[ (A =_{U} B) \simeq (A \simeq B). \]

Abstraction principles for epistemic hyperintensions take, then, the form:
\[ \exists f, g[f(x) = g(x)] \simeq [f(x) \simeq g(x)]. \]

\subsection{Modal Coalgebraic Automata}

Modal coalgebraic automata can be thus characterized. Let a category \( C \) be comprised of a class \( \text{Ob}(C) \) of objects and a family of arrows for each pair of objects \( C(A,B) \) (Venema, 2007: 421). A functor from a category \( C \) to a category \( D \), \( E : C \rightarrow D \), is an operation mapping objects and arrows of \( C \) to objects and arrows of \( D \) (422). An endofunctor on \( C \) is a functor, \( E : C \rightarrow C \) (op. cit.).

A \( E \)-coalgebra is a pair \( A = (A, \mu) \), with \( A \) an object of \( C \) referred to as the carrier of \( A \), and \( \mu : A \rightarrow E(A) \) is an arrow in \( C \), referred to as the transition map of \( A \) (390).

As, further, a coalgebraic model of modal logic, \( A \) can be defined as follows (407):

For a set of formulas, \( \Phi \), let \( \nabla \Phi := \Box \lor \Phi \land \diamond \Phi \), where \( \diamond \Phi \) denotes the set \( \{ \diamond \phi \mid \phi \in \Phi \} \) (op. cit.). Then,

\footnote{See Chapter 3 for further discussion.}
\(\diamond \phi \equiv \nabla \{\phi, \top\}\),
\(\Box \phi \equiv \nabla \emptyset \lor \nabla \phi\) (op. cit.).

\([\nabla \Phi] = \{w \in W \mid R[w] \subseteq \bigcup \{\{\phi\} \mid \phi \in \Phi\} \text{ and } \forall \phi \in \Phi, \{\phi\} \cap R[w] \neq \emptyset\}\)
(Fontaine, 2010: 17).

Let an \(\mathbf{E}\)-coalgebraic modal model, \(\mathcal{A} = \langle S, \lambda, R[\cdot] \rangle\), such that \(S, s \models \nabla \Phi\) if and only if, for all (some) successors \(\sigma\) of \(s \in S\), \(\Phi, \sigma(s) \in \mathbf{E}(\models_{\lambda})\) (Venema, 2007: 407), with \(\mathbf{E}(\models_{\lambda})\) a relation lifting of the satisfaction relation \(\models_{\lambda} \subseteq S \times \Phi\). Let a functor, \(K\), be such that there is a relation \(K! \subseteq K(A) \times K(A')\) (Venema, 2012: 17)). Let \(Z\) be a binary relation s.t. \(Z \subseteq A \times A'\) and \(\wp!Z \subseteq \wp(A) \times \wp(A')\), with
\(\wp!Z := \{((X, X') \mid \forall x \in X \exists x' \in X' \text{ with } (x, x') \in Z \land \forall x' \in X' \exists x \in X \text{ with } (x, x') \in Z\}\)
(op. cit.). Then, we can define the relation lifting, \(K!\), as follows:
\(K! := \{((\pi, X), (\pi', X')) \mid \pi = \pi' \text{ and } (X, X') \in \wp!Z\}\) (op. cit.).

A coalgebraic model of deterministic automata can finally be thus defined (Venema, 2007: 391). An automaton is a tuple, \(\mathcal{A} = \langle A, a_I, C, \delta, F \rangle\), such that \(A\) is the state space of the automaton \(\mathcal{A}\); \(a_I \in A\) is the automaton’s initial state; \(C\) is the coding for the automaton’s alphabet, mapping numerals to properties of the natural numbers; \(\delta: A \times C \rightarrow A\) is a transition function, and \(F \subseteq A\) is the collection of admissible states, where \(F\) maps \(A\) to \(\{1, 0\}\), such that \(F: A \rightarrow 1\) if \(a \in F\) and \(A \rightarrow 0\) if \(a \notin F\) (op. cit.).

Modal automata are defined over a modal one-step language (Fontaine and Venema, 2018: 3.1-3.2; Venema, 2020: 7.2). With \(A\) being a set of propositional variables the set, \(\text{Latt}(X)\), of lattice terms over \(X\) has the following grammar:

\[\pi ::= \bot \mid \top \mid x \mid \pi \land \pi \mid \pi \lor \pi,\]

with \(x \in X\) and \(\pi \in \text{Latt}(A)\) (op. cit.).

The set, \(\text{1ML}(A)\), of modal one-step formulas over \(A\) has the following grammar:

\[\alpha \in A ::= \bot \mid \top \mid \diamond \pi \mid \Box \pi \mid \alpha \land \alpha \mid \alpha \lor \alpha\] (op. cit.).

A modal P-automaton \(\mathcal{A}\) is a triple, \(\langle A, \Theta, a_I \rangle\), with \(A\) a non-empty finite set of states, \(a_I \in A\) an initial state, and the transition map
\(\Theta: A \times \wp P \rightarrow \text{1ML}(A)\)
maps states to modal one-step formulas (op. cit.: 7.3).

The crux of the reconciliation between algebraic models of cognitivism and the formal foundations of modal expressivism is based on the duality

30
between categories of algebras and coalgebras: \( A = \langle A, \alpha: A \to E(A) \rangle \) is dual to the category of algebras over the functor \( \alpha \) (417-418). For a category \( C \), object \( A \), and endofunctor \( E \), define a new arrow, \( \alpha \), s.t. \( \alpha: EA \to A \). A homomorphism, \( f \), can further be defined between algebras \( \langle A, \alpha \rangle \), and \( \langle B, \beta \rangle \). Then, for the category of algebras, the following commutative square can be defined: (i) \( EA \to EB \ (Ef) \); (ii) \( EA \to A \ (\alpha) \); (iii) \( EB \to B \ (\beta) \); and (iv) \( A \to B \ (f) \) (cf. Hughes, 2001: 7-8). The same commutative square holds for the category of coalgebras, such that the latter are defined by inverting the direction of the morphisms in both (ii) \( [A \to EA \ (\alpha)] \), and (iii) \( [B \to EB \ (\beta)] \) (op. cit.)

The significance of the foregoing is twofold. First and foremost, the above demonstrates how a formal correspondence can be effected between algebraic models of cognition and coalgebraic models which provide a natural setting for modal logics and automata. The second aspect of the philosophical significance of modal colagebraic automata is that – as a model of modal expressivism – the proposal is able to countenance fundamental properties in the foundations of mathematics, and circumvent the issues accruing to the attempt so to do by the competing expressivist approaches.

### 2.3 Material Adequacy

The material adequacy of epistemic modal algebras as a fragment of the representational theory of mind is witnessed by the prevalence of possible worlds and hyperintensional semantics – the model theory for which is algebraic (cf. Blackburn et al., 2001: ch. 5) – in cognitive psychology and artificial intelligence.

In Bayesian perceptual psychology, e.g., the visual system is presented with a prior distribution of possibilities concerning the direction of a source of light. The set of possibilities is pointed, as the visual system calculates the likelihood that one of the possibilities is actual, and places a condition thereby on the accuracy of the attribution of properties – such as boundedness and volume – to distal particulars (cf. Mamassian et al., 2000).

In artificial intelligence, the subfield of knowledge representation draws on epistemic logic, where belief and knowledge are interpreted as necessity operators (Meyer and van der Hoeck, 1995; Fagin et al., 1995). Possibility and necessity may receive other interpretations in mental terms, such as that of conceivability and apriority (i.e. truth in all epistemic possibil-
ities, or inconceivability that not \( \phi \). The language of thought hypothesis maintains that thinking occurs in a mental language with a computational syntax and a semantics. The philosophical significance of cognitivism about epistemic modality is that it construes epistemic intensions and hyperintensions as abstract, computational functions in the mind, and thus provides an explanation of the relation that human beings bear to epistemic possibilities. Intensions and hyperintensions are semantically imbued abstract functions comprising the computational syntax of the language of thought. The functions are semantically imbued because they are defined relative to a parameter ranging over either epistemically possible worlds or epistemic states in a state space, and extensions or semantic values are defined for the functions relative to that parameter. Cognitivism about epistemic modality argues that thoughts are composed of epistemic intensions or hyperintensions. Cognitivism about epistemic modality provides a metaphysical explanation or account of the ground of thoughts, arguing that they are grounded in epistemic possibilities and either intensions or hyperintensions which are themselves internal representations comprising the syntax and semantics for a mental language. This is consistent with belief and knowledge being countenanced in an epistemic logic for artificial intelligence, as well. Epistemic possibilities are constitutively related to thoughts, and figure furthermore in the analysis of notions such as apriority and conceivability, as well as belief and knowledge in epistemic logic for artificial intelligence.

The proposal that possible worlds semantics comprises the model for thoughts and propositions is anticipated by Wittgenstein (1921/1974: 2.15-2.151, 3-3.02); Chalmers (2011); and Jackson (2011). Their approaches depart, however, from the one here examined in the following respects.

Wittgenstein (op. cit.: 1-1.1) has been interpreted as endorsing an identity theory of propositions, which does not distinguish between internal thoughts and external propositions (cf. McDowell, 1994: 27; and Hornsby, 1997: 1-3). How the identity theory of propositions is able to accommodate Wittgenstein’s suggestion that a typed hierarchy of propositions can be generated – only if the class of propositions has a general form and the sense of propositions over which operations range is invariant by being individuated by the possibilities figuring as their truth and falsity conditions (cf. Wittgenstein, 1979: 21/11/16, 23/11/16, 7/11/17; and Potter, 2009: 283-285 for detailed discussion) – is an open question. Wittgenstein (1921/1974: 5.5561) writes that 'Hierarchies are and must be independent of reality', although provides no account of how the independence can be effected.
Jackson (2008: 48-50) distinguishes between personal and subpersonal theories by the role of neural science in individuating representational states (cf. Shea, 2013, for further discussion), and argues in favor of a 'personal-level implicit theory' for the possible worlds semantics of mental representations.

Chalmers’ approach comes closest to the one here proffered, because he argues for a hybrid cognitivist-expressivist approach as well, according to which epistemic intensions – i.e. functions from epistemically possible worlds to extensions – are individuated by their inferential roles (2012a: 462-463). Chalmers endorses what he refers to as 'anchored inferentialism', and in particular 'acquaintance inferentialism' for intensions, according to which 'there is a limited set of primitive concepts, and all other concepts are grounded in their inferential role with respect to these concepts', where 'the primitive concepts are acquaintance concepts' (463, 466) and '[a]cquaintance concepts may include phenomenal concepts and observational concepts: primitive concepts of phenomenal properties, spatiotemporal properties, and secondary qualities' (2010b: 11). According to Chalmers, 'anchored inferential role determines a primary intension. The relevant role can be seen as an internal (narrow or short-armed) role, so that the content is a narrow content' (5). The inferences in question are taken to be 'suppositional' inferences, from a base class of truths, PQTI – i.e. truths about physics, consciousness, and indexicality, and a that’s all truth – determining canonical specifications of epistemically possible worlds, to other truths (3). With regard to how suppositional inference, i.e. 'scrutability', plays a role in the definitions of intensions, Chalmers writes that '[t]he primary intension of [a sentence] S is true at a scenario [i.e. epistemically possible world] w iff D epistemically necessitates S, where D is a canonical specification of w*, where 'D epistemically necessitates S iff a conditional of the form 'D → S' is apriori' and the apriori entailment is the relation of scrutability (2006). Chalmers (2012a: 245) is explicit about this: "The intension of a sentence S (in a context) is true at a scenario w iff S is a priori scrutable from D (in that context), where D is a canonical specification of w (that is, one of the epistemically complete sentences in the equivalence class of w) ... A Priori Scrutability entails that this sentence S is a priori scrutable (for me) from a canonical specification D of my actual scenario, where D is something along the lines of PQTI". "The secondary intension of S is true at a world w iff D metaphysically necessitates S*, where 'D metaphysically necessitates S when a subjunctive conditional of the form 'if D had been the case, S would have been the case' is true" (op. cit.). Thus, suppositional inference, i.e. scrutability, determines the
intensions of two-dimensional semantics.

On my approach, intensions and hyperintensions are semantically imbued functions. Intensions and hyperintensions as functions comprise the computational syntax for the language of thought, but they are semantically imbued because they are functions from epistemic possibilities to extensions.\(^5\) This is consistent with the inferences of scrutability playing a role in the individuation of intensions and hyperintensions, but whereas Chalmers grounds inferences in dispositions (2010: 10), I claim that the inferences drawn from the canonical specifications of epistemic possibilities to arbitrary truths are apriori mental computations.

In the the remainder of the paper, I outline an expressivist semantics for epistemic modality and proffer an argument against relativism about epistemic modals. I endeavor, then, to demonstrate the advantages accruing to the present approach to countenancing modal expressivism via modal colagebraic automata, via a comparison of the theoretical strength of the proposal when applied to characterizing the fundamental properties of the foundations of mathematics, by contrast to the competing approaches to modal expressivism and the limits of their applications thereto.

### 2.4 Expressivist Semantics for Epistemic Possibility

Let expressivism about a domain of discourse be the claim that an utterance from that domain expresses a mental state, rather than states a fact (Hawke and Steinert-Threlkeld, 2021). Hawke and Steinert-Threlkeld (op. cit., 480) distinguish between semantic expressivism and pragmatic expressivism. Expressivism about epistemic possibility takes the property expressed by \(\Diamond \phi\) to be \(\{s \subseteq W: s \Vdash \neg \neg \neg p\}\), where \(s\) is a state of information, \(W\) is a set of possible worlds, and \(s \Vdash \phi\) if and only if \(\phi\) is assertible relative to \(s\), if and only if the state of information is compatible with \(\phi\) (op. cit.). Semantic expressivism incorporates a "psychologistic semantics" according to which the value of \(\phi\) is a partial function from information states to truth-values, such that "the mental type expressed by \(\phi\) is characterized in terms of the assertibility relation \(\Vdash\)" and "the definition of \(\Vdash\) is an essential part of that of \([\ ]\)" (481).

---

\(^5\)An anticipation of this proposal is Tichy (1969), who defines intensions as Turing machines.
Pragmatic expressivism rejects the psychologistic semantics condition, and "allows for a gap between the compositional semantic theory and $\vdash$" (op. cit.).

Hawke and Steinert-Threlkeld (op. cit.) argue that satisfying the following conditions is a desideratum of any expressivist account about epistemic possibility (3.5):

(Weak) Wide-scope Free Choice (WFC (3.1)):
\[ \diamond p \lor \diamond \neg p \vdash \diamond p \land \diamond \neg p \]

Disjunctive Inheritance (DIN (3.2)):
\[ (\diamond p \land q) \lor r \vdash [\diamond (p \land q) \land q] \lor r \]

Disjunctive Syllogism and Schroeder’s Constraints (3.4):
\[ \text{DSF } \{\diamond \neg q, p \lor \Box q \not \Delta p\} \]
\[ \text{SCH } \{\diamond \neg p, p \lor \Box q \not \Delta \Box q\} \]

DSF and SCH record the failure of disjunctive syllogism in the presence of epistemic contradictions.

WFC is vindicated by the contention that when someone asserts $p \lor \neg p$, they neither believe $p$ nor believe $\neg p$, and so are in a position to assert both $\diamond p$ and $\diamond \neg p$.

DIN is vindicated by the equivalence of the content of the utterances, e.g.,

(1) Nataly is at home and might be watching a film.
(2) Nataly is at home and might be watching a film at home (3.2).

Hawke and Steinert-Threlkeld’s modal propositional assertibility semantics is then as follows (5.1).

Reading $t \subseteq s$: $[\phi]^t \neq 1$ as *s refutes $\phi^s$:

- if $p$ is an atom: $[p]^s = 1$ iff $s \subseteq V(p)$
  if $p$ is an atom $[p]^s = 0$ iff $s$ refutes $p$

- $[\neg \phi]^s = 1$ iff $[\phi]^s = 0$
- $[\neg \phi]^s = 0$ iff $[\phi]^s = 1$

- $[\phi \land \psi]^s = 1$ iff $[\phi]^s = 1$ and $[\psi]^s = 1$
- $[\phi \land \psi]^s = 0$ iff $s$ refutes $\phi \land \psi$
Unlike Yalcın’s (2007) domain semantics (4.1), Veltman’s (1996) update semantics (4.2), and Moss’ (2015, 2018) probabilistic semantic expressivism (6.2), Hawke and Steinert-Threlkeld’s assertibility semantics satisfies WFC, DIN, DSF, and SCH (Hawke and Steinert-Threlkeld, 2020: 507). As a preliminary, suppose

**Proposition 1** If \( \varphi \) is \( \diamond \)-free, then \( s \models \diamond \varphi \) holds iff there exists \( w \in s \) such that: \( \{w\} \models \varphi \) (op. cit.).

Proof: \( s \models \diamond \varphi \) holds iff \( \llbracket \varphi \rrbracket^s \neq 0 \). \( \llbracket \varphi \rrbracket^s = 0 \) iff \( \llbracket \varphi \rrbracket^w = 0 \) for every \( w \in s \).

So, \( \llbracket \varphi \rrbracket^s \neq 0 \) iff \( \llbracket \varphi \rrbracket^w \neq 0 \) for some \( w \in s \) iff \( \{w\} \models \varphi \) for some \( w \in s \) (op. cit.).

For WFC, suppose that \( s \models \diamond p \lor \diamond \neg p \). So, there exists \( s_1, s_2 \) that cover \( s \) and \( s_1 \models \diamond p \) and \( s_2 \models \diamond \neg p \). By Proposition 1, there exist \( u, v \in s \) such that \( \{u\} \models p \) and \( \{v\} \models \neg p \). Thus, \( s \models \diamond p \) and \( s \models \diamond \neg p \) (op. cit.).
For DIN, suppose that \( s \models (\varnothing p \land q) \lor r \). So, there exists \( s_1, s_2 \), such that \( s = s_1 \cup s_2 \) with \( s_1 \models \varnothing p, s_1 \models q \), and \( s_2 \models r \). For every \( w \in s_1 \), \( \{w\} \models q \). There also exists \( u \in s_1 \) such that \( \{u\} \models p \). Hence, \( \{u\} \models p \land q \) and – by Proposition 1 – \( s_1 \models \varnothing(p \land q) \). Thus \( s \models [\varnothing(p \land q) \land q] \lor r \) (op. cit.).

For DSF and SCH, suppose that there is an \( s \) such that every world in \( s \) is either a \( p \land \neg q \) world or a \( \neg p \land q \) world. Suppose that there exists at least one \( p \land \neg q \) world in \( s \) and at least one \( \neg p \land q \) world in \( s \) (op. cit.).

Relativists about epistemic modals either relativize content or relativize truth to a context of assessment (Starr, 2012: 3; Egan and Weatherson, 2011: 11-14). According to content relativism, epistemic modals express different propositions in different contexts of assessment (Starr, op. cit.). According to truth relativism, epistemic modals express the same proposition, which is true relative to some assessors and false relative to others, such that truth is a three-place relation between a world, a judge, and a proposition, i.e. a centered world and a proposition (Starr, op. cit.: 3, 5). Thus, \( X \) believes that stealing is wrong is an ascription of belief in a centered proposition, i.e. a de se belief (Beddor, 2019). As Egan and Weatherson (2011: 14-15, 17) and Yalcin (2011: 307) point out, utterances with epistemic modals on the truth relativist proposal thus express second-order states (cf. Beddor, op. cit.).

That epistemic modal beliefs are second-order on the truth relativist proposal adduces against the merits of the view. Yalcin (op. cit.: 308) argues that non-human animals can entertain states expressed by epistemic modals, and we here follow him in thinking that, by taking epistemic modal beliefs to be second-order de se ascriptions, the truth relativist proposal would preclude young children and non-human animals from entertaining epistemic possibilities. However, young children and non-human animals, while lacking the capacity to entertain second-order states, nevertheless entertain epistemic possibilities. The foregoing thus adduces in favor of the expressivist proposal that epistemic modals express first-order states of mind.

2.5 Modal Expressivism and the Philosophy of Mathematics

When modal expressivism is modeled via speech acts on a common ground of presuppositions, the application thereof to the foundations of mathematics
is limited by the manner in which necessary propositions are characterized. Because for example a proposition is taken, according to the proposal, to be identical to a set of possible worlds, all necessarily true mathematical formulas can only express a single proposition; namely, the set of all possible worlds (cf. Stalnaker, 1978; 2003: 51). Thus, although distinct set-forming operations will be codified by distinct axioms of a language of set theory, the axioms will be assumed to express the same proposition: The axiom of Pairing in set theory – which states that a unique set can be formed by combining an element from each of two extant sets: \( \exists x \forall u(u \in x \iff u = a \lor u = b) \) – will be supposed to express the same proposition as the Power Set axiom – which states that a set can be formed by taking the set of all subsets of an extant set: \( \exists x \forall u(u \in x \iff u \subseteq a) \). However, that distinct operations – i.e., the formation of a set by selecting elements from two extant sets, by contrast to forming a set by collecting all of the subsets of a single extant set – are characterized by the different axioms is readily apparent. As Williamson (2016: 244) writes: "...if one follows Robert Stalnaker in treating a proposition as the set of (metaphysically) possible worlds at which it is true, then all true mathematical formulas literally express the same proposition, the set of all possible worlds, since all true mathematical formulas literally express necessary truths. It is therefore trivial that if one true mathematical proposition is absolutely provable, they all are. Indeed, if you already know one true mathematical proposition (that 2 + 2 = 4, for example), you thereby already know them all. Stalnaker suggests that what mathematicians really learn are in effect new contingent truths about which mathematical formulas we use to express the one necessary truth, but his view faces grave internal problems, and the conception of the content of mathematical knowledge as contingent and metalinguistic is in any case grossly implausible."

Thomasson (2007) argues for a version of modal expressivism which she refers to as 'modal normativism', according to which alethic modalities are

---

7See chapter 10 for an application of the epistemic interpretation of two-dimensional semantics to account for the modal profile of Orey sentences; i.e. mathematical propositions that are undecidable relative to the axioms of a given language. (For the origins of two-dimensional intensional semantics, see Kamp, 1967; Vlach, 1973; and Segerberg, 1973.) The distinction between epistemic and metaphysical possibilities, as they pertain to the values of mathematical formulas, is anticipated by Gödel's (1951: 11-12) distinction between mathematics in its subjective and objective senses, where the former targets all "demonstrable mathematical propositions", and the latter includes 'all true mathematical propositions".
to be replaced by deontic modalities taking the form of object-language, modal indicative conditionals (op. cit.: 136, 138, 141). The modal indicative conditionals serve to express constitutive rules pertaining, e.g., to ontological dependencies which state that: 'Necessarily, if an entity satisfying a property exists then a distinct entity satisfying a property exists' (143-144), and generalizes to other expressions, such as analytic conditionals which state, e.g., that: 'Necessarily, if an entity satisfies a property, such as being a bachelor, then the entity satisfies a distinct yet co-extensive property, such as being unmarried' (148). A virtue of Thomasson’s interpretation of modal indicative conditionals as expressing both analytic and ontological dependencies is that it would appear to converge with the 'If-thenist' proposal in the philosophy of mathematics. 'If-thenism' is an approach according to which, if an axiomatized mathematical language is consistent, then (i) one can either bear epistemic attitudes, such as fictive acceptance, toward the target system (cf. Leng, 2010: 180) or (ii) the system (possibly) exists [cf. Russell (op. cit.: §1); Hilbert (1899/1980: 39); Menger (1930/1979: 57); Putnam (1967); Shapiro (2000: 95); Chihara (2004: Ch. 10); and Awodey (2004: 60-61)].

However, there are at least two issues for the modal normativist approach in the setting of the philosophy of mathematics. One general issue for the proposal is that the treatment of quantification remains unaddressed, given that there are translations from modal operators, such as figure in modal indicatives, into existential and universal quantifiers.

---

8See Leng (2009), for further discussion. Field (1980/2016: 11-21; 1989: 54-65, 240-241) argues in favor of the stronger notion of conservativeness, according to which consistent mathematical theories must be satisfiable by internally consistent theories of physics. More generally, for a class of assertions, A, comprising a theory of fundamental physics, and a class of sentences comprising a mathematical language, M, any sentences derivable from A+M ought to be derivable from A alone. Another variation on the 'If-thenist' proposal is witnessed in Field (2001: 333-338), who argues that the existence of consistent forcing extensions of set-theoretic ground models adduces in favor of there being a set-theoretic pluriverse, and thus entrains indeterminacy in the truth-values of undecidable sentences. For a similar proposal, which emphasizes the epistemic role of examining how instances of undecidable sentences obtain and fail so to do relative to forcing extensions in the set-theoretic pluriverse, see Hamkins (2012: §7).

9The formal correspondence between modalities and quantifiers is anticipated by Aristotle (De Interpretatione, 9; De Caelo, I.12), who defines the metaphysical necessity of a proposition as its being true at all times. For detailed discussion of Aristotle’s theory, see Waterlow (1982). For a contemporary account of the multi-modal logic for metaphysical and temporal modalities, see Dorr and Goodman (2019). For contemporary accounts of the correspondence between modal operators and quantifiers see von Wright (1952/1957);
the normative indicative conditional approach is that Thomasson’s normative modalities are unimodal. They are thus not sufficiently fine-grained to capture distinctions such as Gödel’s (op. cit.) between mathematics in its subjective and objective senses. Further distinctions between the types of mathematical modality can be delineated which permit epistemic types of mathematical possibility to serve as a guide as to whether a formula is metaphysically mathematically possible. The convergence between epistemic and metaphysical mathematical modalities can be countenanced via a two-dimensional intensional semantics. Thus, by eschewing alethic modalities for unimodal, normative indicatives, the normative modalities are unable to account for the relation between the alethic interpretation of modality and, e.g., logical mathematical modalities treated as consistency operators on languages (cf. Field, 1989: 249-250, 257-260; Leng: 2007; 2010: 258), or for the convergence between epistemic possibilities concerning decidability and their bearing on the metaphysical modal status of undecidable sentences.

According, finally, to Brandom’s (op. cit.) modal expressivist approach, terms are individuated by their rules of inference, where the rules are taken to have a modal profile translatable into the counterfactual forms taken by the transition functions of automata (cf. Brandom, 2008: 142). In order to countenance the metasemantic truth-conditions for the object-level, pragmatic abilities captured by the automata’s counterfactual transition states, Brandom augments a first-order language comprised of a stock of atomic formulas with an incompatibility function (141). An incompatibility function, \( I \), is defined as the incoherence of the union of two sentences, where incoherence is a generalization of the notion of inconsistency to nonlogical vocabulary.

\[
x \cup y \in Inc \iff x \in I(y) \quad (141-142).
\]

Incompatibility is supposed to be a modal notion, such that the union of the two sentences is incompossible (126). A sentence, \( \beta \), is an incompatibility-consequence, \( \models I \), of a sentence, \( \alpha \), if there is no sequence of sentences, \(<\gamma_1, \ldots, \gamma_n>\), such that it can be the case that \( \alpha \models I <\gamma_1, \ldots, \gamma_n> \), yet not

---

\( ^{10} \)See chapters 8 and 10 for further discussion. A precedent is Reinhardt (1974: 199-200), who proposes the use of imaginary sets, classes, and projections, as ‘imaginary experiments’ (204), in order to ascertain the consequences of accepting new axioms for ZF which might account for the reduction of the incompleteness of Orey sentences. See Maddy (1988,b), for critical discussion.
be the case that $\beta \models_I \langle \gamma_1, \ldots, \gamma_n \rangle$ (125). To be incompatible with a necessary formula is to be compatible with everything that does not entail the formula (129-130). Dually, to be incompatible with a possible formula is to be incompatible with everything compatible with something compatible with the formula (op. cit.).

There are at least two, general issues for the application of Brandom’s modal expressivism to the foundations of mathematics.

The first issue is that the mathematical vocabulary – e.g., the set-membership relation, $\in$ – is axiomatically defined. I.e., the membership relation is defined by, inter alia, the Pairing and Power Set axioms of set-theoretic languages. Thus, mathematical terms have their extensions individuated by the axioms of the language, rather than via a set of inference rules that can be specified in the absence of the mention of truth values. Even, furthermore, if one were to avail of modal notions in order to countenance the intensions of the mathematical vocabulary at issue – i.e., functions from terms in intensional contexts to their extensions – the modal profile of the intensions is orthogonal to the properties encoded by the incompatibility function. Fine (2006) avails, e.g., of a dynamic logic in order to countenance the possibility of reinterpreting the intensions at issue, and of thus accounting for variance in the range of the domains of quantifier expressions. The dynamic possibilities are specified as operational conditions on tracking increases in the size of the cardinality of the universe (Fine, 2005). Uzquiano (2015b) argues that it is always possible to reinterpret the intensions of non-logical vocabulary, as one augments one’s language with stronger axioms of infinity and climbs thereby farther up the cumulative hierarchy of sets. The reinterpretations of, e.g., the concept of set are effected by the addition of new large cardinal axioms, which stipulate the existence of larger inaccessible cardinals. However, it is unclear how the incompatibility function – i.e., a modal operator defined via Boolean negation and a generalized condition on inconsistency – might similarly be able to model the intensions pertaining to the ontological expansion of the cumulative hierarchy.

The second issue is that Brandom’s inferential expressivist semantics is not compositional (Brandom, 2008: 135-136). While the formulas of the semantics are recursively formed – because the decomposition of complex formulas into atomic formulas is decidable$^{11}$ – formulas in the language are

---

$^{11}$Let a decision problem be a propositional function which is feasibly decidable, if it is a member of the polynomial time complexity class; i.e., if it can be calculated as a polynomial
not compositional, because they fail to satisfy the subformula property to the effect that the value of a logically complex formula is calculated as a function of the values of the component logical connectives applied to subformulas therein (op. cit.)\(^{12}\).

By contrast to the limits of Brandom’s approach to modal expressivism, modal coalgebraic automata can circumvent both of the issues mentioned in the foregoing. In response to the first issue, concerning the axiomatic individuation and intensional profiles of mathematical terms, functors of modal coalgebraic automata can be interpreted in order to provide a precise delineation of the intensions of the target vocabulary [cf. Author (ms)]. In response, finally, to the second of the above issues, the values taken by modal coalgebraic automata are both decidable and computationally feasible, while the duality of colagebras to Boolean-valued models of modal algebras ensures that the formulas therein retain their compositionality. The decidability of colagebraic automata can further be witnessed by the role of modal coalgebras in countenancing the modal profile of Ω-logical consequence, where – given a proper class of Woodin cardinals – the values of mathematical formulas can remain invariant throughout extensions of the ground models comprising the set-theoretic pluriverse (cf. Woodin, 2010; and chapter 9). The individuation of large cardinals can further be characterized by the functors of modal coalgebraic automata, when the latter are interpreted so as to countenance the elementary embeddings constitutive of large cardinal axioms in the category of sets.

### 2.6 Concluding Remarks

In this essay, I have endeavored to account for a mathematically tractable background against which to model both modal cognitivism and modal expressivism. I availed, to that end, of the duality between epistemic modal algebras and modal coalgebraic automata. Epistemic modal algebras were shown to comprise a materially adequate fragment of the language of thought, function of the size of the formula’s input [see Dean (2015) for further discussion].

\(^{12}\)Note that Incurvati and Schlöder (2020) advance a multilateral inferential expressivist semantics for epistemic modality which satisfies the subformula property. (Thanks here to Luca Incurvati.) Incurvati and Schlöder (2021) extend the semantics to normative vocabulary, but it is an open question whether the semantics is adequate for mathematical vocabulary as well.
given that models thereof figure in both cognitive psychology and artificial intelligence. It was then shown how the approach to modal expressivism here proffered, as regimented by the modal coalgebraic automata to which the epistemic modal algebras are dual, avoids the pitfalls attending to the competing modal expressivist approaches based upon both the inferentialist approach to concept-individuation and the approach to codifying the speech acts in natural language via intensional semantics. The present modal expressivist approach was shown, e.g., to avoid the limits of the foregoing in the philosophy of language, as they concerned the status of necessary propositions; the inapplicability of inferentialist-individuation to mathematical vocabulary; and failures of compositionality. Countenancing modal expressivism via modal coalgebraic automata was shown, then, to be able to account for both the intensions of mathematical terms and possible reinterpretations thereof; for the modal profile of \( \Omega \)-logical consequence in the category of sets; and for the elementary embeddings constitutive of large cardinal axioms in set-theoretic languages.
Chapter 3

Cognitivism about Epistemic Modality

3.1 Introduction

This essay aims to vindicate the thesis that cognitive computational properties are abstract objects implemented in physical systems.¹ A recent approach to the foundations of mathematics is Homotopy Type Theory.² In Homotopy Type Theory, homotopies can be defined as equivalence relations on intensional functions. In this essay, I argue that homotopies can thereby figure in abstraction principles for epistemic intensions, i.e. functions from epistemically possible worlds to extensions.³ Homotopies for epistemic intensions thus comprise identity criteria for some cognitive mechanisms. The philosophical significance of the foregoing is twofold. First, the proposal demonstrates how epistemic modality is a viable candidate for a fragment of the language of thought.⁴ Second, the proposal serves to delineate one

¹Cf. Turing (1950); Putnam (1967b); Newell (1973); Fodor (1975); and Pylyshyn (1978).
²Cf. The Univalent Foundations Program (2013).
³For the first proposal to the effect that abstraction principles can be used to define abstracta such as cardinal number, see Frege (1884/1980: 68; 1893/2013: 20). For the locus classicus of the contemporary abstractionist program, see Hale and Wright (2001).
⁴Given a metalanguage, a precedent to the current approach – which models thoughts and internal representations via possible worlds model theory – can be found in Wittgenstein (1921/1974: 2.15-2.151, 3-3.02).
conduit for our epistemic access to epistemic intensions as abstract objects.\footnote{The proposal that epistemic intensions might be sui generis abstract objects, not reducible to sets, is proffered by Chalmers (2011: 101) who writes: 'It is even possible to introduce a special sort of abstract object corresponding to these intensions. Of course these abstract objects cannot be sets of ordered pairs. But we might think of an intension formally as an abstract object which when combined with an arbitrary scenario yields a truth value (or an extension).'}

In Section 2, I provide an abstraction principle for epistemic intensions, by availing of the equivalence relations countenanced in Homotopy Type Theory. In Section 3, I describe how models of Epistemic Modal Algebra are availed of when perceptual representational states are modeled in Bayesian

Bealer (1982) proffers a non-modal algebraic logic for intensional entities – i.e., properties, relations, and propositions – which avails of a λ-definable variable-binding abstraction operator (op. cit.: 46-48, 209-210). Bealer reduces modal notions to logically necessary conditions-cum-properties, as defined in his non-modal algebraic logic (207-209). The present approach differs from the foregoing by: (i) countenancing a modal algebra, on an epistemic interpretation thereof; (ii) treating the abstraction operator as a Fregean function from concepts to objects, rather than as a λ-operator; (iii) availing of the univalence axiom in Homotopy Type Theory – which collapses identity and isomorphism – in order to provide an equivalence relation for the abstraction principle pertinent to (ii); and (iv) demonstrating how the model is availed of in various branches of the cognitive sciences, such that Epistemic Modal Algebra may be considered a viable candidate for the language of thought.

Katz (1998) proffers a view of the epistemology of abstracta, according to which the syntax and the semantics for the propositions are innate (35). Katz suggests that the proposal is consistent with both a Fregean approach to propositions, according to which they are thoughts formed by the composition of senses, and a Russellian approach, according to which they are structured tuples of non-conceptual entities (36). He endorses an account of senses according to which they are correlated to natural language sentence types (114-115).

One difference between Katz’s proposal and the one here presented is that Katz rejects modal approaches to propositions, because the latter cannot distinguish between distinct contradictions (38fn.6). Following, Lewis (1973: I.6), the present approach does not avail of impossible worlds which distinguish between distinct contradictions. For approaches to epistemic space and conceivability which do admit of impossible worlds, see Jago (2009; 2014); Berto (2014); Berto and Schoonen (2018); and Priest (2019). A second difference is that, on Katz’s approach, the necessity of mathematical truths is argued to consist in reductio proofs, such that the relevant formulas will be true on all interpretations, and thus true of logical necessity (39). However, the endeavor to reduce the necessity of mathematical truths to the necessity of logical consequence would result in the preclusion, both of cases of informal proofs in mathematics, which can, e.g., involve diagrams (cf. Azzouni, 2004; Giaquinto, 2008: 1.2), and of mathematical truths which obtain in axiomatizable, yet non-logical mathematical languages such as Euclidean geometry. Finally, Katz rejects abstraction principles, and thus implicit definitions for abstract objects (105-106).
perceptual psychology; when speech acts are modeled in natural language semantics; and when knowledge, belief, intentional action, and rational intuition are modeled in philosophical approaches to the nature of propositional attitudes. This provides abductive support for the claim that Epistemic Modal Algebra is both a compelling and materially adequate candidate for a fragment of the language of thought. In Section 4, I argue that the proposal (i) resolves objections to the relevant abstraction principles advanced by both Dean (2016) and Linnebo and Pettigrew (2014). Section 5 provides concluding remarks.

3.2 An Abstraction Principle for Epistemic Intensions

In this section, I specify a homotopic abstraction principle for intensional functions. Intensional isomorphism, as a jointly necessary and sufficient condition for the identity of intensions, is first proposed in Carnap (1947: §14). The isomorphism of two intensional structures is argued to consist in their logical, or L-, equivalence, where logical equivalence is co-extensive with the notions of both analyticity (§2) and synonymy (§15). Carnap writes that: ‘[A]n expression in S is L-equivalent to an expression in S’ if and only if the semantical rules of S and S’ together, without the use of any knowledge about (extra-linguistic) facts, suffice to show that the two have the same extension’ (p. 56), where semantical rules specify the intended interpretation of the constants and predicates of the languages (4).\(^6\) The current approach differs from Carnap’s by defining the equivalence relation necessary for an abstraction principle for epistemic intensions on Voevodsky’s (2006) Univalence Axiom, which collapses identity with isomorphism in the setting of intensional type theory.\(^7\)

---

\(^6\)For criticism of Carnap’s account of intensional isomorphism, based on Carnap’s (1937: 17) ‘Principle of Tolerance’ to the effect that pragmatic desiderata are a permissible constraint on one’s choice of logic, see Church (1954: 66-67).

\(^7\)Note further that, by contrast to Carnap’s approach, epistemic intensions are here distinguished from linguistic intensions (cf. Chapter 6, for further discussion), and the current work examines the philosophical significance of the convergence between epistemic intensions and formal, rather than natural, languages. For a translation from type theory to set theory – which is of interest to, inter alia, the definability of epistemic intensions in the setting of set theory (cf. Chapter 10, below) – see Linnebo and Rayo (2012). For

---

46
Topological Semantics

In the topological semantics for modal logic, a frame is comprised of a set of points in topological space, a domain of propositions, and an accessibility relation:

\[ F = (X,R); \]

\[ X = \{ x \in X \}; \] and

\[ R = \{ R_{xy} \}_{x,y \in X} \] iff \( R_x \subseteq X \times X \), s.t. if \( R_{xy} \), then \( \exists o \subseteq X \), with \( x \in o \) s.t. \( \forall y \in o \) and \( R_{xy} \),

where the set of points accessible from a privileged node in the space is said to be open.\(^8\) A model defined over the frame is a tuple, \( M = (F,V) \), with \( V \) a valuation function from subsets of points in \( F \) to propositional variables taking the values 0 or 1. Necessity is interpreted as an interiority operator on the space:

\[ M,x \models □φ \text{ iff } \exists o \subseteq X, \text{ with } x \in o, \text{ such that } \forall y \in o M,y \models φ. \]

Homotopy Theory

Homotopy Theory countenances the following identity, inversion, and concatenation morphisms, which are identified as continuous paths in the topology. The formal clauses, in the remainder of this section, evince how homotopic morphisms satisfy the properties of an equivalence relation.\(^9\)

\[ p : [0,1] \to X, \text{ with } p(0) = x \text{ and } p(1) = y; \]

\[ f : X_1 \to X_2; \]

\[ g : X_1 \to X_2; \]

\[ H : X_1 \times \{0,1\} \to X_2, H_{x,0} = f(x) \text{ and } H_{x,1} = g(x). \]

---

8In order to ensure that the Kripke semantics matches the topological semantics, \( X \) must further be Alexandrov; i.e., closed under arbitrary unions and intersections. Thanks here to Peter Milne.

9The definitions and proofs at issue can be found in the Univalent Foundations Program (op. cit.: ch. 2.0-2.1).
Reflexivity
\forall x,y:A \forall p(p : x =_A y) : \tau(x,y,p), with A and \tau designating types, 'x:A' interpreted as 'x is a token of type A', refl a dependent function of reflexivity, and U designating a universe of elements, e:
\forall \alpha:A \exists e(\tau(\alpha, \alpha, \text{refl}_\alpha))

The Induction Principle
If:
\forall x,y:A \forall p(p : x =_A y) \exists \tau[\tau(x,y,p)] \land \forall \alpha:A \exists e(\tau(\alpha, \alpha, \text{refl}_\alpha))
Then:
\forall x,y:A \exists p(p : x =_A y) \exists e[\text{ind}_A(\tau, e, x, y, p) : \tau(x,y,p), such that
\text{ind}_A(\tau, e, \alpha, \alpha, \text{refl}_\alpha) \equiv e(\alpha)]

Symmetry
\forall A \forall x,y:A \exists H_\Sigma(x=y \rightarrow y=x)
H_\Sigma := p \mapsto p^{-1}, such that
\forall x:A(\text{refl}_x \equiv \text{refl}_x^{-1})

Transitivity
\forall A \forall x,y:A \exists H_T(x=y \rightarrow y=z \rightarrow x=z)
H_T := p \mapsto q \mapsto p \cdot q, such that
\forall x:A[\text{refl}_x \cdot \text{refl}_x \equiv \text{refl}_x]

Homotopic Abstraction
For all type families A,B, there is a homotopy:
H := [(f \sim g) \equiv \prod_{x:A}(f(x) = g(x))], where
\prod_{f:A \rightarrow B}[(f \sim f) \land (f \sim g \rightarrow g \sim f) \land (f \sim g \rightarrow g \sim h \rightarrow f \sim h)], such that, via Voevodsky’s (2006) Univalence Axiom, for all type families A,B:U, there is a function:
idtoeqv : (A =_U B) \rightarrow (A \simeq B),
which is itself an equivalence relation:
\((A =_U B) \simeq (A \simeq B)\).

Epistemic intensions take the form,

\[ \text{pri}(x) = \lambda s.\llbracket x\rrbracket^{s,s}, \]

with \(s\) an epistemically possible state.

Abstraction principles for epistemic intensions take, then, the form:

- \(\exists f,g[f(x) = g(x)] \simeq [f(x) \simeq g(x)]\).

### 3.3 Examples in Philosophy and Cognitive Science

The material adequacy of epistemic modal algebras as a fragment of the the language of thought is witnessed by the prevalence of possible worlds semantics – the model theory for which is algebraic (cf. Blackburn et al., 2001: ch. 5) – in cognitive psychology. Possible worlds model theory is availed of in the computational theory of mind, Bayesian perceptual psychology, and natural language semantics.

Marcus (2001) argues that mental representations can be treated as algebraic rules characterizing the computation of operations on variables, where the values of a target domain for the variables are universally quantified over and the function is one-one, mapping a number of inputs to an equivalent number of outputs (35-36). Models of the above algebraic rules can be defined in both classical and weighted, connectionist systems: Both a single and multiple nodes can serve to represent the variables for a target domain (42-45). Temporal synchrony or dynamic variable-bindings are stored in short-term working memory (56-57), while information relevant to long-term variable-bindings are stored in registers (54-56). Examples of the foregoing algebraic rules on variable-binding include both the syntactic concatenation of morphemes and noun phrase reduplication in linguistics (37-39, 70-72), as well as learning algorithms (45-48). Conditions on variable-binding are further examined, including treating the binding relation between variables and values as tensor products – i.e., an application of a multiplicative axiom for variables and their values treated as vectors (53-54, 105-106). In order to account for recursively formed, complex representations, which he refers to as structured propositions, Marcus argues instead that the syntax
and semantics of such representations can be modeled via an ordered set of registers, which he refers to as 'treelets' (108).

A strengthened version of the algebraic rules on variable-binding can be accommodated in models of epistemic modal algebras, when the latter are augmented by cylindrifications, i.e., operators on the algebra simulating the treatment of quantification, and diagonal elements. By contrast to Boolean Algebras with Operators, which are propositional, cylindric algebras define first-order logics. Intuitively, valuation assignments for first-order variables are, in cylindric modal logics, treated as possible worlds of the model, while existential and universal quantifiers are replaced by, respectively, possibility and necessity operators (\(\Diamond\) and \(\Box\)) (Venema, 2013: 249). For first-order variables, \(\{v_i\mid i < \alpha\}\) with \(\alpha\) an arbitrary, fixed ordinal, \(v_i = v_j\) is replaced by a modal constant \(d_{i,j}\) (op. cit: 250). The following clauses are valid, then, for a model, \(M\), of cylindric modal logic, with \(E_{i,j}\) a monadic predicate and \(T_i\) for \(i,j < \alpha\) a dyadic predicate:

\[
\begin{align*}
M,w \models p & \iff w \in V(p); \\
M,w \models d_{i,j} & \iff w \in E_{i,j}; \\
M,w \models \Diamond_i \psi & \iff \text{there is a } v \text{ with } w \in T_i v \text{ and } M,v \models \psi \ (252).
\end{align*}
\]

Cylindric frames need further to satisfy the following axioms (op. cit.: 254):

1. \(p \rightarrow \Diamond_i p\)
2. \(p \rightarrow \Box_i \Diamond_i p\)
3. \(\Diamond_i \Diamond_i p \rightarrow \Diamond_i p\)
4. \(\Diamond_i \Diamond_j p \rightarrow \Diamond_j \Diamond_i p\)
5. \(d_{i,i}\)
6. \(\Diamond_i(d_{i,j} \land p) \rightarrow \Box_i(d_{i,j} \rightarrow p)\)

[Translating the diagonal element and cylindric (modal) operator into, respectively, monadic and dyadic predicates and universal quantification: \(\forall xyz[(T_i xy \land E_{i,j} y \land T_i xz \land E_{i,j} z) \rightarrow y = z]\) (op. cit.)]

7. \(d_{i,j} \iff \Diamond_k(d_{i,k} \land d_{k,j})\).

Finally, a cylindric modal algebra of dimension \(\alpha\) is an algebra, \(A = \langle A, +, \cdot, -, 0, 1, \Diamond_i, d_{ij}\rangle_{i,j<\alpha}\), where \(\Diamond_i\) is a unary operator which is normal (\(\Diamond_i 0 = 0\)) and additive \([\Diamond_i(x + y) = \Diamond_i x + \Diamond_i y]\) (257).

The philosophical interest of cylindric modal algebras to Marcus’ cognitive models of algebraic variable-binding is that variable substitution is

---

\(^{10}\)See Henkin et al (op. cit.: 162-163) for the introduction of cylindric algebras, and for the axioms governing the cylindrification operators.
treated in the modal algebras as a modal relation, while universal quantification is interpreted as necessitation. The interest of translating universal generalization into operations of epistemic necessitation is, finally, that – by identifying epistemic necessity with apriority – both the algebraic rules for variable-binding and the recursive formation of structured propositions can be seen as operations, the implicit knowledge of which is apriori.

In Bayesian perceptual psychology, the problem of underdetermination is resolved by availing of a gradational possible worlds model. The visual system is presented with a set of possibilities with regard, e.g., to the direction of a light source. So, for example, the direction of light might be originating from above, or it might be originating from below. The visual system computes the constancy, i.e. the likelihood that one of the possibilities is actual.\textsuperscript{11} The computation of the perceptual constancy is an unconscious statistical inference, as anticipated by Helmholtz’s (1878) conjecture.\textsuperscript{12} The constancy places, then, a condition on the accuracy of the attribution of properties – such as boundedness and volume – to distal particulars.\textsuperscript{13}

In the program of natural language semantics in empirical and philosophical linguistics, the common ground or 'context set' is the set of possibilities presupposed by a community of speakers.\textsuperscript{14} Kratzer (1979: 121) refers to cases in which the above possibilities are epistemic as an 'epistemic conversational background', where the epistemic possibilities are a subset of objective or circumstantial possibilities (op. cit.). Modal operators are then defined on the space, encoding the effects of various speech acts in entraining updates on the context set.\textsuperscript{15} So, e.g., assertion is argued to provide a truth-conditional update on the context set, whereas there are operator updates, the effects of which are not straightforwardly truth-conditional and whose semantic values must then be defined relative to an array of intensional parameters (including

\textsuperscript{11}Cf. Mamassian et al. (2002).
\textsuperscript{12}For the history of the integration of algorithms and computational modeling into contemporary visual psychology, see Johnson-Laird (2004).
\textsuperscript{13}Cf. Burge (2010), and Rescorla (2013), for further discussion. A distinction ought to be drawn between unconscious perceptual representational states – as targeted in Burge (op. cit.) – and the inquiry into whether the properties of phenomenal consciousness have accuracy-conditions – where phenomenal properties are broadly construed, so as to include, e.g., color-phenomenal properties, as well as the property of being aware of one’s perceptual states.
\textsuperscript{14}Cf. Stalnaker (1978).
\textsuperscript{15}Cf. Kratzer (op. cit.); Stalnaker (op. cit.); Lewis (1980); Heim (1992); Veltman (1996); von Fintel and Heim (2011); and Yalcin (2012).
Finally, Epistemic Modal Algebra, as a fragment of the language of thought, is able to delineate the fundamental structure of the propositional attitudes targeted in 20th century philosophy; notably knowledge, belief, intentional action, and rational intuition. In Chapter 14, I argue, e.g., that the types of intention – acting intentionally; referring to an intention as an explanation for one’s course of action; and intending to pursue a course of action in the future – can be modeled as modal operators, whose semantic values are defined relative to an array of intensional parameters. E.g., an agent can be said to act intentionally iff her 'intention-in-action' receives a positive semantic value, where a necessary condition on the latter is that there is at least one world in her epistemic modal space at which – relative to a context of a particular time and location, which constrains the admissibility of her possible actions as defined at a first index, and which subsequently constrains the outcome thereof as defined at a second index – the intention is realized:

\[ \llbracket \text{Intenton-in-Action}(\phi) \rrbracket_w = 1 \text{ only if } \exists w' \llbracket \phi \rrbracket_{w',c(\ell,t),a,o} = 1. \]

The agent’s intention to pursue a course of action at a future time – i.e., her 'intention-for-the-future' – can receive a positive value only if there is a possible world and a future time, relative to which the possibility that a state, \( \phi \), is realized can be defined. Thus:

\[ \llbracket \text{Intention-for-the-future}(\phi) \rrbracket_w = 1 \text{ only if } \exists w' \forall t \exists t'[t < t' \land \llbracket \phi \rrbracket_{w',t'} = 1]. \]

In the setting of epistemic logic, epistemic necessity can further be modeled in a relational semantics encoding the property of knowledge, whereas epistemic possibility might encode the property of belief (cf. Hintikka, 1962; Fagin et al., 1995; Meyer and van der Hoek, 1995; Williamson, 2009). Finally, in Chapter 12, I treat Gödel’s (1953) conception of rational propositional intuition as a modal operator in the setting of a bimodal, dynamic provability logic, and demonstrates how – via correspondence theory – the notion of 'intuition-of', i.e. a property of awareness of one’s cognitive states, can be shown to be formally equivalent to the notion of 'intuition-that', i.e. a modal operator concerning the value of the propositional state at issue.\(^{16}\)

\(^{16}\)The correspondence results between modal propositional and first-order logic are advanced in van Benthem (1983; 1984/2003) and Janin and Walukiewicz (1996). Availing of correspondence theory in order to account for the relationship between the notions of 'intuition-of' and 'intuition-that' resolves an inquiry posed by Parsons (1993: 233). As a dynamic interpretational modality, rational intuition can further serve as a guide to possible reinterpretations both of quantifier domains (cf. Fine, 2006) and of the extensions
3.4 Objections and Replies

Dean (2016) raises two issues for a proposal similar to the foregoing, namely that algorithms – broadly construed – can be defined via abstraction principles which specify equivalence relations between implementations of computational properties in isomorphic machines.\textsuperscript{17} Dean’s candidate abstraction principle for algorithms as abstracts is: that the algorithm implemented by $M_1 = \text{the algorithm implemented by } M_2$ iff $M_1 \simeq M_2$.\textsuperscript{18} Both issues target the uniqueness of the algorithm purported to be identified by the abstraction principle.

The first issue generalizes Benacerraf’s (1965) contention that, in the reduction of number theory to set theory, there must be, and is not, a principled reason for which to prefer the identification of natural numbers with von Neumann ordinals (e.g., $2 = \{\emptyset, \{\emptyset\}\}$), rather than with Zermelo ordinals (i.e., order-types of well-orderings).\textsuperscript{19} The issue is evinced by the choice of whether to define algorithms as isomorphic iterations of state transition functions (cf. Gurevich, 1999), or to define them as isomorphic recursions of functions which assign values to a partially ordered set of elements (cf. Moschovakis, op. cit.). Linnebo and Pettigrew (2014: 10) argue similarly that, for two ‘non-rigid’ structures which admit of non-trivial automorphisms, one can define a graph which belies their isomorphism. E.g., let an abstraction principle be defined for the isomorphism between $S$ and $S^*$, such that

$$\forall S, S^* [A_S = A_{S^*} \text{ iff } \langle S, R_1 \ldots R_n \rangle \simeq \langle S^*, R^*_1 \ldots R^*_n \rangle].$$

However, if there is a graph, $G$, such that:

$S = \{v_1, v_2\}$, and $R = \{\langle v_1, v_2 \rangle, \langle v_2, v_1 \rangle\}$,

of mathematical vocabulary such as the membership-relation (cf. Uzquiano, 2015a). This provides an account of Gödel’s (op. cit.; 1961) suggestion that rational intuition can serve as a guide to conceptual elucidation.

\textsuperscript{17}Fodor (2000: 105, n.4) and Piccinini (2004) note that the identification of mental states with their functional roles ought to be distinguished from identifying those functional roles with abstract computations. Conversely, a computational theory of mind need not be committed to the identification of abstract, computational operations with the functional organization of a machine. Identifying abstract computational properties with the functional organization of a creature’s mental states is thus a choice point, in theories of the nature of mental representation.

\textsuperscript{18}Cf. Moschovakis (1998).

\textsuperscript{19}Cf. Zermelo (1908/1967) and von Neumann (1923/1967). Well-orderings are irreflexive, transitive, binary relations on all non-empty sets, which define a least or distinguished element in the sets.
then one can define an automorphism, $f : G \simeq G$, such that $f(v_1) = v_2$ and $f(v_2) = v_1$, such that $S^* = \{v_1\}$ while $R^* = \{\langle v^*_1, v^*_1 \rangle \}$. Then $S^*$ has one element via the automorphism, while $S$ has two. So, $S$ and $S^*$ are not, after all, isomorphic.

The second issue is that complexity is crucial to the identity criteria of algorithms. Two algorithms might be isomorphic, while the decidability of one algorithm is proportional to a deterministic \textit{polynomial} function of the size of its input – with $k$ a member of the natural numbers, $N$, and \textsc{Time} referring to the relevant complexity class: $\bigcup_{k \in N} \text{\textsc{Time}}(n^k)$ – and the decidability of the second algorithm will be proportional to a deterministic \textit{exponential} function of the size of its input – $\bigcup_{k \in N} \text{\textsc{Time}}(2^{n^k})$. The deterministic polynomial time complexity class is a subclass of the deterministic exponential time complexity class. However, there are problems decidable by algorithms only in polynomial time (e.g., the problem of primality testing, such that, for any two natural numbers, the numbers possess a greatest common divisor equal to 1), and only in exponential time (familiarly from logic, e.g., the problem of satisfiability – i.e., whether, for a given formula, there exists a model which can validate it – and the problem of validity – i.e. whether a satisfiable formula is valid).\textsuperscript{20}

Both issues can be treated by noting that Dean’s discussion targets abstraction principles for the very notion of a computable function, rather than for abstraction principles for cognitive computational properties. It is a virtue of homotopic abstraction principles for cognitive intensional functions that both the temporal complexity class to which the functions belong, and the applications of the model, are subject to variation. Variance in the cognitive roles, for which Epistemic Modal Algebra provides a model, will crucially bear on the nature of the representational properties unique to the interpretation of the intensional functions at issue. Thus, e.g., when the internal representations in the language of thought – as modeled by Epistemic Modal Algebra – subserve perceptual representational states, then their contents will be individuated by both the computational constancies at issue and the external, environmental properties – e.g., the properties of lightness and distance – of the perceiver.\textsuperscript{21} A further virtue of the foregoing is that variance in the coding of Epistemic Modal Algebras – i.e. in the types of informa-

\textsuperscript{20}For further discussion, see Dean (2015).
\textsuperscript{21}The computational properties at issue can also be defined over non-propositional information states, such as cognitive maps possessed of geometric rather than logical structure. See, e.g., O’Keefe and Nadel (1978); Camp (2007); and Rescorla (2009).
tion over which the intensional functions will be defined – by contrast to a restriction of the language of thought to mathematical languages such as Peano arithmetic, permits homotopic abstraction principles to circumvent the Burali-Forti paradox for implicit definitions based on isomorphism.\textsuperscript{22}

The examples of instances of Epistemic Modal Algebra – witnessed by the possible worlds models in Bayesian perceptual psychology, linguistics, and philosophy of mind – provide abductive support for the existence of the intensional functions specified in homotopic abstraction principles. The philosophical significance of independent, abductive support for the existence of epistemic modalities in the philosophy of mind and cognitive science is that the latter permits a circumvention of the objections to the abstractionist foundations of number theory that have accrued since its contemporary founding (cf. Wright, 1983). Eklund (2006) suggests, e.g., that the existence of the abstract objects which are the referents of numerical term-forming operators might need to be secured, prior to assuming that the abstraction principle for cardinal number is true. While Hale and Wright (2009) maintain, in response, that the truth of the relevant principles will be prior to the inquiry into whether the terms defined therein refer, they provide a preliminary endorsement of an ’abundant’ conception of properties, according to which identifying the sense of a predicate will be sufficient for predicate reference.\textsuperscript{23} One aspect of the significance of empirical and philosophical instances of models of Epistemic Modal Algebra is thus that, by providing independent, abductive support for the truth of the homotopic abstraction principles for epistemic intensions, the proposal remains neutral on the status of ’sparse’ versus ’abundant’ conceptions of properties. Another aspect of the philosophical significance of possible worlds semantics being availed of in Bayesian vision science and empirical linguistics is that it belies the purportedly naturalistic grounds for Quine’s (1963/1976) scepticism of \textit{de re} modality.

\textsuperscript{22}Cf. Burali-Forti (1897/1967). Hodes (1984a) and Hazen (1985) note that abstraction principles based on isomorphism with unrestricted comprehension entrain the paradox.

\textsuperscript{23}For identity conditions on abundant properties – where the domain of properties, in the semantics of second-order logic, is a subset of the domain of objects, and the properties are definable in a metalanguage by predicates whose satisfaction-conditions have been fixed – see Hale (2013a). For a generalization of the abundant conception, such that the domain of properties is isomorphic to the powerset of the domain of objects, see Cook (2014).
3.5 Concluding Remarks

In this essay, the equivalence relations countenanced in Homotopy Type Theory were availed of, in order to specify an abstraction principle for intensional, computational properties. The homotopic abstraction principle for epistemic intensions provides an epistemic conduit into our knowledge of intensions as abstract objects. Because intensional functions in Epistemic Modal Algebra are deployed as core models in the philosophy of mind, Bayesian visual psychology, and natural language semantics, there is independent abductive support for the truth of homotopic abstraction. Epistemic modality may thereby be recognized as both a compelling and a materially adequate candidate for the fundamental structure of mental representational states, and as thus comprising a fragment of the language of thought.
Chapter 4

Two-dimensional Truthmaker Semantics

4.1 Introduction

Philosophical applications of two-dimensional intensional semantics have demonstrated that an account of representation which is sensitive to an array of parameters can play a crucial role in explaining the values of linguistic expressions (Kamp, 1967; Kaplan, 1979); the role of speech acts in affecting shared contexts of information (Stalnaker, 1978; Lewis, 1980,a/1998; MacFarlane, 2005); the relationship between conceivability and metaphysical possibility (Chalmers, 1996); the limits of subjectivism in ethics (Peacocke, 2003); and the viability of modal realism (Russell, 2010).

In order to circumvent issues for the modal analysis of counterfactuals (2012a), and to account for the general notion of aboutness and a subject matter (2015), a hyperintensional, ‘truthmaker’ semantics has recently been developed by Fine (2017a,b). In this essay, I examine the status of two-dimensional indexing in truthmaker semantics, and specify the two-dimensional profile of the grounds for the truth of a formula (Section 2.2). I proceed, then, to outline three novel interpretations of the two-dimensional, hyperintensional framework, beyond the interpretations of multiply indexed intensional semantics that are noted above. The first interpretation provides a formal setting in which to define the distinction between fundamental and derivative truths (Section 3.1). The second interpretation concerns the interaction between the two-dimensional profile of the verifiers for a proposition,
subjective probability, and decision theory (Section 3.2). Finally, a third interpretation of the two-dimensional hyperintensional framework concerns the types of intentional action. I demonstrate, in particular, how multiply indexed truthmaker semantics is able to resolve a puzzle concerning the role of intention in action (Section 3.3). Section 4 provides concluding remarks.

4.2 Two-dimensional Truthmaker Semantics

4.2.1 Intensional Semantics

In his (1979), Evans endeavors to account for the phenomenon of the contingent apriori by distinguishing between two types of modality. In free logic, closed formulas may receive a positive, classical semantic value when the terms therein have empty extensions (op. cit.: 166). Suppose that the name, 'Plotinus', is introduced via the reference fixer, 'the author of the The Enneads'. Then the sentence, 'if anyone uniquely is the author of The Enneads, then Plotinus is the author of the The Enneads' is 'epistemically equivalent' to the sentence, 'if anyone uniquely is the author of The Enneads, then the author of the The Enneads is the author of the The Enneads' (cf. Hawthorne, 2002). Informative identity statements – such as that 'Plotinus = the author of The Enneads' – are thus taken to be epistemically equivalent to vacuously true identity statements – e.g., 'Plotinus = Plotinus' (op. cit.: 177). The apriority of the vacuously true identity statement is thus argued to be a property of the informative identity statement, as well. However, the informative identity statement is contingent. For example, it is metaphysically possible that the author of The Enneads is Plato, rather than Plotinus.

Evans argues that the foregoing 'superficial' type of contingency at issue is innocuous, by distinguishing it from what he refers to as a 'deep' type of contingency according to which a sentence is possibly true only if it is made true by a state of affairs (185). The distinction between the types of modality consists in that superficial contingency records the possible values of a formula when it embeds within the scope of a modal operator – e.g., possibly, x is red and possibly x is blue – whereas deep contingency records whether the formula is made true by a metaphysical state of affairs. In light of the

---

1A premise in the argument is that definite descriptions are non-referring, although – in free logic – still enable the sentences in which they figure to bear a positive, classical value. See Evans (op. cit.: 167-169).
approach to apriority which proceeds via the free-logical, epistemic equivalence of vacuous and informative identity statements, a formula may thus be apriori and yet superficially contingent.\textsuperscript{2} Evans (op. cit.: 183-184; 2004: 11-12) goes further and – independently developing work in two-dimensional intensional semantics by Kamp (1967), Vlach (1973), and Segerberg (1973) – treats the actuality operator as a rigidifier, such that the value of actually $\phi$ determines the counterfactual value of possibly $\phi$.

Two-dimensional intensional semantics provides a framework for regimenting the thought that the value of a formula relative to one parameter determines the value of the formula relative to another parameter. The semantics assigns truth-conditions to formulas, and semantic values to the formula’s component terms. The conditions of the formulas and the values of their component terms are assigned relative to the array of intensional parameters. So, e.g., a term may be defined relative to a context; and the value of the term relative to the context will determine the value of the term relative to an index. For example, according to Kaplan (1979), an utterance’s character is a mapping from the utterance’s context of evaluation to the utterance’s content. According to Stalnaker (op. cit.; 2004), having distinct functions associated with the value of an utterance provides one means of reconciling the necessity of a formula presupposed by speakers with the contingency of the values of assertions made about that formula. According to Chalmers (op. cit.), the value of a formula relative to a context, which ranges over epistemically possible worlds, determines the value of a formula relative to an index, which ranges over metaphysically possible worlds. According to Lewis (op. cit.), the context may be treated as a concrete situation ranging over individuals, times, locations, and worlds; and the index may be treated as ranging over information states relative to the context. Accord-

\textsuperscript{2} Evans’ approach is defined within a single space of metaphysically possible worlds. However, one may define the value of a formula relative to two spaces: A space of epistemic possibilities and a space of metaphysical possibilities. By contrast to securing apriority by (i) eliding the values of informative and vacuous identity statements in a free logic within a single space of metaphysical possibilities, and then (ii) arguing that apriori identity statements are superficially contingent because possibly false, an alternative approach argues that an identity statement is contingent apriori if and only if it is (i) apriori, because the statement is necessarily true in epistemic modal space, while the statement is (ii) contingent, because possibly the statement is false in metaphysical modal space. However, a point of convergence between the above approaches is that neither avails of the notion of deep contingency – i.e., truthmaking – in order to reconcile the apriority of a formula with the possibility of its falsehood.

59
ing to MacFarlane (op. cit.), formulas may receive their value relative to a context ranging over two distinct agents; the context determines the value of an index ranging over their states of information; and the value of the formula may yet be defined relative to a third parameter ranging over the states of an independent, third assessor. Finally, in decision theory, the value of a formula relative to a context, which ranges over a time, location, and agent, constrains the value of the formula relative to a first index on which a space of the agent’s possible acts is built, and the latter will subsequently constrain the value of the formula relative to a second index on which a space of possible outcomes may be built.

Primary, secondary, and 2D intensions can be defined as follows:

- **Primary Intension:**
  \[ \text{pri}(x) = \lambda c.[x]^{c,c} \], with \( c \) an epistemically possible world;

- **Secondary Intension:**
  \[ \text{sec}_{v\alpha}(x) = \lambda w.[x]^{v\alpha,w} \], with \( w \) a metaphysically possible world;

- **2D-Intension:**
  \[ 2D(x) = \lambda c\lambda w[x]^{c,w} = 1. \]

4.2.2 Truthmaker Semantics

A hyperintensional, ‘truthmaker’ semantics has recently been developed by Fine (2017a,b).\(^3\) Truthmaker semantics has been applied, in order to explain the verification-conditions which ground the truth of parts of propositions, rather than of the propositions in their entirety.\(^4\)

Truthmaker semantics is defined over a state space, \( F = \langle S, \sqsubseteq \rangle \), where \( S \) is a set of states comprising a world, and \( \sqsubseteq \) is a parthood relation on \( S \) comprising a partial order, such that it is reflexive \( (a \sqsubseteq a) \), anti-symmetric \( ((a \sqsubseteq b) \land \neg(b \sqsubseteq a)) \), and transitive \( (a \sqsubseteq b, b \sqsubseteq c; a \sqsubseteq c) \) (Fine, 2017a: 19).

\(^3\)The logic for the semantics is classical. Fine (2014) develops a truthmaker semantics for intuitionistic logic.

A proposition $P \subseteq S$ is verifiable if $P$ is non-empty, and is otherwise unverifiable (20).

$s,t \subseteq S$ are compatible if their fusion $s \sqcup t$ is a possible state of $S$ (op. cit.).

$s$ overlaps $t$ if $\exists u(u \sqsubseteq S)$, s.t. $u \sqsubseteq s$ and $u \sqsubseteq t$, denoted $s \cap t$ (op. cit.).

A model, $M$, over $F$ is a tuple, $M = (F,D,V)$, where $D$ is a domain of closed formulas (i.e. propositions), and $V$ is an assignment function mapping propositions $P \in D$ to pairs of subsets of $S$, $\{1,0\}$, i.e. the verifier and falsifier of $P$, such that $[P]^+ = 1$ and $[P]^− = 0$ (35).

The verification-rules in truthmaker semantics are then the following:

- $s \vdash P$ if $s \in [P]^+$ (s verifies $P$, if $s$ is a truthmaker for $P$ i.e. if $s$ is in $P$’s extension);
- $s \vdash\neg P$ if $s \in [P]^−$ (s falsifies $P$, if $s$ is a falsifier for $P$ i.e. if $s$ is in $P$’s anti-extension);
- $s \vdash P$ if $s \vdash P$ (s verifies $P$ if $s$ falsifies $P$);
- $s \vdash \neg P$ if $s \vdash \neg P$ (s falsifies $P$ if $s$ verifies $P$);
- $s \vdash P \land Q$ if $\exists u,u \vdash P, u \vdash Q$, and $s = t \cap u$ (s verifies $P$ and $Q$, if $s$ is the fusion of states, $t$ and $u$, $t$ verifies $P$, and $u$ verifies $Q$);
- $s \vdash P \land Q$ if $s \vdash P$ or $s \vdash Q$ (s falsifies $P$ and $Q$, if $s$ falsifies $P$ or $s$ falsifies $Q$);
- $s \vdash P \lor Q$ if $s \vdash P$ or $s \vdash Q$ (s verifies $P$ or $Q$, if $s$ verifies $P$ or $s$ verifies $Q$);
- $s \vdash P \lor Q$ if $\exists u,u \vdash P, u \vdash Q$, and $s = t \cap u$ (s falsifies $P$ or $Q$, if $s$ is the state overlapping the states $t$ and $u$, $t$ falsifies $P$, and $u$ falsifies $Q$);
- $s$ exactly verifies $P$ if and only if $s \vdash P$ if $s \in [P]$;
- $s$ inexactly verifies $P$ if and only if $s \not\vdash P$ if $\exists s' \sqsubseteq S, s' \vdash P$; and
- $s$ loosely verifies $P$ if and only if, $\forall t, s.t. s \sqcup t, s \sqcup t \vdash P$ (35-36).

Differentiated contents may be defined as follows. A state $s \subseteq S$ is differentiated only if $s$ is the fusion of distinct parts, s.t. $s = s_1 \sqcup s_2$. $s$ is thereby comprised of three parts: An initial state, $s_1$; an additional state, $s_2$; and a total state, $s$. The three states correspond accordingly to three

---

5Fine (op. cit.: 8, 12) avails of product spaces in his discussion of content and subject matter, though we continue here to work with a single space for ease of exposition.
The substraction of Q from P, P – Q, is defined such that p – q = p – (q \cap p) (17).

Finally, subject matters may be defined as follows.

A positive subject matter, p+, expresses a verifiable proposition, \[ [P]^+ \] (20-21).

A negative subject matter, p−, expresses a falsifiable proposition, \[ [P]^− \] (21).

A comprehensive subject matter expresses the fusion of the subject matters both verified and falsified by the fusion of a number of states:

\[ p_{1,+,−} = p_{1,+} \cap p_{1,−} = (s \vdash P \land s \vdash \neg P); \]
\[ p_{2,+,−} = p_{2,+} \cap p_{2,−} = (s \vdash P \land s \vdash \neg P); \]
\[ p_{1,2,+,−} = p_{1,2,+} \cap p_{1,2,−} = (s \vdash P \land s \vdash \neg P) \] (op. cit.).

A differentiated subject matter expresses the fusion of the subject matters that are either verified or falsified by the fusion of a number of states:

\[ p_{1,+/−} = p_{1,+} \cup p_{1,−} = (s \vdash P \vee s \vdash \neg P); \]
\[ p_{2,+/−} = p_{2,+} \cup p_{2,−} = (s \vdash P \vee s \vdash \neg P); \]
\[ p_{1,2,+/−} = p_{1,2,+} \cup p_{1,2,−} = (s \vdash P \vee s \vdash \neg P) \] (op. cit.).

Informally, propositions P and Q express the same subject matters, p and q, when the following conditions hold:

P is exactly about Q if p = q;

P is partly about Q if p and q overlap, such that \( \exists u \subseteq S(u \vdash R); \forall s_1, s_2 \subseteq S, s_1 \vdash P, s_2 \vdash Q; \) and \( u = s_1 \cap s_2, \) such that \( R = P \cap Q; \)

P is entirely about Q if p \( \subseteq \) q; and

P is about Q in its entirety if p \( \supseteq \) q (5).

### 4.2.3 Two-dimensional Truthmaker Semantics

In order to account for two-dimensional indexing, we augment the model, M, with a second state space, S*, on which we define both a new parthood relation, \( \sqsubseteq^* \), and partial function, V*, which serves to map propositions in D to pairs of subsets of S*, \( \{1,0\} \), i.e. the verifier and falsifier of P, such that \[ [P]^+ = 1 \] and \[ [P]^− = 0. \] Thus, M = \( \langle S, S^*, D, \sqsubseteq, \sqsubseteq^*, V, V^* \rangle \). The two-dimensional hyperintensional profile of propositions may then be recorded by defining the value of P relative to two parameters, c,i: c ranges over subsets of S, and i ranges over subsets of S*.
(*) M,s ∈ S, s* ∈ S* ⊨ P iff:
(i) ∃c, [P]_{c,c} = 1 if s ∈ [P]^+; and
(ii) ∃i, [P]_{c,i} = 1 if s* ∈ [P]^+

(Distinct states, s, s*, from distinct state spaces, S, S*, provide a two-dimensional verification for a proposition, P, if the value of P is provided a truthmaker by s. The value of P as verified by s determines the value of P as verified by s*).

We say that P is hyper-rigid iff:

(*) M,s ∈ S, s* ∈ S* ⊨ P iff:
(i) ∀c', [P]_{c,c'} = 1 if s ∈ [P]^+; and
(ii) ∀i, [P]_{c,i} = 1 if s* ∈ [P]^+

The foregoing provides a two-dimensional hyperintensional semantic framework within which to interpret the values of a proposition. In order to account for partial contents, we define the values of subpropositional entities relative again to tuples of states from the distinct state spaces in our model:
s is a two-dimensional exact truthmaker of P if and only if (*);
s is a two-dimensional inexact truthmaker of P if and only if ∃s' ⊑ S, s → s', s' ⊨ P and such that
∃c, [P]_{c,c} = 1 if s' ∈ [P]^+; and
∃i, [P]_{c,i} = 1 if s* ∈ [P]^+;
s is a two-dimensional loose truthmaker of P if and only if, ∃t, s.t. s \sqcup t, s \sqcup t ⊨ P:
∃c, [P]_{c,c} = 1 if s' ∈ [P]^+; and
∃i, [P]_{c,i} = 1 if s* ∈ [P]^+.

• [P]^{c,i} is exactly about [Q]^{c,i} if \text{f}_{1-1}[p^{c,i} \iff q^{c,i}]
  (Suppose that the values of P and of Q are two-dimensionally determined, as above. Then P is exactly about Q if there is a bijection between the two-dimensionally individuated subject matters that they express);

• [P]^{c,i} is partly about [Q]^{c,i} if p and q overlap, s.t. \exists u ⊑ S, s.t. u ⊨ R, and \forall s_1, s_2 ⊑ S, s_1 ⊨ P, s_2 ⊨ Q, and u = s_1 \cap s_2 such that R^{c,c} = P \cap Q. A neighborhood function, A, maps u to a state s* in i where s* ⊨ R^{c,i}.
• $[P]^c_i$ is entirely about $[Q]^c_i$ if $p^c_i \leq q^c_i$

(Suppose that the values of $P$ and of $Q$ are two-dimensionally determined. Then $P$ is entirely about $Q$ if there is a surjection from the subject matter of $Q$ onto the subject matter of $P$);

• $[P]^c_i$ is about $[Q]^c_i$ in its entirety if $p^c_i \Rightarrow q^c_i$

(Suppose that the values of $P$ and of $Q$ are two-dimensionally determined. Then $P$ is about $Q$ in its entirety if there is an injection from the subject matter of $P$ onto the subject matter of $Q$).

4.3 New Interpretations

The two-dimensional account of truthmaker semantics provides a general framework in which a number of interpretations of the state spaces at issue can be defined. The framework may accommodate, e.g., the so-called 'semantic' and 'metasemantic' interpretations of the framework. The semantic interpretation targets, as noted, the bearing of contextual parameters on the values of terms, and provides an account of validity relative to the spaces defined in the term’s intensional parameters (cf. Kaplan, op. cit.). The metasemantic interpretation accommodates, by contrast, the update effects of contingently true assertions, with the necessary propositions of which speakers’ shared information states might be comprised (cf. Stalnaker, op. cit.). The framework may further be provided an ‘epistemic’ interpretation, in order to countenance hyperintensional distinctions in the relations between conceivability, i.e. the space of an agent’s epistemic states, and metaphysical possibility, i.e. the state space of facts (cf. Chalmers, op. cit.). Chapters 4 and 8 outline an epistemic two-dimensional truthmaker semantics in detail. In this section, I advance three novel interpretations of two-dimensional semantics, as witnessed by the new relations induced by the interaction between two-dimensional indexing and hyperintensional value assignments. The three interpretations concern (i) the distinction between fundamental and derivative truths; (ii) probabilistic grounding in the setting of decision theory; and (iii) the structural contents of the types of intentional action.
4.3.1 Fundamental and Derivative Truths

The first novel interpretation concerns the distinction between fundamental and derivative truths. In the foregoing model, the value of the subject matter expressed by a proposition may be verified by states in a first space, which determine, then, whether the proposition is verified by states in a second space. Allowing the first space to be interpreted so as to range over fundamental facts and the second space to be interpreted so as to range over derivative facts permits a precise characterization of the determination relations between the fundamental and derivative grounds for a truth.

Suppose, e.g., that the fundamental facts concern the computational characterization of a subject’s mental states, and let the fundamental facts comprise the first state space. Let the derivative facts concern states which verify whether the subject is consciously aware of their mental representations, and let the derivative facts comprise the second state space. Finally, let $\phi$ be a formula in an experimental task which expresses that there is a particular valence for the contrast-level of a stimulus.\(^6\) The formula’s having a truthmaker in the first space – where the states of which range, as noted, over the subject’s psychofunctional facts – will determine whether the formula has a truthmaker in the second space – where the states of which range over the mental representations of which the subject is consciously aware. If the deployment of some attentional functions provides a necessary condition on the instantiation of phenomenal awareness, then the role of the state of the attentional function in the first space in verifying $\phi$ will determine whether $\phi$ is subsequently verified relative to the second space. Intuitively: Attending to a stimulus with a particular value will constrain whether a truthmaker can be provided for being consciously aware of the stimulus. If the computational facts at issue are fundamental, and the phenomenal facts at issue are derivative, then a precise characterization may be provided of the two-dimensional relations between the verifiers which target fundamental and derivative truths.

Note that hyperintensional truthmaker semantics is consistent with a necessitist modal ontology, according to which necessarily everything is necessarily something. An apparent tension might be thought to arise by augmenting a first-order formal language with both a truthmaker principle – to the effect that a proposition is true only if there is something which entails

\(^6\)In experimental psychology, stimuli which represent contrast gradients are referred to as Gabor patches.
it $[A \rightarrow \exists x \Box [\exists y(x = y \rightarrow A)]]$ – and the principle of the necessary necessity of being – which codifies the thought that necessarily everything is necessarily something $[\Box \forall x \Box y(x = y)]$. The apparent tension is that conjoining the truthmaker principle to the necessity of being entails the necessity of truth (cf. Williamson, 2013a: 400). The contingency of truth is belied by there necessarily being a (possible) object which serves as either a truthmaker or a falsifier for every proposition. However, the foregoing language is consistent with there being contingent truths; e.g., formulas comprised of non-logical vocabulary, such as that there are tigers. Williamson replies by arguing against the truthmaker principle (op. cit.). He observes, however, that a hyperintensional version of the truthmaker principle is consistent with the necessary necessity of being – $A \rightarrow \exists P(A$ because $P)$ – which is an instance of the principle of the priority of being to truth: namely, that for all formulas $\pi A$, $[T\pi (A$ because $A)]$ (400-401). Consistently, then, with the necessary necessity of being, hyperintensional truthmaker semantics is one means of interpreting the explanatory relation – $A$ 'because' $B$ – in the principle of the priority of being to truth.

4.3.2 Decision Theory

A second novel interpretation of two-dimensional truthmaker semantics concerns the types of intentional action, and the interaction of the latter with decision theory. As noted in the foregoing, two-dimensional intensional semantics may be availed of in order to explain how the value of a formula relative to a context ranging over an agent and time will determine the value of the formula relative to an index ranging over a space of admissible actions made on the basis of the formula, where the value of the formula relative to the context and first index will determine the value of the formula relative to a second index, ranging over a space of outcomes.

One notable feature of the decision-theoretic interpretation is that it provides a natural setting in which to provide a gradational account of truth-making. A proposition and its component expressions are true, just if they are verified by states in a state space, such that the state and its parts fall within the proposition’s extension. In decision theory, a subject’s expectation that the proposition will occur is recorded by a partial belief function, mapping the proposition to real numbers in the $\{0,1\}$ interval. The subject’s desire that the proposition occurs is recorded by a utility function, the quantitative values of which – e.g., 1 or 0 – express the qualitative value of the
proposition’s occurrence. The evidential expected utility of a proposition’s occurrence is calculated as the probability of its obtaining conditional on an agent’s action, as multiplied by the utility to the agent of the proposition’s occurrence. The causal expected utility of the proposition’s occurrence is calculated as the probability of its obtaining, conditional on both the agent’s acts and background knowledge of the causal efficacy of their actions, multiplied by the utility of the proposition’s occurrence.

There are three points at which a probabilistic construal of the foregoing may be defined. One point concerns the objective probability that the proposition will be verified, i.e. the chance thereof. The second point concerns subjective probability with which a subject partially believes that the proposition will obtain. A third point concerns the probability that an outcome will occur, where the space of admissible outcomes will be constrained by a subject’s acts. An agent’s actions will, in the third case, constrain the admissible verifiers in the space of outcomes, and thus the probability that the verifier for the proposition will obtain as an outcome.\(^7\)

In order formally to countenance the foregoing, we define a probability measure on a state space, such that the probability measure satisfies the Kolmogorov axioms: normality \([\text{Pr}(T) = 1]\); non-negativity \([\text{Pr}(\phi) \geq 0]\); additivity \([\text{For disjoint } \phi \text{ and } \psi [\text{Pr}(\phi \cup \psi) = \text{Pr}(\phi) + \text{Pr}(\psi)]\)]; and conditionalization \([\text{Pr}(\phi | \psi) = \text{Pr}(\phi \cap \psi) / \text{Pr}(\psi)]\). In order to account for the interaction between objective probability and the verification-conditions in truthmaker semantics, we avail, then, of a regularity condition in our earlier model, \(M\), in which the assignment function, \(V\), maps propositions \(P \in D\) to pairs of subsets of \(S\), \(\{1,0\}\), i.e. the verifier and falsifier of \(P\), such that \(\llbracket P \rrbracket^+ = \{0,1\}\) and \(\llbracket P \rrbracket^- = 1 - P\). In our gradational truthmaker semantics, a state, \(s\), verifies a

\(^7\)A proponent of metaphysical indeterminacy might further suggest that the verifiers are themselves gradational; thus, rather than target the probability of a verifier’s realization, the proponent of metaphysical indeterminacy will suggest that a proposition \(P\) is made true only to a certain degree, such that both of the proposition’s extension and anti-extension will have non-negative, real values. One objection to the foregoing account of metaphysical indeterminacy for truthmakers is, however, that the metalogic for many-valued logic is classical (cf. Williamson, 2014a). A distinct approach to metaphysical indeterminacy is proffered by Barnes and Williams (2011), who argue that metaphysical indeterminacy consists in persistently unpointed models, i.e. a case in which it is unclear which among a set of worlds is actual, even upon filtering the set with precisifications. A proponent of metaphysical indeterminacy for probabilistic truthmaker semantics might then argue both that the realization of a verifier has a gradational value and that it is indeterminate which of the states which can verify a given formula is actual.
proposition, P, if the probability that s is in P’s extension is greater than or equal to .5:

\[ s \models P \text{ if } \Pr(s \in [P]^+) \geq .5. \]

A state, s, falsifies a proposition P if the probability that s is in P’s extension is less than .5 iff the probability that s is in P’s anti-extension is greater than or equal to .5

\[ s \not\models P \text{ if } \Pr(s \in [P]^-) \geq .5 \]
\[ \text{iff } \Pr(s \in [P]^+) < .5. \]

The subjective probability with regard to the proposition’s occurrence is expressed by a probability measure satisfying the Kolmogorov axioms as defined on a second state space, i.e., a space whose points are interpreted as concerning the subject’s states of information. The formal clauses for partial belief in truthmaker semantics are the same as in the foregoing, save that the probability measures express the mental states of an agent, by being defined on the space of their states of information.

Finally, the interaction between objective and subjective probability measures in hyperintensional semantics may be captured in two ways.

One concerns the interaction between the chance of a proposition’s occurrence, the subject’s partial belief that the proposition will occur, and the spaces for the subjects actions and outcomes. The formal clause for the foregoing will then be as follows:

\[ M,s \models [P]^{c'(a,o)} > .5, \]

where c ranges over the space of physical states, and a probability measure recording objective chance is defined thereon; c’ ranges over the space of an agent’s states of information, and the value of P relative to c’ determines the value of P relative to the space of the agent’s acts, a, where the latter determines the space of admissible outcomes concerning P’s occurrence, o. Thus, the parameters, c’,a,o possess a hyperintensional two-dimensional profile, whereas the space of physical states, c, constrains the verification of the proposition’s occurrence, without determining the values of the subject’s partial beliefs and their subsequently conceivable actions and outcomes.

Directly accounting for the relation between c and c’ – i.e., specifying a norm on the relation between chances and credences – is the second means by which to account for how objective gradational truthmakers interact with a subject’s partial beliefs about whether propositions are verified. Following Lewis (1980,b/1987), a candidate chance-credence norm may be what he
refers to as the 'principal principle'. The principal principle states that an agent’s partial belief that a proposition will be verified, conditional on the objective chance of the proposition’s occurrence and the admissible evidence, will be equal to the objective chance of the proposition’s occurrence itself:

\[ \Pr_s(P \mid \text{ch}(P) \land E) = \text{ch}(P). \]

### 4.3.3 Intentional Action

A third novel interpretation of two-dimensional intensional semantics provides a natural setting in which to delineate the structural content of the types of intentional action. For example, the mental state of intending to pursue a course of action may be categorized as falling into three types, where intending—that is treated as a modal operator defined on an agent’s space of states of information. One type targets a unique structural content for the state of acting intentionally, such that an agent intends to bring it about that \( \phi \) just if the intention satisfies a clause which mirrors that outlined in the last paragraph:

- \( \llbracket \text{Intenton-in-Action}(\phi) \rrbracket_w = 1 \) only if \( \exists w' \llbracket \phi \rrbracket_{w', c(=t,t), a, o} = 1. \)

A second type of intentional action may be recorded by a future-directed operator, such that an agent intends to \( \phi \) only if they intend to pursue a course of action in the future, only if there is a world and a future time relative to which the agent’s intention is satisfied:

- \( \llbracket \text{Intention-for-the-future}(\phi) \rrbracket_w = 1 \) only if \( \exists w' \forall t \exists t'[t < t' \land \llbracket \phi \rrbracket_{w', t'} = 1] \).

---

8See Pettigrew (2012), for a justification of a generalized version of the principal principle based on Joyce’s (1998) argument for probabilism. Probabilism provides an accuracy-based account of partial beliefs, defining norms on the accuracy of partial beliefs with reference only to worlds, metric ordering relations, and probability measures thereon. The proposal contrasts to pragmatic approaches, according to which a subject’s probability and utility measures are derivable from a representation theorem, only if the agent’s preferences with regard to a proposition’s occurrence are consistent (cf. Ramsey, 1926). Probabilism states, in particular, that, if there is an ideal subjective probability measure, the ideality of which consists e.g. in its matching objective chance, then one’s probability measure ought to satisfy the Kolmogorov axioms, on pain of there always being a distinct probability measure which will be metrically closer to the ideal state than one’s own.
Finally, a third type of intentional action concerns reference to the intention as an explanation for one’s course of action. In Chapter 14, I regiment the structural content of this type of intention as a modal operator which receives its value only if a hyperintensional grounding operator which takes scope over a proposition and an action, receives a positive semantic value. The formal properties of the grounding operator follows the presentation in Fine (2012b), according to which the operator corresponds to relations of parthood, which satisfy the logical properties of reflexivity and permit bijections between the states at issue, namely, actions and intentional descriptions.

The varieties of subject matter, as defined in two-dimensional truthmaker semantics, can be availed of in order to enrich the present approach. Having multiple state spaces from which to define the verifiers of a proposition enables a novel solution to issues concerning the interaction between action and explanation. The third type of intentional action may be regimented, as noted, by the agent’s reference to an intention as an explanation for her course of action.

The foregoing may also be availed of, in order to provide a novel solution to an issue concerning the interaction between involuntary and intentional action. The issue is as follows. Wittgenstein (1953/2009; 621) raises the inquiry: ‘When I raise my arm, my arm goes up. Now the problem arises: what is left over if I subtract the fact that my arm goes up from the fact that I raise my arm?’ Because the arm’s being raised has at least two component states, namely, the arm’s going up and whatever the value of the variable state might be, the answer to Wittgenstein’s inquiry is presumably that the agent’s intentional action is the value of the variable state, such that a combination of one’s intentional action and one’s arm going up is sufficient for one’s raising one’s arm. The aforementioned issue with the foregoing concerns how precisely to capture the notion of partial content, which bears on the relevance of the semantics of the component states and the explanation of the unique state entrained by their combination.

Given our two-dimensional truthmaker semantics, a reply to Wittgenstein’s inquiry which satisfies the above desiderata may be provided. Let W express a differentiated subject matter, whose total content is that an agent’s arm is raised. W expresses the total content that an agent’s arm is raised, because W is comprised of an initial content, U (that one’s arm goes up), and an additional content, R (that one intends to raise one’s arm).

The verifier for W may be interpreted as a two-dimensional loose truth-
maker. Let \( c \) range over an agent’s motor states, \( S \). Let \( i \) range over an agent’s states of information, \( S^* \). We define a modal operator for intentional action on the state space of the agent’s motor actions. The value of the modal operator for intentional action is positive just if a selection function, \( g \), is a mapping from the powerset of motor actions in \( S \) to a unique state \( s' \) in \( S \). This specifies the initial, partial content, \( U \), that one’s arm goes up. An intention may then be defined as a unique state, \( s^* \), in the agent’s state of information, \( S^* \). The state, \( s^* \), specifies the additional, partial content \( R \), that one intends to raise one’s arm.

Formally:

\[
s \vdash U \text{ only if } \exists s' \sqsubseteq S, \text{ such that } g: s \rightarrow s', \text{ s.t. } s' \vdash U,
\]

\[
\exists s^*, s^* \vdash R, \text{ and }
\]

\[
W = U \sqcup R.
\]

The two-dimensional loose truthmaker for one’s arm being raised may then be defined as follows:

\[
\exists c_{s \rightarrow s'} [W]^{c\cdot} = 1 \text{ if } s' \in [W]^+, \text{ and }
\]

\[
\exists i_{s*} [W]^{c\cdot} = 1 \text{ if } s^* \in [W]^+.
\]

Intuitively, the value of the total content that one’s arm is raised is defined relative to a set of motor states – where a first intentional action selects a series of motor states which partly verify that one’s arm goes up. The value of one’s arm being raised, relative to (the intentionally modulated) motor state of one’s arm possibly going up, determines the value of one’s arm being raised relative to the agent’s distinct intention to raise their arm. The agent’s first intention selects among the admissible motor states, and – all else being equal – the motor states will verify the fact that one’s arm goes up. Recall that the value of a formula relative to a context determines the value of the formula relative to an index. As follows, the priority of the motor act to the subsequent intention to raise one’s arm is thus that it must first be possible for one’s arm to go up in order to determine whether the subsequent intention to raise one’s arm can be satisfied.\(^9\)

The fusion of (i) the state corresponding to the initial partial content that one’s arm goes up, and (ii) the state corresponding to the additional partial content that one intends to raise one’s arm, is sufficient for the verification of (iii) the state corresponding to the total content that one’s arm is raised.

\(^9\)The role of the first intention in acting as a selection function on the space of motor actions corresponds to the comparator functions stipulated in the contemporary cognitive science of action theory. For further discussion of the comparator model, see Frith et al. (2000) and Pacherie (2012).
4.4 Concluding Remarks

In this essay, I have endeavored to establish foundations for the interaction between two-dimensional indexing and hyperintensional semantics. I examined, then, the philosophical significance of the framework by developing three, novel interpretations of two-dimensional truthmaker semantics, in light of the new relations induced by the model.

The first interpretation enables a rigorous characterization of the distinction between fundamental and derivative truths. The second interpretation evinces how the elements of decision theory are definable within the two-dimensional hyperintensional setting, and a novel account was then outlined concerning the interaction between probability measures and hyperintensional grounds. The third interpretation of two-dimensional hyperintensional semantics concerns the structural content of the types of intentional action. Finally, I demonstrated how the hyperintensional array of state spaces, relative to which propositions may be verified, may serve to resolve a previously intransigent issue concerning the role of intention in action.
Chapter 5

Non-Transitive Self-Knowledge: Luminosity via Modal $\mu$-Automata

This essay provides a novel account of self-knowledge, which avoids the epistemic indeterminacy witnessed by the invalidation of modal axiom 4 in epistemic logic; i.e. the KK principle: $\Box\phi \rightarrow \Box\Box\phi$. The essay argues, by contrast, that – despite the invalidation of modal axiom 4 on its epistemic interpretation – states of epistemic determinacy might yet be secured by countenancing self-knowledge on the model of fixed points in monadic second-order modal logic, i.e. the modal $\mu$-calculus.

Counterinstances to modal axiom 4 – which records the property of transitivity in labeled transition systems\(^1\) – have been argued to occur within various interpretations of the sorites paradox. Suppose, e.g., that a subject is presented with a bounded continuum, the incipient point of which bears a red color hue and the terminal point of which bears an orange color hue. Suppose, then, that the cut-off points between the points ranging from red to orange are indiscriminable, such that the initial point, a, is determinately red, and matches the next apparent point, b; b matches the next apparent point, c; and thus – by transitivity – a matches c. Similarly, if b matches c, and c matches d, then b matches d. The sorites paradox consists in that iterations of transitivity would entail that the initial and terminal points in the bounded continuum are phenomenally indistinguishable. However, if one

\(^1\)Cf. Kripke (1963).
takes transitivity to be the culprit in the sorites, then eschewing the principle would entail a rejection of the corresponding modal axiom (4), which records the iterative nature of the relation.² Given the epistemic interpretation of the axiom – namely, that knowledge that a point has a color hue entails knowing that one knows that the point has that color hue – a resolution of the paradox which proceeds by invalidating axiom 4 subsequently entrains the result that one can know that one of the points has a color hue, and yet not know that they know that the point has that color hue (Williamson, 1990: 107-108; 1994: 223-244; 2001: chs. 4-5). The non-transitivity of phenomenal indistinguishability corresponds to the non-transitivity of epistemic accessibility. As Williamson (1994: 242) writes: 'The example began with the non-transitive indiscriminability of days in the height of the tree, and moved on to a similar phenomenon for worlds. It seems that this can always be done. Whatever x, y and z are, if x is indiscriminable from y, and y from z, but x is discriminable from z, then one can construct miniature worlds \( w_x, w_y \) and \( w_z \) in which the subject is presented with x, y and z respectively, everything else being relevantly similar. The indiscriminability of the objects is equivalent to the indiscriminability of the corresponding worlds, and therefore to their accessibility. The latter is therefore a non-transitive relation too.' The foregoing result holds, furthermore, in the probabilistic setting, such that the evidential probability that a proposition has a particular value may be certain – i.e., be equal to 1 – while the iteration of the evidential probability operator – recording the evidence with regard to that evidence – is yet equal to 0. Thus, one may be certain on the basis of one’s evidence that a proposition has a particular value, while the higher-order evidence with regard to one’s evidence adduces entirely against that valuation (Williamson, 2014).

In the foregoing argument, ‘safety’ figures as a necessary condition on knowledge, and is codified by margin-for-error principles of the form: \( \forall x \forall \phi [K^{m+1} \phi(x) \to K^m \phi(x+1)] \)’ (Williamson, 2001: 128; Gómez-Torrente, 2002: 114). Intuitively, the safety condition ensures that if one knows that a predicate is satisfied, then one knows that the predicate is satisfied in relevantly similar worlds. Williamson targets the inconsistency of margin-for-error principles, the luminosity principle \( [\forall x \forall \phi [\phi(x) \to K \phi(x)]] \), and the characterization of the sorites as occurring when an object satisfies a property, such that similar objects would further do so. The triad evinces that the luminosity principle

²For more on non-transitivist approaches to the sorites, see Zardini (2019).
is false, given the plausibility of margin-for-error principles and the characterization of the sorites. In cases, further, in which conditions on knowledge are satisfied, epistemic indeterminacy is supposed to issue from the non-transitivity of the accessibility relation on worlds (1994: 242).

One of the primary virtues of the present proposal is thus that it targets the property of transitivity directly, because transitivity both engenders the sorites paradox on the assumption that the states are known and the property is codified by the epistemic modal axiom for transitivity, i.e., 4 or the KK principle. By so doing, it permits a uniform interpretation of transitivity in the sorites – as codified by the KK principle – such that it applies not only to epistemic accessibility relations whose obtaining is relevant to the safety condition, but further to the logical property and its explanatory role in engendering the paradox.

A second virtue adducing in favor of the foregoing, 'epistemicist' approach to vagueness – which takes the latter to be a phenomenon of epistemic indeterminacy – is that vagueness can be explained without having to revise the underlying logic. The epistemicist approach is consistent with classical logical laws, such as e.g. the law of excluded middle; and thus it can determinately be the case that a point has a color hue; determinately be the case that the next subsequent point has a distinguishable color hue; and one can in principle know where in the continuum the cut-off between the two points lies – yet vagueness will consist in the logical limits – i.e. the non-transitivity – of one’s state of knowledge. Thus, one will not in principle be able to know that they know the point at which the color hues are dissimilar.

In this essay, I endeavor to provide a novel account which permits the retention of both classical logic as well as a modal approach to the phenomenon of vagueness, while salvaging the ability of subjects to satisfy necessary conditions on self-knowledge. I will argue that – despite the invalidity of modal axiom 4, given the non-transitivity of the similarity relation – a distinct means of securing an iterated state of knowledge concerning one’s first-order knowledge that a particular state obtains is by availing of fixed point, non-deterministic automata in the setting of coalgebraic modal logic. Propositional modal logic is equivalent to the bisimulation-invariant fragment of fixed point monadic second-order logic.\(^3\) The fixed point higher-order modal logic is referred to as the modal $\mu$-calculus, where $\mu(x)$ is an operator recording a least fixed point. Despite the non-transitivity of sorites phenomena

\[^3\text{Cf. Janin and Walukiewicz (1996).}\]
– such that, on its epistemic interpretation, the subsequent invalidation of modal axiom 4 entails structural, higher-order epistemic indeterminacy – the modal $\mu$-calculus provides a natural setting in which a least fixed point can be defined with regard to the states instantiated by non-deterministic modal automata. In virtue of recording iterations of particular states, the least fixed points witnessed by non-deterministic modal automata provide, then, an escape route from the conclusion that the invalidation of the KK principle provides an exhaustive and insuperable obstruction to self-knowledge. Rather, the least fixed points countenanced in the modal $\mu$-calculus provide another conduit into subjects’ knowledge to the effect that they know that a state has a determinate value. Thus, because of the fixed points definable in the modal $\mu$-calculus, the non-transitivity of the similarity relation is yet consistent with necessary conditions on epistemic determinacy and self-knowledge, and the states at issue can be luminous to the subjects who instantiate them.

In the remainder of the essay, we introduce labeled transition systems, the modal $\mu$-calculus, and non-deterministic Kripke (i.e., $\mu$-) automata. We recount then the sorites paradox in the setting of the modal $\mu$-calculus, and demonstrate how the existence of fixed points enables there to be iterative phenomena which ensure that – despite the invalidation of modal axiom 4 – iterations of mental states can be secured, and can thereby be luminous.

A labeled transition system is a tuple comprised of a set of worlds, $S$; a valuation, $V$, from $S$ to its powerset, $\wp(S)$; and a family of accessibility relations, $R$. So $\text{LTS} = \langle S, V, R \rangle$ (cf. Venema, 2012: 7). A Kripke coalgebra combines $V$ and $R$ into a Kripke functor, $\sigma$; i.e. the set of binary morphisms from $S$ to $\wp(S)$ (op. cit.: 7-8). Thus for an $s \in S$, $\sigma(s) := [\sigma_V(s), \sigma_R(s)]$ (op. cit.). Satisfaction for the system is defined inductively as follows: For a formula $\phi$ defined at a state, $s$, in $S$,

- $[\phi]^S = V(s)$
- $[\neg \phi]^S = S - V(s)$
- $[\bot]^S = \emptyset$
- $[T]^S = S$
- $[\phi \lor \psi]^S = [\phi]^S \cup [\psi]^S$
- $[\phi \land \psi]^S = [\phi]^S \cap [\psi]^S$
- $[\Diamond_s \phi]^S = [R_s][\phi]^S$
- $[\Box_s \phi]^S = [R_s][\phi]^S$, with

$^4$Alternatively, $M, s \models \phi$ if $s \in V(\phi)$ (9).
\( (R_s)(\phi) := \{ s' \in S \mid R_s[s'] \cap \phi \neq \emptyset \} \) and
\[ [R_s](\phi) := \{ s' \in S \mid R_s[s'] \subseteq \phi \} \] (9)
\[ [\mu x. \phi] = \bigcap \{ U \subseteq S \mid [\phi]_{\tau[U]} \subseteq U \} \] (Fontaine, 2010: 18)
\[ [\nu x. \phi] = \bigcup \{ U \subseteq S \mid U \subseteq [\phi]_{\tau[U]} \} \] (op. cit.; Fontaine and Place, 2010),
where \( \tau[x \mapsto U] \) is the assignment \( \tau' \) s.t. \( \tau'(x) = U \) and \( \tau'(y) = \tau(y) \), for all variables \( y \neq x \) (op. cit.).

A Kripke coalgebra can be represented as the pair \( (S, \sigma : S \rightarrow KS) \) (Venema, 2020: 8.1)

In our Kripke colagebra, we have \( M_s \models \langle \pi^* \rangle \phi \iff (\phi \lor \circ_s \langle \pi^* \rangle \phi) \) (Venema, 2012: 25). \( \langle \pi^* \rangle \phi \) is thus said to be the fixed point for the equation, \( x \iff \phi \lor \circ x \), where the value of the formula is a function of the value of \( x \) conditional on the constancy in value of \( \phi \) (38). The smallest solution of the formula, \( x \iff \phi \lor \circ x \), is written \( \mu x. \phi \lor \circ x \) (25). The value of the least fixed point is, finally, defined more specifically thus:
\[ [\mu x. \phi \lor \circ x] = V(\phi) \cup \langle R \rangle ([\mu x. \phi \lor \circ x]) \] (38).

A non-deterministic automaton is a tuple \( A = \langle A, \delta, Acc, a_1 \rangle \), with \( A \) a finite set of states, \( a_1 \) being the initial state of \( A \); \( \delta \) is a transition function s.t. \( \delta : A \rightarrow \phi(A) \); and \( Acc \subseteq A \) is an acceptance condition which specifies admissible conditions on \( \delta \) (60, 66).

Let two Kripke models \( A = \langle A, a \rangle \) and \( S = \langle S, s \rangle \), be bisimilar if and only if there is is a non-empty binary relation, \( Z \subseteq A \times S \), which is satisfied, if:

(i) For all \( a \in A \) and \( s \in S \), if \( aZs \), then \( a \) and \( s \) satisfy the same proposition letters;
(ii) The forth condition. If \( aZs \) and \( R_{\Delta}a, v_1 \ldots v_n \), then there are \( v'_1 \ldots v'_n \) in \( S \), s.t.
- for all \( i \ (1 \leq i \leq n) \ \nu_i Zv'_i \), and
- \( R'_{\Delta}a, v'_1 \ldots v'_n \);
(iii) The back condition. If \( aZs \) and \( R'_{\Delta}a, v'_1 \ldots v'_n \), then there are \( v_1 \ldots v_n \) in \( A \), s.t.
- for all \( i \ (1 \leq i \leq n) \ \nu_i Zv'_i \) and
- \( R_{\Delta}a, v_1 \ldots v_n \) (cf. Blackburn et al, 2001: 64-65).

Bisimulations may be redefined as relation liftings. We let, e.g., a Kripke functor, \( K \), be such that there is a relation \( K! \subseteq K(A) \times K(A') \) (Venema, 2020: 81). Let \( Z \) be a binary relation s.t. \( Z \subseteq A \times A' \) and \( \phi!Z \subseteq \phi(A) \times \phi(A') \), with
℘!Z := \{ (X,X') \mid \forall x \in X \exists x' \in X' \text{ with } (x,x') \in Z \wedge \forall x' \in X' \exists x \in X \text{ with } (x,x') \in Z \} \text{ (op. cit.).}

Then, we can define the relation lifting, K!, as follows:

K! := \{ ([\pi, X], (\pi', X')) \mid \pi = \pi' \text{ and } (X,X') \in \wp!Z \} \text{ (op. cit.).}

Finally, given the Kripke functor, K, K can be defined as the \( \mu \)-automaton, i.e., the tuple \( A = \langle A, \delta, a_I \rangle \), with \( a_I \in A \) defined again as the initial state in the set of states \( A \); and \( \delta \) defined as a mapping such that \( \delta : A \to \wp(\mathbf{K}A) \), where the \( \exists \) subscript indicates that \( (a,s) \in A \times S \to \{ (a',s) \in \mathbf{K}(A) \times S \mid a' \in \delta(a) \} \) (93).

The philosophical significance of the foregoing can now be witnessed by defining the \( \mu \)-automata on an alphabet; in particular, a non-transitive set comprising a bounded real-valued, ordered sequence of chromatic properties. Although the non-transitivity of the ordered sequence of color hues belies modal axiom 4, such that one can know that a particular point in the sequence has a particular value although not know that one knows that the point satisfies that value, the chromatic values, \( \phi \), in the non-transitive set of colors nevertheless permits every sequential input state in the \( \mu \)-automaton to define a fixed point. In order for there to be least and greatest fixed points, there must be monotone operators defined on complete lattices. As Venema (2020: A-2) writes: "A partial order is a structure \( \mathbb{P} = \langle P, \leq \rangle \) such that \( \leq \) is a reflexive, transitive and antisymmetric relation on \( P \). Given a partial order \( \mathbb{P} \), an element \( p \in \mathbb{P} \) is an upper bound (lower bound, respectively) of a set \( X \subseteq P \) if \( p \geq x \) for all \( x \in X \) (\( p \leq x \) for all \( x \in X \)). If the set of upper bounds of \( X \) has a minimum, this element is called the least upper bound, supremum, or join of \( X \), notation: \( \lor X \). Dually, the greatest lower bound, infimum, or meet of \( X \), if existing, is denoted as \( \land X \) . . . A partial order \( \mathbb{P} \) is called a lattice if every two-element subset of \( P \) has both an infimum and a supremum; in this case, the notation is as follows: \( p \land q := \land\{p,q\} \), \( p \lor q := \lor\{p,q\} \) . . . A partial order \( \mathbb{P} \) is called a complete lattice if every subset of \( P \) has both an infimum and a supremum . . . A complete lattice will usually be denoted as a structure \( C = \langle C, \lor, \land \rangle \)." 'Let \( \mathbb{P} \) and \( \mathbb{P}' \) be two partial orders and let \( f : P \to P' \) be some map. Then \( f \) is called monotone or order preserving if \( f(x) \leq' f(y) \) whenever \( x \leq y \) . . . '(3.1). 'Let \( \mathbb{P} = \langle P, \leq \rangle \) be a partial order, and let \( f : P \to P \) be some map. Then an element \( p \in \mathbb{P} \) is called a prefixpoint of \( f \) if \( f(p) \leq p \), a postfixpoint of \( f \) if \( p \leq f(p) \), and a fixpoint if \( f(p) = p \). The sets of prefixpoints, postfixpoints, and fixpoints of \( f \) are denoted respectively as \( \text{PRE}(f) \), \( \text{POS}(f) \) and \( \text{FIX}(f) \). / In case the set of fixpoints of \( f \) has a least (respectively greatest) member, this element is denoted LFP. \( f \) (GFP.\( f \), respectively)' (3-2). The Knaster-Tarski Theorem says, then, that,
for a complete lattice, $\mathbb{C} = \langle \mathbb{C}, \lor, \land \rangle$, with $f: \mathbb{C} \to \mathbb{C}$ being monotone, $f$ has both a least and greatest fixpoint, $\text{LFP}.f = \land \text{PRE}(f)$, and $\text{GFP}.f = \lor \text{POS}(f)$ (op. cit.).

The epistemic approach to vagueness relies, as noted, on the epistemic interpretation of the modal operator, such that the invalidation of transitivity and modal axiom 4 ($\Box \phi \rightarrow \Box \Box \phi$) can be interpreted as providing a barrier to a necessary condition on self-knowledge. Crucially, $\mu$-automata can receive a similar epistemic interpretation. An epistemic interpretation of a $\mu$-automaton is just such that the automaton operates over epistemically possible worlds. The automaton can thus be considered a model for an epistemic agent. The transition function accounts for the transition from one epistemic state to another, e.g., as one proceeds along the stages of a continuum. A fixed point operator on a given epistemic state, e.g., $\Box(\phi)$ where $\Box$ is interpreted so as to mean knowledge-that, amounts to one way to iterate the state. If one knows a proposition $\phi$, the least fixed point operation, $\mu(\phi)$, records an iteration of the epistemic state, knowledge of knowledge, and similarly for belief. Thus, interpreting the $\mu$-automaton epistemically permits the fixed points relative to the arbitrary points in the ordered continuum to provide a principled means – distinct from the satisfaction of the KK principle – by which to account for the pertinent iterations of epistemic states unique to an agent’s self-knowledge.
Part II: Conceivability and Metaphysical Possibility
Chapter 6

Conceivability, Haecceities, and Essence

6.1 Introduction

In this essay, I endeavor to provide an account of how the epistemic interpretation of two-dimensional intensional semantics can be sensitive to haecceities and essential properties more generally. Let a model, M, be comprised of a set of epistemically possible worlds C; a set of metaphysically possible worlds W; a domain, D, of terms and formulas; binary relations defined on each of C and W; and a valuation function mapping terms and formulas to subsets of C and W, respectively. So, \( M = \langle C, W, D, R_C, R_W, V \rangle \). A term or formula is epistemically necessary or apriori iff it is inconceivable for it to be false (\( \Box \iff \neg \Diamond \neg \)). A term or formula is negatively conceivable iff nothing rules it out apriori (\( \Diamond \iff \neg \Box \neg \)). A term or formula is positively conceivable only if the term or formula can be perceptually imagined. According to the epistemic interpretation of two-dimensional intensional semantics, the semantic value of a term or formula can then be defined relative to two parameters, a context and an index.\(^1\) The context ranges over the set of epistemically possible worlds, and the index ranges over the set of metaphysically possible worlds. The value of the term or formula relative to the context determines the value of the term or formula relative to the index. Thus, the epistemically possible value of the term or formula constrains the metaphysically possible value of the term or formula; and so conceivability might, given the

foregoing, serve as a guide to metaphysical possibility.

Roca-Royes (2011) and Chalmers (2010; 2011; 2014) note that, on the above semantics, epistemic possibility cannot track the difference between the metaphysical modal profile of a non-essential proposition – e.g., that there is a shooting star – and the metaphysical modal profile of an essential definition, such as a theoretical identity statement – e.g., that water = H2O. Another principle of modal metaphysics to which epistemic possibilities are purported to be insensitive is haecceity comprehension; namely, that $\Box \forall x, y \exists \Phi (\Phi x \iff x = y)$.

The aim of this note is to redress the contention that epistemic possibility cannot be a guide to the principles of modal metaphysics. I will argue that the interaction between the two-dimensional intensional framework and the mereological parthood relation enables the epistemic possibility of haecceity comprehension to entail the metaphysical possibility of haecceity comprehension. Further, the super-rigidity of essential properties entrains that the epistemic possibility of their obtaining entails the metaphysical possibility of their obtaining.

In Section 2, I examine a necessary condition on admissible cases of conceivability entailing metaphysical possibility in the two-dimensional intensional framework, focusing on the property of super-rigidity. I argue that – despite the scarcity of properties which satisfy the super-rigidity condition – metaphysical properties such as the parthood relation and essential properties do so. In Section 3, I address objections to two dogmas of the semantic rationalism underpinning the epistemic interpretation of two-dimensional intensional semantics. The first dogma states that distinctions can be delineated between linguistic intensions and conceptual epistemic intensions, while the second dogma records that there are criteria on the basis of which formal from informal domains, unique to the extensions of various concepts, can be distinguished, such that the modal profiles of those concepts would thus be determinate. I examine the Julius Caesar problem as a test case. I specify, then, a two-dimensional intensional formula encoding the relation between the epistemic possibility of haecceity comprehension and its metaphysical possibility, and I generalize the approach to essential properties. In Section 4, I address objections from the indeterminacy of ontological principles relative to the space of epistemic possibilities, and from the consistency of epistemic modal space. Section 5 provides concluding remarks.
6.2 Super-rigidity

Mereological parthood satisfies a crucial condition in the epistemic interpretation of two-dimensional intensional semantics. The condition is called super-rigidity, and its significance is that, unless the semantic value for a term is super-rigid, i.e. maps to the same extension throughout the classes of epistemic and metaphysical possibilities, the extension of the term in epistemic modal space risks diverging from the extension of the term in metaphysical modal space.

There appear to be only a few expressions which satisfy the super-rigidity condition. Such terms include those referring to the properties of phenomenal consciousness, to the parthood relation, and perhaps to the property of friendship (Chalmers, 2012: 367, 374). Other candidates for super-rigidity are taken to include metaphysical terms such as 'cause' and 'fundamental'; numerical terms such as 'one'; and logical constants such as '$\land$' (Chalmers, op. cit.). However, there are counterexamples to each of the foregoing proposed candidates.

Against the super-rigidity of 'fundamental', Fine (2001: 3) argues that a proposition is fundamental if and only if it is real, while Sider (2011: 112, 118) argues that a proposition is fundamental iff it possesses a truth-condition (in a 'metaphysical semantics', stated in perfectly joint-carving terms) for the sub-propositional entities – expressed by quantifiers, functions, predicates – comprising the target proposition. The absolute joint-carving terms are taken to include logical vocabulary (including quantifiers), metaphysical predicates such as mereological parthood, and physical predicates.

Against the super-rigidity of 'cause', Sider (op.cit.: 8.3.5) notes that a causal deflationist might argue that causation is non-fundamental. By contrast, a causal nihilist might argue that causation is non-fundamental as well, though for the distinct reason that there is no causation. So, while both the deflationist and nihilist believe that 'cause' does not carve at the joints – the nihilist can still state that there is a related predicate, 'cause*', such that they can make the joint-carving claim that 'Nothing causes* anything', whereas the deflationist will remain silent, and maintain that no broadly causal locutions carve at the joints.

Against the super-rigidity of 'two', Benacerraf (1965) notes that, in the reduction of number theory to set theory, there must be, and is not, a principled reason for which to prefer the identification of natural numbers with von Neumann ordinals (e.g., $2 = \{\emptyset, \emptyset\}$), rather than with Zermelo ordinals.
(i.e., an order-type of a well-ordering $2 = \{\emptyset\}$).\(^2\)

Against the super-rigidity of the logical connective, $\land$, the proponent of model-theoretic validity will prefer a definition of the constant according to which, for propositions $\phi$ and $\psi$ and a model, $M$, $M$ validates $\phi \land \psi$ iff $M$ validates $\phi$ and $M$ validates $\psi$. By contrast, the proponent of proof-theoretic validity will prefer a distinct definition which makes no reference to truth, according to which $\land$ is defined by its introduction and elimination rules: $\phi, \psi \vdash \phi \land \psi$; $\phi \land \psi \vdash \phi$; $\phi \land \psi \vdash \psi$.

Finally, terms for physical entities such as 'tensor field' might have a rigid intension mapping the term to the same extension in metaphysical modal space, and a non-rigid intension mapping the term to distinct extensions in epistemically possible space, such that what is known about the term is contingent and might diverge from its necessary metaphysical profile.\(^3\) That physical terms are not super-rigid might be one way to challenge the soundness of the conceivability argument to the effect that, if it is epistemically possible that truths about consciousness cannot be derived from truths about physics, then the dissociation between phenomenal and physical truths is metaphysically possible (cf. Chalmers, 2010: 151).

Crucially for the purposes of this note, there appear to be no clear counterexamples to the claim that mereological parthood is super-rigid. If this is correct, then mereological parthood in the space of epistemic modality can serve as a guide to the status of mereological parthood in metaphysical modal space. The philosophical significance of the foregoing is that it belies the contention proffered by Roca-Royes (op. cit.) and Chalmers (op. cit.) concerning the limits of conceivability-based modal epistemology. The super-rigidity of the parthood relation ensures that the interaction between the conceivability of mereological parthood, which records the existence of haecceities, can serve as a guide to the metaphysical modal profile of haecceity comprehension. I argue further that essential properties are super-rigid, such that the conceivability of essential properties obtaining can be a guide to the metaphysical possibility of essential properties obtaining.

\(^2\)Cf. Zermelo (1908/1967) and von Neumann (1923/1967). Well-orderings are irreflexive, transitive, binary relations on all non-empty sets, defining a least element in the sets.

\(^3\) A 'tensor field' is a function from $m$ '1-forms' at a spacetime point, $p$, and $n$ vectors at $p$, to the real numbers. A 1-form is a function, $\omega$, s.t. $\omega$ maps four vectors to the real numbers, and satisfies the condition that for vectors $\geq 2$, $\mu$, $\tau$, and real numbers $\alpha$ and $\beta$: $\omega(\alpha \mu + \beta \tau) = \alpha \omega(\mu) + \beta \omega(\tau)$. Cf. Arntzenius (2012): 72.
6.3 Two Dogmas of Semantic Rationalism

The tenability of the foregoing depends upon whether objections to what might be understood as the two dogmas of semantic rationalism can be circumvented.\textsuperscript{4}

6.3.1 The First Dogma

The first dogma of semantic rationalism mirrors Quine’s (1951) contention that one dogma of the empiricist approach is the distinction that it records between analytic and synthetic claims. The analogous dogma in the semantic rationalist setting is that a distinction can be drawn between linguistic epistemic intensions – witnessed by differences in the cognitive significance of two sentences or terms which have the same extension, e.g., with \( x = 2 \), ‘\( x^2 \)’ and ‘\( 2x \)’ – by contrast to conceptual epistemic intensions – e.g., those which denote the properties of phenomenal consciousness. The distinction coincides with two interpretations of two-dimensional intensional semantics. As noted, the epistemic interpretation of two-dimensional intensional semantics takes the value of a formula relative to a context ranging over epistemically possible worlds to determine the extension of the formula relative to an index ranging over metaphysically possible worlds (cf. Chalmers, op. cit.). According to the metasemantic interpretation, a sentence, such as that ‘water = H\(_2\)O’, is metaphysically necessary, whereas assertions made about metaphysically necessary sentences record the non-ideal epistemic states of agents and are thus contingent (cf. Stalnaker, 1978, 2004). The first dogma is thus to the effect that there are distinct sets of worlds – sets of non-linguistic conceptual possibilities and of linguistic presuppositions, respectively – over which the context ranges in the epistemic and metasemantic interpretations.

If no conditions on the distinctness between linguistic and conceptual epistemic intensions can be provided, then variance in linguistic intension might adduce against the uniqueness of the conceptual intension. Because of the possible proliferation of epistemic intensions, conditions on the super-rigidity of the formulas and terms at issue might thereby not be satisfiable. The significance of the first dogma of semantic rationalism is that it guards against the collapse of conceptual and linguistic epistemic intensions, and thus the collapse of language and thought.

\textsuperscript{4}Thanks to Josh Dever for the objections.
A defense of the first dogma of semantic rationalism might, in response, be proffered, in light of the status of higher-order distributive plural quantification in natural language semantics. Plural quantifiers are distributive, if the individuals comprising the plurality over which the quantifier ranges are conceived of singly, rather than interpreting the quantifier such that it ranges over irreducible collections. Natural language semantics permits plural quantification into both first and second-level predicate position. However, there are no examples of plural quantification into third-level predicate position in empirical linguistics, despite that examples thereof can be readily countenanced in intended models of formal languages. As follows, higher-order plural quantification might adduce in favor of the first dogma of semantic rationalism, to the effect that linguistic and conceptual epistemic intensions can be sufficiently distinguished.

6.3.2 The Second Dogma

The second dogma of semantic rationalism mirrors Quine’s (op. cit.) contention that another dogma of empiricism is the reduction of the meaning of a sentence to the empirical data which verifies its component expressions. The analogous dogma in the semantic rationalist setting states that individuation-conditions on concepts can be provided in order to distinguish between concepts unique to formal and informal domains. The significance of the second dogma of semantic rationalism is that whether the objects falling under a concept belong to a formal domain of inquiry will subsequently constrain its modal profile.

In the space of epistemic possibility, it is unclear, e.g., what reasons there might be to preclude implicit definitions such as that the real number of the x’s is identical to Julius Caesar (cf. Frege, 1884/1980: 56; Clark, 2007) by contrast to being identical to a unique set of rational numbers as induced via Dedekind cuts. It is similarly unclear how to distinguish, in the space of epistemic possibility, between formal and informal concepts, in order to provide a principled account of when a concept, such as the concept of ‘set’, can be defined via the axioms of the language in which it figures, by contrast to concepts such as ‘water’, where definitions for the latter might target the observational, i.e. descriptive and functional, properties thereof. Scrutability from a canonical description of an epistemically possible world i.e. scenario, characterized by the set of truths, PQTI – i.e. physical, phenomenal, and indexical truths and a ‘that’s-all’ truth – to an arbitrary sentence, fixes an
epistemic intension (Chalmers, 2012: 245). However, physical, phenomenal, and indexical truths are orthogonal to truths about necessarily non-concrete objects such as abstracta.\footnote{For challenges to the indexing account of mathematical explanation, see Baker and Colyvan (2011). For more on mathematical explanation and its relation to scientific truths, see Mancosu (2008); Pincock (2012); Lange (2017); and Baron et al (2020).} How then are the epistemic intensions for abstracta fixed? The most obvious maneuver would be to add mathematical truths to the scrutability base from which sentences about mathematical objects can be inferred. Doing so leaves a number of questions open, however. In this section, I thus provide an explanation of how formal and informal domains can be distinguished, determining thereby the modal profiles of the relevant concepts.

The concept of mereological parthood provides a borderline case. While the parthood relation can be axiomatized so as to reflect whether it is irreflexive, non-symmetric, and transitive, its status as a formal property is more elusive. The fact, e.g., that an order-type is part of the sequence of ordinal numbers impresses as being necessary, while yet the fact that a number of musicians comprise the parts of a chamber ensemble might impress as being contingent.

The Julius Caesar problem, and the subsequent issue of whether there might be criteria for delineating formal from informal concepts in the space of epistemic modality, may receive a unified response. The ambiguity with regard to whether the parthood relation is formal – given that its relata can include both formal and informal objects – is similar to the ambiguity pertaining to the nature of real numbers. As Frege (1893/2013: 161) notes: 'Instead of asking which properties an object must have in order to be a magnitude, one needs to ask: how must a concept be constituted in order for its extension to be a domain of magnitudes [...] a thing is a magnitude not in itself but only insofar as it belongs, with other objects, to a class that is a domain of magnitudes’. Frege defines a magnitude as the extension of a relation on arbitrary domains (op. cit.). The concept of a magnitude is then referred to as a 'Relation', and domains of magnitudes are defined as classes of Relations (162). Bypassing the rational numbers, Frege defines, then, the real numbers as relations on – namely, ratios of – magnitudes; and thus refers to the real numbers as 'Relations on Relations', because the extension of the higher-order concept of real number is taken to encompass the extension of the lower-order concept of classes of Relations, i.e., domains of magnitudes...
The interest of Frege’s definition of the concept of real number is that explicit mention must be made therein to a domain of concrete entities to which the number is supposed, as a type of measurement, to be applied.

In response: The following implicit definitions—i.e., abstraction principles—can be provided for the concept of real number, where the real numbers are defined as sets, or Dedekind cuts, of rational numbers. Following Shapiro (2000), let $F, G,$ and $R$ denote rational numbers, such that concepts of the reals can be specified as follows: $\forall F, G[\text{C}(F) = \text{C}(G) \iff \forall R(F \leq R \iff G \leq R)]$. Concepts of rational numbers can themselves be obtained via an abstraction principle in which they are identified with quotients of integers—$[Q(m,n) = Q(p,q) \iff n = 0 \land q = 0 \lor n \neq 0 \land q \neq 0 \land m \times q = n \times p]$; concepts of the integers are obtained via an abstraction principle in which they are identified with differences of natural numbers—$[D(\langle x,y \rangle) = D(\langle z,w \rangle) \iff x + w = y + z]$; concepts of the naturals are obtained via an abstraction principle in which they are identified with pairs of finite cardinals—$\forall x, y, z, w[\langle x, y \rangle (=P) = \langle z, w \rangle (=P) \iff x = z \land y = w]$; and concepts of the cardinals are obtained via Hume’s Principle, to the effect that cardinals are identical if and only if they are equinumerous—$\forall A \forall B \exists R[\forall x: A = Nx: B \equiv \exists R[\forall x[Ax \to \exists y(By \land Rx y \land \forall z(Bz \land Rx z \to y = z)]]) \land \forall y[By \to \exists x(Ax \land Rx y \land \forall z(Az \land Rx z \to x = z)]])$.

Frege notes that ‘we can never [...] decide by means of [implicit] definitions whether any concept has the number Julius Caesar belonging to it, or whether that same familiar conqueror of Gaul is a number or not’ (1884/1980: 56). A programmatic line of response endeavors to redress the Julius Caesar problem by appealing to sortal concepts, where it is an essential property of objects that they fall in the extension of the concept (cf. Hale and Wright, 2001: 389, 395). In order further to develop the account, I propose to avail of recent work in which identity conditions are interpreted so as to reflect relations of essence and explanatory ground. The role of the essentiality operator will be to record a formal constraint on when an object falls under a concept ‘in virtue of the nature of the object’ (Fine, 1995: 241-242). The role of the grounding operator will be to record a condition on when two objects are the same, entraining a hyperintensional type of implicit definition for concepts which is thus finer-grained and less susceptible to error through misidentification.

In his (2015), Fine treats identity criteria as generic statements of ground. By contrast to material identity conditions which specify when two objects are identical, criterial identity conditions explain in virtue of what the two
objects are the same. Arbitrary, or generic, objects are then argued to be constitutive of criterial identity conditions. Let a model, \( M \), for a first-order language, \( L \), be a tuple, where \( M = \langle I, A, R, V \rangle \), with \( I \) a domain of concrete and abstract individuals, \( A \) a domain of arbitrary objects, \( R \) a dependence relation on arbitrary objects, and \( V \) a non-empty set of partial functions from \( A \) to \( I \) (cf. Fine, 1985). The arbitrary objects in \( A \) are reified variables. The dependence relation between any \( a \) and \( b \) in \( A \) can be interpreted as a relation of ontological dependence (op. cit.: 59-60). Informally, from \( a \in A \) s.t. \( F(a) \), one can infer \( \forall x. F(x) \) and \( \exists x. F(x) \), respectively (57). Then, given two arbitrary objects, \( x \) and \( y \), with an individual \( i \) in their range, ‘\( [(x = i \land y = i) \rightarrow x = y] \)’, such that \( x \) and \( y \) mapping to a common individual explains in virtue of what they are the same (Fine, 2015).

Abstraction principles for, e.g., the notion of set, as augmented so as to record distinctions pertaining to essence and ground, can then be specified as follows:

- Given \( x, y \), with \( \text{Set}(x) \land \text{Set}(y) \): 
  \[ \forall z (z \in x \equiv z \in y) \leftrightarrow_{x,y} (x = y) \]

  (Intuitively, where the ‘given’ expression is a quantifier ranging over the domain of variables-as-arbitrary objects: Given \( x, y \), whose values are sets, it is essential to \( x \) and \( y \) being the same that they share the same members); and

- Given \( x, y \), with \( \text{Set}(x) \land \text{Set}(y) \): 
  \[ \forall z (z \in x \equiv z \in y) \rightarrow_{x,y} (x = y) \]

  (Intuitively: Given arbitrary objects, \( x, y \), whose values are sets, the fact that \( x \) and \( y \) share the same members grounds the fact that they are the same).

Combining both of the above directions yields the following hyperintensional, possibly asymmetric, biconditional:

- Given \( x, y \), with \( \text{Set}(x) \land \text{Set}(y) \): 
  \[ \forall z (z \in x \equiv z \in y) \leftrightarrow_{x,y} (x = y) \].

A reply to the Julius Caesar problem for real numbers might then avail of the foregoing metaphysical implicit definitions, such that the definition would record the essentiality to the reals of the property of being necessarily non-concrete as well as provide a grounding-condition:

- Given \( F, G [C(F) = C(G) \leftrightarrow_{F,G} \forall R (F \leq_R \leftrightarrow G \leq_R)] \), and
\[ \forall X/F \exists Y [\neg E(Y) \land \Box(X = Y)] \]

(Intuitively: Given arbitrary objects, F, G, whose values are the real numbers: It is essential to the F’s and the G’s that the concept of the Fs is identical to the concept of the G’s iff (i) F and G are identical subsets of a limit rational number, R, and (ii) with E(x) a concreteness predicate, necessarily for all real numbers, X, necessarily there is a non-concrete object Y, to which necessarily X is identical; i.e., the reals are necessarily non-concrete. The foregoing is conversely the ground of the identification.)\(^6\)

Heck (2011: 129) notes that the Caesar problem incorporates an epistemological objection: "Thus, one might think, there must be more to our apprehension of numbers than a mere recognition that they are the references of expressions governed by HP [Hume’s Principle – HK]. Any complete account of our apprehension of numbers as objects must include an account of what distinguishes people from numbers. But HP alone yields no such explanation. That is why Frege writes: ‘Naturally, no one is going to confuse [Caesar] with the [number zero]; but that is no thanks to our definition of [number]' (Gl, 62)."

The condition of being necessarily non-concrete in the metaphysical definition for real numbers – which includes conditions of essence and ground – provides a reply to the foregoing epistemological objection, i.e. the required account, beyond the abstraction principle, of what distinguishes people from numbers.

### 6.3.3 Mereological Parthood

The above proposal can then be generalized, in order to countenance the abstract profile of the mereological parthood relation. By augmenting the axioms for parthood in, e.g., classical mereological parthood with a clause to the effect that it is essential to the parthood relation that it is necessarily non-concrete, parthood can thus be understood to be abstract; and truths in which the relation figures would thereby be necessary.

\(^6\)Rosen and Yablo (2020) also avail of real, or essential, definitions in their attempt to solve the Caesar problem, although their real definitions do not target grounding-conditions. The need for a grounding-condition is mentioned in Wright (2020: 314, 318). The approach here developed, of solving the Caesar problem by availing of metaphysical definitions, was arrived at independently of Rosen and Yablo (op.cit.) and Wright (op. cit.).
• Given $x$: $\Phi(x) \land \Box \forall x \exists y [\neg E(y) \land \Box (x = y)] \leftrightarrow \Gamma(x)$ where

$\Gamma(x) := x$ is the parthood relation, $<$, which is irreflexive, asymmetric, and transitive, and where the relation satisfies the axioms of classical extensional mereology codified by the predicate, $\Phi(x)$ (cf. Cotnoir, 2014):

**Weak Supplementation:** $x < y \rightarrow \exists z [(z < y \lor z = y) \land \neg \exists w (w < z \lor w = x)]$, and

**Unrestricted Fusion:** $\forall xx \exists y [F(y, xx)]$,

with the axiom of Fusion defined as follows:

**Fusion:** $F(t, xx) := (xx < t \lor xx = t) \land \forall y [(y < t \lor y = t) \rightarrow (y < xx \lor y = xx)]$

Fusions are themselves abstracta, formed by a fusion-abstraction principle. The abstraction principle states that two singular terms – in which an abstraction operator, $\sigma$, from pluralities to fusions figures as a subformula – are identical, if and only if the fusions overlap the same locations (cf. Cotnoir, ms). Let a topological model be a tuple, comprised of a set of points in topological space, $\mu$; a domain of individuals, $D$; an accessibility relation, $R$; and a valuation function, $V$, assigning distributive pluralities of individuals in $D$ to subsets of $\mu$:

$M = \langle \mu, D, R, V \rangle$;

$R = R(xx, yy)_{xx, yy \in \mu}$ iff $R_{xx} \subseteq \mu_{xx} \times \mu_{xx}$, s.t. if $R(xx, yy)$, then $\exists o \subseteq \mu$, with $xx \in o$ s.t. $\forall yy \in o R(xx, yy)$, where the set of points accessible from a privileged node in the space is said to be open; and $V = f(ii \in D, m \in \mu)$.

Necessity is interpreted as an interiority operator on the space:

$M, xx \models \Box \phi$ iff $\exists o \subseteq \mu$, with $xx \in o$, such that $\forall yy \in o M, yy \models \phi$.

The following fusion abstraction principle can then be specified:

Given $xx, yy, F[\sigma(xx, F) = \sigma(yy, F) \leftrightarrow_{xx, yy} [f(xx, m_1) \cap f(yy, m_1) (\neq \emptyset)]]$.

(Intuitively, given arbitrary objects whose values are the pluralities, $xx, yy$: It is essential to $xx$ and $yy$ that fusion-abstracts – formed by mapping the pluralities to the abstracta – are identical, because the fusions overlap the same nonstationary – i.e., $\neq \emptyset$ – locations. The converse is the determinative ground of the identification.)

The foregoing constraints on the formality of the parthood relation – both being necessarily non-concrete and figuring in pluralities which serve

---

$^7\mu$ is further Alexandrov; i.e., closed under arbitrary unions and intersections.
to individuate fusions as abstract objects – are sufficient then for redressing the objections to the second dogma of semantic rationalism; i.e., that individuation-conditions are wanting for concepts unique to formal and informal domains, which would subsequently render the modal profile of such concepts indeterminate. That relations of mereological parthood are abstract adduces in favor of the claim that the values taken by the relation are necessary. The significance of both the necessity of the parthood relation, as well as its being abstract rather than concrete and thus being in some sense apriori, is that there are thus compelling grounds for taking the relation to be super-rigid, i.e., to be both epistemically and metaphysically necessary.

Finally, a third issue, related to the second dogma is that, following Dummett (1963/1978: 195-196), the concept of mereological parthood might be taken to exhibit a type of ‘inherent vagueness’, in virtue of being indefinitely extensible. Dummett (1996: 441) defines an indefinitely extensible concept as being such that: ‘if we can form a definite conception of a totality all of whose members fall under the concept, we can, by reference to that totality, characterize a larger totality all of whose members fall under it’. It will thus be always possible to increase the size of the domain of elements over which one quantifies, in virtue of the nature of the concept at issue; e.g., the concept of ordinal number is such that ordinals can continue to be generated, despite the endeavor to quantify over a complete domain, in virtue of iterated applications of the successor relation, and the concept of real number is such that the reals can continue to be generated via elementary embeddings. Bernays’ (1942) theorem states that class-valued functions from classes to sub-classes are not onto, where classes are non-sets (cf. Uzquiano, 2015a: 186-187). A generalization of Bernays’ theorem can be recorded in plural set theory, where the cardinality of the sub-pluralities of an incipient plurality will always be greater than the size of that incipient plurality. If one takes the cardinal height of the cumulative hierarchy to be fixed, then one way of tracking the variance in the cardinal size falling in the extension of the concept of mereological parthood might be by redefining the intension thereof (Uzquiano, 2015b). Because it would always be possible to reinterpret the concept’s intension in order to track the increase in the size of the plural universe, the intension of the concept would subsequently be non-rigid; and

---

the concept would thus no longer be super-rigid.

One way in which the objection might be countered is by construing the variance in the intension of the concept of parthood as tracking temporal modal properties, rather than metaphysical modal properties. Then, the relation can be necessary while satisfying full S5 – i.e., modal axioms K [(□φ → ψ) → (□φ → □ψ)], T (□φ → φ), and E (¬□φ → □¬□φ) – despite that there can be variations in the size of the quantifier domains over which the relation and its concept are defined. Let ↑ be an intensional parameter which indexes and stores the relevant formulas at issue to a particular world (cf. Hodes, 1984). The ↓-symbol is an operator which serves to retrieve, as it were, that indexed information. Adding multiple arrows is then akin to multiple-indexing: The value of a formula, as indexed to a particular world, will then constrain the value of that formula, as indexed – via the addition of the new arrows – to different worlds. Interpreting the operators temporally permits there to be multiple-indexing in the array of intensional parameters relative to which a formula gets its value, while the underlying logic for metaphysical modal operators can be S5, partitioning the space of worlds into equivalence classes. Formally:

\[ \uparrow_1 \forall x \exists \downarrow_1 \phi \uparrow_2 \exists y [\phi(x) \downarrow_1 \land \phi(y) \downarrow_2]. \]

The clause states that, relative to a first temporal parameter in which all of the x’s satisfying the sethood predicate are quantified over, there is – relative to a distinct temporal parameter – another element which satisfies that predicate. Crucially, differences in the intensional temporal indices, as availed of in order to record variance, at different times, in the size of the cumulative hierarchy of elements falling in the range of the parthood relation, is yet consistent with the cardinality of the elements in the domain falling in the range of the relation being fixed, such that the valuation of the relation can yet be necessary.

6.3.4 Summary

In this section, I addressed objections to two dogmas of the semantic rationalism underpinning the epistemic interpretation of two-dimensional intensional semantics. In response to objections to the first dogma – according to which no distinctions can be delineated between linguistic intensions and conceptual epistemic intensions – I noted that higher-order plural terms are conceptually tractable although they have no analogue in natural language semantics. In response to objections to the second dogma – according to which criteria on
distinguishing formal from informal domains unique to the extensions of various concepts are lacking, which subsequently engenders indeterminacy with regard to the modal profiles of those concepts – I availed of generic criterial identity conditions, in which it is essential to identical arbitrary representatives of objects that they satisfy equivalence relations which are conversely ground-theoretically determinative of the identification, and further essential thereto that they satisfy the predicate of being necessarily non-concrete. The extensions of indefinitely extensible concepts can further be redefined relative to distinct temporal intensional parameters, despite that the background modal logic for the intensions of the concepts partitions the domain of worlds into equivalence classes, and thus satisfies S5. Thus, parthood can be deemed a necessary, because abstract, relation, despite (i) temporal variance in the particular objects on which the parthood relation is defined; and (ii) variance in the cardinality of the domain in which those objects figure, relative to which the concept’s intensions are defined.

When \( \Phi = x \sqsubseteq xx, \Box x y \exists x (\Phi x \iff x = y) \). By the super-rigidity of the parthood relation, the target two-dimensional intensional formula can, finally, be stated as follows:

If it is epistemically possible that \( \Phi x \), then it is metaphysically possible that \( \Phi x \). Formally:

\[ \forall c, w \in W[\Phi x]_{c, w} = 1 \iff \exists c', w' \in W[\Phi x]_{c', w'} = 1. \]

Thus, the epistemic possibility of haecceity comprehension constrains the value of the metaphysical possibility of haecceity comprehension, and – in response to Roca-Royes and Chalmers – there is a case according to which conceivability is a guide to a principle of modal metaphysics.

There can be a generalization of the foregoing approach to the essential properties of individuals, as well, assuming that such properties are super-rigid. Following Fine (2000), suppose there is an operator, \( \Box_F \), where \( \Box_F A \) is read 'it is true in virtue of the nature of the nature of (some or all) of the F’s that A' where 'each of the objects mentioned in A is involved in the nature of one of the F’s' (op. cit.: 543). \( \Box_F \) satisfies the axioms KTE and necessitation:

\[ \Box_F A \rightarrow A, \]
\[ \Box_F (A \rightarrow B) \rightarrow (\Box_F A \rightarrow \Box_F B), \]
\[ \neg \Box_F A \rightarrow \Box_F \neg A, \text{ F rigid}, \]

F is rigid "if it is a rigid predicate symbol or is of the form \( \lambda x \bigwedge_{1 \leq i \leq n} A_i, n \geq 0 \), where each formula \( A_i \), \( i = 1, \ldots, n \), is either of the form \( P x \) or of the form \( x = y \) for some variable y distinct from \( x' \) (545), and
|E| stands for $\lambda x(x\eta E) \times$ the first variable not free in E, where $x\eta E$ stands for $\bigvee_{1 \leq i \leq m} x_i \vee \bigvee_{1 \leq i \leq m} P_j x_i$,

A $\vdash [\Box A]$, and

$F \subset G \rightarrow (\Box_F A \rightarrow \Box_G A)$ (546).

A model $M$ is a quadruple $\langle W, I, \preceq, \phi \rangle$ where $W$ is a non-empty set of worlds, $I$ is a function taking each $w \in W$ into a non-empty set of individuals $I_w$, $\preceq$ is a reflexive transitive dependence relation on $\bigcup_{w \in W}$ with respect to which each world is closed ($a \in I_w$ and $a \preceq b$ implies $b \in I_w$), and $\phi$ is a valuation function taking each constant $a$ into an individual $\phi(a)$ of some $I_w (w \in W)$, each rigid predicate symbol $H$ into a subset $\phi(H)$ of some $I_w$, and each world $w$ and pure $n$-place predicate symbol $F$ into a set $\phi(F, w)$ of $n$-tuples of $I_w$, where a pure predicate involves no reference to any object (544, 547-548).

For a subset $J$ of $\bigcup I_w$, the closure $c(J)$ of $J$ in $M$ is $\{b: a \preceq b$ for some $a \in J\}$ (548).

$M$ is a model with $E$ a sentence or closed predicate whose constants are $a_1, \ldots, a_m$ and whose rigid predicate symbols are $P_1, \ldots, P_n$ (op. cit.). The objectual content $[E]^M$ of $E$ in $M$ is then $\{\phi(a_1), \ldots, \phi(a_m)\} \cup \{\phi(P_1), \ldots, \phi(P_n)\}$ and $E$ is defined in $M$ at $w \in W$ if $[E]^M \subseteq I_w$ (op. cit.).

Then the semantics for $\Box_F$ can be defined as follows:

$w \models \Box_F A$ iff (i) $[A]^M \subseteq c(F_w)$, and (ii) $v \models A$ whenever $I_v \supseteq F_w$, where $F_w$ is $\phi(w, F)$ (op. cit.).

$\Box_F$ can be defined relative two parameters, the first ranging over epistemically possible worlds considered as actual, and the second ranging over metaphysically possible worlds, such that the conceivability of it being true in virtue of the nature of the nature of (some or all) of the $F$’s that $A$ entails the metaphysical possibility of it being true in virtue of the nature of the nature of (some or all) of the $F$’s that $A$:

$\forall c \in C, w \in W[\Box_F A]_{c, w} = 1$ iff $\exists c' \in C, w' \in W[\Box_F]_{c', w'} = 1$.

In the remainder of the paper, I will examine issues pertaining to the determinacy of epistemic possibilities.

6.4 Determinacy and Consistency

In his (2014), Chalmers argues for the law of excluded middle, such that it is either apriori derivable using the material conditional – i.e. ‘scrutable’ – that p or scrutable that $\neg p$, depending on the determinacy of p. Chalmers refers to
the case in which \( p \) must be determinate, entailing determinate scrutability, as the Hawthorne model, and the case in which it can be indeterminate, entailing indeterminate scrutability, as the Dorr model (259).\textsuperscript{9} Chalmers argues that, for any \( p \), one can derive \( 'p \iff \text{it is scrutable that } p' \) from \( 'p \iff \text{it is true that } p' \) (262). However, \( 'p \iff \text{it is scrutable that } p' \) is unrestrictedly valid only on Dorr’s, and not Hawthorne’s, model (op. cit.).\textsuperscript{10}

Chalmers suggests that the relevant notion of consistency might be a property of epistemic possibilities rather than metaphysical possibilities. However, there are general barriers to establishing the consistency of the space of epistemic modality.

One route to securing the epistemic interpretation of consistency is via Chalmers’ conception of idealized epistemic possibility. Conceivability is ideal if and only if nothing rules it out apriori upon unbounded rational reflection (2012: 143). The rational reflection pertinent to idealized conceivability can be countenanced modally, normatively, and so as to concern the notion of epistemic entitlement. An idealization is (i) modal iff it concerns what it is metaphysically possible for an agent to know or believe; (ii) normative iff it concerns what agents ought to believe; and (iii) warrant-involving iff it concerns the propositions which agents are implicitly entitled to believe (2012: 63). It is unclear whether any of (i)-(iii) in the foregoing would either mandate belief in the claim that \( p \land \text{it is indeterminate whether } p' \) is true, or explain in virtue of what the conjuncts are consistent. More general issues for the consistency of epistemically possible worlds, even assuming that the idealization conditions specified in (i)-(iii) are satisfied, include Yablo’s (1993) paradox, and Gödel’s (1931) incompleteness theorems. Yablo’s paradox is as

\textsuperscript{9}Cf. Dorr (2003: 103-4) and Hawthorne (2005: sec. 2).

\textsuperscript{10}Chalmers rejects the epistemicist approach to indeterminacy, which reconciles the determinacy in the value of a proposition with the epistemic indeterminacy concerning whether the proposition is known (op. cit.: 288). Consider, e.g., a color continuum, beginning with a determinate color hue of red and terminating with a determinate color hue of orange. By transitivity, if the determinate hue of red, \( x \), is phenomenally similar to the next point in the continuum, \( y \), and \( y \) is phenomenally similar to the next point, \( z \), then \( x \) is phenomenally similar to \( z \). However, iterating transitivity would entail that the terminal color hue is red and not orange. Thus, if the culprit in the sorites paradox is the property of transitivity, then the modal axiom which encodes transitivity (namely 4: \( \Box\phi \rightarrow \Box\Box\phi \)) is false. The epistemic interpretation of the axiom states that if one knows that \( \phi \), then one knows that one knows that \( \phi \). Thus, rejecting axiom 4 entails that the cut-off points in a sorites series are knowable, although one cannot know that one knows them. For further discussion, see Williamson (1994).
follows:

(S1) For all \( k > 1 \), \( S_k \) is false;
(S2) For all \( k > 2 \), \( S_k \) is false;
...

(Sn) For all \( k > n \), \( S_k \) is false;
(Sn+1) For all \( k > n + 1 \), \( S_k \) is false.

(Sn) says that (Sn+1) is false. Yet (Sn+1) is true. Contradiction.\(^{11}\)

Gödel’s incompleteness theorems can be thus outlined. Relative to a choice of (i) coding for an \( \omega \)-complete, recursively axiomatizable language, \( L \) – i.e. a mapping between properties of numbers and properties of terms and formulas in \( L \); (ii) a predicate, phi; and (iii) a fixed-point construction:

Let phi express the property of 'being provable', and define (iii) s.t., for any consistent theory \( T \) of \( L \), there are sentences, \( p_{phi} \), corresponding to each formula, phi(x), in \( T \), s.t. for \('m' := p_{phi}, \nabs{T} p_{phi} \iff \phi(m).\)

One can then construct a sentence, \( 'm' := \neg \phi(m) \), such that \( L \) is incomplete (the first incompleteness theorem).

Crucially, moreover, \( L \) cannot prove its own consistency:

If:
\nabs{T} 'm' \iff \neg \phi(m),

Then:
\nabs{T} C \rightarrow m.

So, \( L \) is consistent only if \( L \) is inconsistent (the second incompleteness theorem).

Another issue concerning the consistency of \( p \land \) it is indeterminate whether \( p \) – let alone the foregoing general issues concerning the consistency of epistemic modal space – is that Chalmers (2009: 102) endorses the indeterminacy of metaphysical proposals such as unrestricted fusion and, presumably, the necessity of parthood, with regard to which the epistemic interpretation of consistency would be irrelevant (264).

To redress the issue, the metaphysical indeterminacy of ontological proposals might be treated as in Barnes and Williams (2011), for whom metaphysical indeterminacy consists in there being an unpointed set of metaphysically possible worlds; i.e., a set of metaphysical possibilities, \( P \), such that precisifications concerning the determinacy in the values of the elements of

\(^{11}\)For further discussion, see Cook (2014).
P leave it unsettled which possibility is actual (116, 124). If so, then metaphysical indeterminacy will provide no new objection to the viability of the two-dimensional intensional framework, because the conditions on ascertaining the actuality of the epistemic possibility in the context – relative to which a formula receives a value, and thus crucially determines the value of the formula relative to an index which ranges over metaphysically possible worlds – are themselves indeterminate (cf. Yablo, 2008).

The more compelling maneuver might instead be to restrict the valid apriori material entailments to determinately true propositions; and to argue, against Chalmers’s preferred ontological anti-realist methodology, that the necessity of parthood is both epistemically and metaphysically determinately true, if true at all. The (determinate) truth of the proposition might then be corroborated both by the consistency of its augmentation to the logic underlying the intensional semantics, and perhaps in virtue of other abductive criteria – such as strength, simplicity, and compatibility with what is known – on the tenability of the proposal.

6.5 Concluding Remarks

One of the primary objections to accounting for the relationship between conceivability and metaphysical possibility via the epistemic interpretation of two-dimensional intensional semantics is that epistemic possibilities are purportedly insensitive to modal metaphysical propositions, concerning, e.g., the haecceitistic properties of individuals and essential properties. In this essay, I have endeavored to redress the foregoing objection. Further objections, from both the potential indeterminacy in, and inconsistency of, the space of epistemic possibilities, were then shown to be readily answered. In virtue of the super-rigidity of the parthood relation and essential properties, conceivability can thus serve as a guide to haecceity comprehension principles in modal metaphysics as well as to the obtaining of essential properties.
Chapter 7

Grounding, Conceivability, and the Mind-Body Problem

This essay argues that Chalmers’s (1996; 2010) two-dimensional conceivability argument against the derivation of phenomenal truths from physical truths risks being obviated by a hyperintensional regimentation of the ontology of consciousness.

Chalmers (2010a) provides the following argument against the identification of phenomenal truths with physical and functional truths. Let M be a model comprised of a domain D of formulas; C a set of epistemic possibilities; W a set of metaphysical possibilities; R_c and R_w, accessibility relations on C and W, respectively; and V a valuation function assigning formulas to subsets of C and W. So, M = ⟨D,C,W,R_c,R_w,V⟩. Let P denote the subset of formulas in the domain concerning fundamental physics, as well as both neurofunctional properties such as oscillations of neural populations, and psychofunctional properties such as the retrieval of information from memory stores. Let Q denote the subset of formulas in the domain concerning phenomenal consciousness. A formula is epistemically necessary or apriori (□), if and only if it has the same value at all points in C, if and only if it is impossible, i.e. inconceivable, for the formula to a variant value (¬⋄¬). A formula is negatively conceivable (⋄) if and only if nothing rules it out apriori (¬□¬) (144). A formula is metaphysically necessary if and only if it has the same value at all points in W. A formula is said to be ‘super-rigid’, if and only if it is both epistemically and metaphysically necessary, and thus has the same value at all points in epistemic and metaphysical modal space (2012: 474).
The two-dimensional conceivable argument against physicalism proceeds as follows.

The physicalist thesis states that:
\[ P \rightarrow Q. \]
Suppose, however, that the physicalist thesis is false. Thus,
1. \[ \neg(P \rightarrow Q). \]
By the definition of the material conditional,
2. \[ \neg(\neg P \lor Q). \]
By the De Morgan rules for negation,
3. \[ \neg\neg P \land \neg Q. \]
By double negation elimination,
4. \[ P \land \neg Q. \]

‘\( P \land \neg Q \)’ can receive a truth value relative to two parameters, a context, \( C \), and an index, \( W \). In two-dimensional intensional semantics, the value of the formula relative to the context determines the value of the formula relative to the index. Let the context range over a space of epistemic possibilities and let the index range over a space of metaphysical possibilities. Then,

\[ [P \land \neg Q]^{c,w} = 1 \iff \exists c' \in C \exists w' \in W [P \land \neg Q]^{c',w'} = 1. \]

The foregoing clause codifies the thought that, if it is epistemically possible that the truths about physics and functional organization obtain while the truths about consciousness do not, then the dissociation between \( P \) and \( Q \) is metaphysically possible as well. The argument depends on the assumption that propositions about consciousness and physics are super-rigid, such that the epistemic possibility concerning such truths can serve as a guide to the metaphysical possibility thereof.

If the conceivable argument is sound, then the physicalist thesis – that all phenomenal truths are derivable from physical and functional truths –

---

1 For the formal equivalence, given the definition of the material conditional, see Chalmers (2010a: 169).
2 For the clause for the two-dimensional intension, see Chalmers and Rabern (2014: 212). Chalmers’ informal characterization of the argument proceeds as follows:
1. \( P \land \neg Q \) is conceivable.
2. If \( P \land \neg Q \) is conceivable, \( P \land \neg Q \) is [epistemically, i.e.] 1-possible.
3. If \( P \land \neg Q \) is 1-possible, \( P \land \neg Q \) is [metaphysically, i.e.] 2-possible.
4. If \( P \land \neg Q \) is 2-possible, then materialism is false.
Thus,
5. Materialism is false (2010: 149).
is possibly false. The foregoing argument entrains, thereby, the metaphysical possibility of a property-based version of dualism between phenomenal consciousness and fundamental physics.

One of the standard responses to Chalmers’s conceivability argument is to endeavor to argue that there are ‘strong’ necessities, i.e. cases according to which the necessity of the physical and phenomenal formulas throughout epistemic and metaphysical modal space is yet consistent with the epistemic possibility that the formulas have a different value.³ Note, however, that strong necessities are ruled-out, just if one accepts the normal duality axioms for the modal operators: i.e., it is necessary that φ if and only if it is impossible for φ to be false: □φ iff ¬¬φ. Thus, the epistemic necessity of φ rules out the epistemic possibility of not-φ by fiat. So, proponents of the strong necessity strategy are committed to a revision of the classical duality axioms.

Another line of counter-argument proceeds by suggesting that the formulas and terms at issue are not super-rigid. Against the super-rigidity of physical truths, one might argue, for example, that our knowledge of fundamental physics is incomplete, such that there might be newly discovered phenomenal or proto-phenomenal truths in physical theories from which the truths about consciousness might be derived.⁴ More contentiously, the epistemic profile of consciousness – as recorded by the concepts comprising our thoughts thereof, or by the appearance of its instantiation – might be dissociable from its actual instantiation. A variation on this reply takes our concepts of phenomenal consciousness still to refer to physical properties (cf.

³As Chalmers (2010a: 166-167) writes, ‘Before proceeding, it is useful to clarify [the general conceivability-possibility thesis] CP by making clear what a counterexample to it would involve . . . Let us say that a negative strong necessity is a statement S such that S is [epistemically]-necessary and [metaphysically]-necessary but ¬S is negatively conceivable’. For a case-by-case examination of purported examples of strong necessities, see Chalmers (op. cit.: 170-184; 2014a). Because it is epistemically possible for there to be scenarios in which there is no consciousness, the target neighborhood of epistemically possible worlds is that in which the conditions on there being phenomenal consciousness are assumed to obtain. (Thanks here to Chalmers (p.c.).) Thus, the notion of epistemic necessity will satisfy conditions on real world validity, rather than general validity. In the latter case, a formula is necessary if and only if it has the same value in all worlds in a model. In the former case, the necessity at issue will hold throughout the neighborhood, where a neighborhood function assigns the subset of worlds in which consciousness obtains to a privileged world in the model.⁴See Seager (1995) and Strawson (2006) for the panpsychist proposal. Proponents of the pan-protopsychist approach include Stoljar (2001, 2014) and Montero (2010).
A related line of counter-argument relies on the assumption that phenomenal concepts are entities which are themselves physically reducible (cf. Balog, 1999).

Finally, a counter-argument to the conceivability argument that has yet to be advanced in the literature is that its underlying logic might be non-classical. Thus, for example – by relying on double negation elimination in the inference from line 3 to 4 above – the two-dimensional conceivability argument is intuitionistically invalid. A novel approach might further consist in arguing that epistemic modality might be governed by the Routley-Meyer semantics for relevant logic. Relevant validity can be defined via a ternary relation, such that \( [\phi \rightarrow \psi]^\alpha = 1 \) iff \( [\phi]^\beta \leq [\psi]^\gamma \) and \( R(\alpha,\beta,\gamma) \), where the parameters, \( \alpha, \beta, \) and \( \gamma \), range over epistemic possibilities. The philosophical interest of relevant logic is that it eschews the principle of disjunctive syllogism; i.e., \( \forall \phi, \psi[[((\phi \lor \psi) \land \neg \phi) \rightarrow \psi] \) and \( \forall \phi, \psi[[\phi \land (\neg \phi \lor \psi)] \rightarrow \psi]. \)

In this essay, I will pursue a line of argument which is novel and distinct from the foregoing. I argue, in turn, that the conceivability argument can be circumvented, when the relationship between the truths about fundamental physics and the truths about phenomenal consciousness is analyzed in a classical, hyperintensional setting. Suppose, for example, that the physicalist thesis is defined using hyperintensional, grounding operators rather than metaphysical necessitation. Then, the epistemic and metaphysical possibility that \( \neg (P \rightarrow Q) \) is classically valid, although targets a less fine-grained metaphysical connection between physical and phenomenal truths. Even if \( P \)'s grounding \( Q \) still entails the metaphysical necessitation of \( Q \) by \( P \), the epistemic-intensional value of \( \neg (P \rightarrow Q) \) – will be an insufficient guide to the metaphysical-hyperintensional value of the proposition. So, even if the intension for \( \text{‘consciousness’} \) is rigid in both epistemic and metaphysical modal space, the epistemic intension recording the value of the proposition will be blind to its actual metaphysical value, because the latter will be hyperintensional.

In the remainder of this essay, I will outline the regimentation of the

---

5 Cf. Routley and Meyer (1972a,b; 1973).
6 For the logic and operator-based semantics for the notion of explanatory ground, see Fine (2012b,c).
proposals in the ontology of consciousness using hyperintensional grounding operators, rather than the resources of modality and identity.\textsuperscript{7} By contrast to the modal approach underlying the conceivability argument, the hyperintensional regimentation targets the properties of reflexivity and bijective mappings, in order to countenance novel, ontological dependence relations between the properties of consciousness and physics, which are finer-grained than necessitation.\textsuperscript{8}

Following Fine (2012b,c), let a polyadic operator have a ground-theoretic interpretation, only if the profile induced by the interpretation concerns the hyperintensional truth-making connection between an antecedent set of truths or properties and the relevant consequent. Let a grounding operator be weak if and only if it induces reflexive grounding; i.e., if and only if it is sufficient for the provision of its own ground. A grounding operator is strict if and only if it is not weak. A grounding operator is full if and only if it uniquely provides the explanatory ground for a fact. A grounding operator is part if and only if it - along with other facts - provide the explanatory ground for a fusion of facts.

Combinations of the foregoing explanatory operators may also obtain: $x < y$ iff $\phi$ is a strict full ground for $\psi$; $x \leq y$ iff $\phi$ is a weak full ground for $\psi$; $x \prec y$ iff $\phi$ is a strict part ground for $\psi$; $x \preceq y$ iff $\phi$ is a weak part ground for $\psi$; $x \triangleleft y \land \neg(y \preceq x)$ iff $\phi$ is a strict partial ground for $\psi$; $x \triangleleft^* y$ iff $x_1, \ldots, x_n \leq y$, iff $\phi$ is a partial strict ground for $\psi$; $x \prec^* z$ iff $[\phi \prec^* \psi \land \psi \preceq \mu]$ iff

\textsuperscript{7}Cf. Khudairi (ms), for the regimentation and for further discussion.

\textsuperscript{8}The claim that necessitation must be present in cases in which there is grounding is open to counterexample. Because, e.g., hyperintensional dependencies can obtain in only parts of, rather than entirely within, a world, the hyperintensional dependencies need not reflect necessitation. For further discussion of the grounding-necessitation thesis, see Rosen (2010) and Skiles (2015).
φ is a part strict ground for some further fact, μ.\(^9\)

The proposals in the metaphysics of consciousness can then be regimented in the hyperintensional framework as follows.

- **Functionalism** (modally: truths about consciousness are identical to truths about neuro- or psychofunctional role):

  Functional truths (F) ground truths about consciousness (Q) if and only if the grounding operator is:
  - strict full, s.t. F < Q
  - distributive (i.e. bijective between each truth-ground and grounded truth), s.t. \( \exists f_{1-1}(F, Q) \)

- **Phenomenal Realist Type Identity** (modally: truths about consciousness are identical to truths about biological properties, yet phenomenal

\(^9\)The derivation is induced by the following proof-rules:

- **Subsumption**

  \( (<, \leq) : \)

  \[ [(x_1, \ldots, x_n < y)] \to (x \leq y) \]

  \( (<, <) : \)

  \[ [(x_1, \ldots, x_n) < y] \to (x < y) \]

  \( (\prec, \leq) : \)

  \[ (x < y) \to (x \leq y) \]

  \( (\leq, \leq) : \)

  \[ (x \leq y) \to (x \leq y) \]

- **Distributivity/Bijection:**

  \( \forall x \in X, y \in Y \)

  \[ [G[(\ldots x \ldots)(\ldots y \ldots)]], \text{s.t.} \]

  \( f_{1-1} : [x_1 \to y_1], \ldots, f_{1-1} : [x_n \to y_n] \)

104
properties are – in some sense – non-reductively real.\textsuperscript{10}

Biological truths (B) ground truths about consciousness (Q) if and only if the grounding operator is:
- strict partial, s.t. $B \preceq Q \land \neg Q \preceq B$
- distributive, s.t. $\exists f_{1-1} (B, Q)$; and
- truths about consciousness are weak part (i.e. the set partly reflexively grounds itself), s.t. $Q \preceq Q$

- Property Dualism (modally: truths about consciousness are identical neither to functional nor biological truths, yet are necessitated by physical truths):

Physical truths (P) ground truths about consciousness (Q) if and only if the grounding operator is:
- $P \preceq Q$
- non-distributive, s.t. $\neg \exists f_{1-1} (P, Q)$; and
- truths about consciousness are weak part, s.t. $Q \preceq Q$

- Panpsychism (in Non-constitutive guise: Phenomenal properties are the intrinsic realizers of extrinsic functional properties and their roles; in Constitutive guise: (i) fundamental microphysical entities are functionally specified and they instantiate microphenomenal properties, where microphenomenal properties are the realizers of the fundamental microphysical entity’s role/functional specification; and (ii) microphenomenal properties constitute the macrophenomenal properties of macrophysical entities):

Truths about consciousness (Q) ground truths about functional role (F) if and only if the grounding operator is:
- strict full, s.t. $Q < F$; and
- non-distributive, s.t. $\neg \exists f_{1-1} (Q, F)$

\textsuperscript{10}See, e.g., Smart (1959: 148-149), for an attempt to account for how phenomenal properties and biological properties can be identical, while phenomenal properties might yet have distinct higher-order properties.
The philosophical significance of the hyperintensional regimentation of the ontology of consciousness is at least three-fold. First, the regimentation permits one coherently to formulate Phenomenal Reality Type Identity. Leibniz’s law states that for all propositional variables x,y and for all properties R, x = y iff (Rx ⇐⇒ Ry). According to the Phenomenal Realist Type Identity proposal, phenomenal properties are identical to biological properties, while phenomenal properties are in some sense non-reductively real. Thus, in the modal setting, Phenomenal Realist Type Identity belies Leibniz’s law, on the assumption that the latter can be applied to intensional entities. One virtue of the hyperintensional regimentation is thus that it avoids this result, by providing a framework with the expressive resources sufficient to formulate the non-reductive Type Identity proposal.

Second, the hyperintensional grounding regimentation evinces how functionalist approaches to the ontology of consciousness can be explanatory, because the identification of phenomenal properties with functional organization can be defined via the foregoing ground-theoretic explanatory properties. Block (2015) suggests that – by contrast to Phenomenal Realist Type Identity – identifying phenomenal properties with functional roles cannot sufficiently account for the ground-theoretic explanation of the identity. Block distinguishes between metaphysical and ontological versions of physicalism. Block’s ‘ontological physicalism’ is a reductive, functionalist theory, and eschews of explanation by restricting the remit of its theory to ‘what there is’; i.e. to specifying identity statements between entities in the domain of quantification (114). By contrast, Block’s ‘metaphysical physicalism’ – namely, Phenomenal Realist Type Identity – purports to account for the nature of the entities figuring in theoretical identity statements via availing of relations of explanatory, ontological dependence (op. cit.).

Block poses the following consideration against the functionalist (117). Suppose that there is a counterpart of a human organism with isomorphic functional properties, but comprised of distinct biological properties. Suppose that the functional isomorph instantiates phenomenal properties. Block argues that the functional isomorph ‘is like us superficially, but not in any deep property that can plausibly be one that scientists will one day tell us is the physical ground of consciousness […] So there is a key question that that kind of reductive physicalism – ontological physicalism – does not ask nor answer: what is it that creatures with the same phenomenology share that grounds that phenomenology’ (op. cit.)? The foregoing does not provide an argument that the neuro- and psycho-functionalist must provide an account.
of in virtue of what phenomenal properties are instantiated. Rather, Block suggests only that functionalist proposals do not sufficiently inquire into the realizers of the functional roles that they specify. He suggests that this theoretical approach would be insufficient, if one were to seek an explanation of the psychofunctional correlations between phenomenal property types and the relevant functional roles.

The second theoretical virtue of the hyperintensional regimentation is thus that it demonstrates how Block’s analysis might be circumvented. Functionalism can be regimented within the logic of hyperintensional ground; and can therefore satisfy the formal requirements on explaining in virtue of what phenomenal truths ontologically depend upon functional truths [cf. Khudairi (op. cit.)].

Third, and most crucially: The regimentation demonstrates how metaphysically possible relations between consciousness and physics cannot be witnessed by epistemic constraints, when the latter are recorded by the conceivability – i.e., the epistemic possibility – thereof. Propositional epistemic modality is blind to the hyperintensional, metaphysical dependencies holding between phenomenal and physical truths. Thus, the two-dimensional conceivability argument against the derivation of phenomenal truths from physical truths risks being obviated by a hyperintensional regimentation of the ontology of consciousness.

One way to resolve the foregoing issue is to provide a hyperintensional semantics for epistemic space, such that epistemic space can track metaphysical space when the latter is itself hyperintensionally defined via e.g. grounding

11 Of pertinence to the foregoing is another distinction drawn by Fine (2015b), between material and criterial identity conditions. While material identity conditions imply the identity of the objects in question, criterial identity conditions explain in virtue of what the objects in question are the same. In order to countenance criterial identity conditions, Fine avails of his earlier work on arbitrary objects (2015b: 306; cf. Fine, 1985). Let a model, $M$, for a first-order language, $L$, be a tuple, where $M = \langle I, A, R, V \rangle$, with $I$ a domain of concrete and abstract individuals, $A$ a domain of arbitrary objects, $R$ a dependence relation on arbitrary objects, and $V$ a non-empty set of partial functions from $A$ to $I$ (cf. Fine, 1985). The arbitrary objects in $A$ can be conceived of as reified variables, and the dependence relation between any $a$ and $b$ in $A$ may be interpreted as a relation of ontological dependence (op. cit.: 59-60). A criterial identity condition for, e.g., sets, can then be stated as follows: Given arbitrary $x,y$, with $\text{Set}(x) \land \text{Set}(y)$: $[\forall z(z \in x \equiv z \in y) \rightarrow_{x,y} (x = y)]$. (Intuitively: Given arbitrary objects, $x$, $y$, whose values are sets, the fact that $x$ and $y$ share the same members grounds the fact that they are the same.) A crucial point of departure between the foregoing and the approach proffered in this essay is Fine’s ontology of arbitrary objects, to which the present proposal need make no appeal.
operators. Following chapter 4, we thus provide a hyperintensional epistemic two-dimensional truthmaker semantics by which conceivability can be a guide to metaphysical possibility in the hyperintensional setting. According to truthmaker semantics for epistemic logic, a modalized state space model is a tuple \( \langle S, P, \leq, v \rangle \), where \( S \) is a non-empty set of states, i.e. parts of the elements in \( A \) in the foregoing epistemic modal algebra \( U \), \( P \) is the subspace of possible states where states \( s \) and \( t \) are compatible when \( s \sqcup t \in P \), \( \leq \) is a partial order, and \( v: \text{Prop} \to (2^S \times 2^S) \) assigns a bilateral proposition \( \langle p^+, p^- \rangle \) to each atom \( p \in \text{Prop} \) with \( p^+ \) and \( p^- \) incompatible (Hawke and Özgün, forthcoming: 10-11). Exact verification \( (\vdash) \) and exact falsification \( (\dashv) \) are recursively defined as follows (Fine, 2017a: 19; Hawke and Özgün, forthcoming: 11):

\[
\begin{align*}
s \vdash p \text{ if } s \in \llbracket p \rrbracket^+ & \quad (s \text{ verifies } p, \text{ if } s \text{ is a truthmaker for } p \text{ i.e. } s \text{ is in } p \text{'s extension}); \\
s \vdash p \text{ if } s \in \llbracket p \rrbracket^- & \quad (s \text{ falsifies } p, \text{ if } s \text{ is a falsifier for } p \text{ i.e. } s \text{ is in } p \text{'s anti-extension}); \\
s \vdash \neg p \text{ if } s \vdash p & \quad (s \text{ verifies not } p, \text{ if } s \text{ falsifies } p); \\
s \vdash \neg p \text{ if } s \vdash p & \quad (s \text{ falsifies not } p, \text{ if } s \text{ verifies } p); \\
s \vdash p \land q \text{ if } \exists t, u, t \vdash p, u \vdash q, \text{ and } s = t \sqcap u & \quad (s \text{ verifies } p \text{ and } q, \text{ if } s \text{ is the fusion of states, } t \text{ and } u, t \text{ verifies } p, \text{ and } u \text{ verifies } q); \\
s \vdash p \land q \text{ if } s \vdash p \text{ or } s \vdash q & \quad (s \text{ falsifies } p \text{ and } q, \text{ if } s \text{ falsifies } p \text{ or } s \text{ falsifies } q); \\
s \vdash p \lor q \text{ if } s \vdash p \text{ or } s \vdash q & \quad (s \text{ verifies } p \text{ or } q, \text{ if } s \text{ verifies } p \text{ or } s \text{ verifies } q); \\
s \vdash p \lor q \text{ if } \exists t, u, t \vdash p, u \vdash q, \text{ and } s = t \sqcap u & \quad (s \text{ falsifies } p \text{ or } q, \text{ if } s \text{ is the state overlapping the states } t \text{ and } u, t \text{ falsifies } p, \text{ and } u \text{ falsifies } q); \\
\vdash A\phi \text{ if and only if for all } t \in P \text{ there is a } t' \in P \text{ such that } t' \sqsubseteq t \in P \text{ and } t' \vdash \phi; \\
\vdash A\phi \text{ if and only if there is a } t \in P \text{ such that for all } u \in P \text{ either } t \sqcup u \notin P \text{ or } u \vdash \phi, \text{ where } A\phi \text{ or } \square \phi \text{ denotes the apriority of } \phi.
\end{align*}
\]
In order to account for two-dimensional indexing, we augment the model, M, with a second state space, S*, on which we define both a new parthood relation, \( \leq^* \), and partial function, V*, which serves to map propositions in a domain, D, to pairs of subsets of S*, \( \{1,0\} \), i.e. the verifier and falsifier of p, such that \([P]^+ = 1\) and \([p]^- = 0\). Thus, \( M = \langle S, S*, D, \leq, \leq^*, V, V^* \rangle \).

The two-dimensional hyperintensional profile of propositions may then be recorded by defining the value of p relative to two parameters, c,i: c ranges over subsets of S, and i ranges over subsets of S*.

\[ (*) \text{ M,} s \in S, s^* \in S^* \vdash p \text{ iff:} \]

\( (i) \exists c_s \llbracket p \rrbracket^c = 1 \text{ if } s \in \llbracket p \rrbracket^+; \) and

\( (ii) \exists i_s \llbracket p \rrbracket^c,i = 1 \text{ if } s^* \in \llbracket p \rrbracket^+ \)

(Distinct states, s,s*, from distinct state spaces, S,S*, provide a multi-dimensional verification for a proposition, p, if the value of p is provided a truthmaker by s. The value of p as verified by s determines the value of p as verified by s*).

We say that p is hyper-rigid iff:

\[ (*) \text{ M,} s \in S, s^* \in S^* \vdash p \text{ iff:} \]

\( (i) \forall c_s' \llbracket p \rrbracket'^c = 1 \text{ if } s \in \llbracket p \rrbracket^+; \) and

\( (ii) \forall i_s' \llbracket p \rrbracket'^c,i = 1 \text{ if } s^* \in \llbracket p \rrbracket^+ \)

The foregoing provides a two-dimensional hyperintensional semantic framework within which to interpret the values of a proposition. In order to account for partial contents, we define the values of subpropositional entities relative again to tuples of states from the distinct state spaces in our model:

s is a two-dimensional exact truthmaker of p if and only if (*)
s is a two-dimensional inexact truthmaker of p if and only if \( \exists s' \sqsubseteq S, s \rightarrow s' \), s' \( \vdash p \) and such that

\( \exists c_{s'} \llbracket p \rrbracket^c = 1 \text{ if } s' \in \llbracket p \rrbracket^+; \) and

\( \exists i_{s'} \llbracket p \rrbracket^c,i = 1 \text{ if } s^* \in \llbracket p \rrbracket^+; \)

s is a two-dimensional loose truthmaker of p if and only if, \( \exists t, s \sqcup t, s \sqcup t \vdash p; \)

\( \exists c_{s \sqcup t} \llbracket p \rrbracket^c = 1 \text{ if } s' \in \llbracket p \rrbracket^+; \) and

\( \exists i_{s \sqcup t} \llbracket p \rrbracket^c,i = 1 \text{ if } s^* \in \llbracket p \rrbracket^+; \)

Epistemic (primary), subjunctive (secondary), and 2D hyperintensions can be defined as follows, where hyperintensions are functions from states to extensions, and intensions are functions from worlds to extensions:
• Epistemic Hyperintension:
\[\text{pri}(x) = \lambda s.\llbracket x \rrbracket^s \mathpzc{s},\] with \(s\) a state in the state space defined over the foregoing epistemic modal algebra, \(U\);

• Subjunctive Hyperintension:
\[\text{sec}_{\mathpzc{v} \mathpzc{a}}(x) = \lambda i.\llbracket x \rrbracket^{\mathpzc{v} \mathpzc{a}, i},\] with \(i\) a state in metaphysical state space \(I\);

• 2D-Hyperintension:
\[2\mathpzc{D}(x) = \lambda s \lambda w \llbracket x \rrbracket^{s, i} = 1.\]
Part III: Epistemic Modality and the Philosophy of Mathematics
Chapter 8
Epistemic Modality, Necessitism, and Abstractionism

Modal notions have been availed of, in order to argue in favor of nominalist approaches to mathematical ontology. Field (1989) argues, for example, that mathematical modality can be treated as a logical consistency operator on a set of formulas comprising an empirical theory, such as Newtonian mechanics, in which the mathematical vocabulary has been translated into the vocabulary of physical geometry.\(^1\) Putnam (1967a), Parsons (1983), Chihara (1990), and Hellman (1993) argue that intensional models both of first- and second-order arithmetic and of set theory motivate an eliminativist approach to mathematical ontology. On this approach, reference to mathematical objects can be eschewed, and possibly the mathematical structures at issue are nothing.\(^2\)

This essay aims to provide modal foundations for mathematical platonism, i.e., the proposal that mathematical terms for sets; functions; and the natural, rational, real, and imaginary numbers refer to abstract – necessarily non-concrete – objects. Intensional constructions of arithmetic and set theory have been developed by, inter alia, Fine (1981); Parsons (op. cit.); Shapiro (1985); Myhill (1985); Reinhardt (1988); Nolan (2002); Linnebo

\(^1\)For a generalization of Field’s nominalist translation scheme to the differential equations in the theory of General Relativity, see Arntzenius and Dorr (2012).

\(^2\)For further discussion of modal approaches to nominalism, see Burgess and Rosen (1997: II, B-C) and Leng (2007; 2010: 258).
(2013); and Studd (2013). Williamson (2013a) emphasizes that mathematical languages are extensional, although in Williamson (2016) he argues that Orey sentences, such as the generalized continuum hypothesis \(2^{\aleph_\alpha} = \aleph_{\alpha+1}\) which are currently undecidable relative to the axioms of the language of Zermelo-Fraenkel Set Theory with choice as augmented by large cardinal axioms, are yet possibly decidable.\(^3\) This chapter and chapter 10 argue that the epistemic interpretation of two-dimensional semantics provides a novel approach to the epistemology of mathematics, such that if the decidability of mathematical axioms is epistemically possible, then their decidability is metaphysically possible.\(^4\) Epistemic mathematical modality, suitably constrained, can thus serve as a guide to metaphysical mathematical modality.\(^5\) Hamkins and Löwe (2007; 2013) argue that the modal logic of set-forcing extensions and the corresponding logic for their ground models satisfy at least S4.2, i.e., axioms K [\(\square(\phi \rightarrow \psi) \rightarrow (\square\phi \rightarrow \square\psi)\)]; T (\(\square\phi \rightarrow \phi\)); 4 (\(\square\phi \rightarrow \square\square\phi\)); and G (\(\Diamond\square\phi \rightarrow \square\Diamond\phi\)). While the foregoing approaches are consistent with realism about mathematical objects, they are nevertheless not direct arguments thereof. The aim of this essay is to redress the foregoing lacuna, and thus to avail of the resources of modal ontology and epistemology in order to argue for the reality of mathematical entities and truth.

In Section 2, I outline the elements of the abstractionist foundations of mathematics. In Section 3, I examine Hale and Wright (2009)’s objections to the merits and need, in the defense of mathematical platonism and its epistemology, of the thesis of Necessitism, underlying the thought that whatever can exist actually does so. The Necessitist thesis is codified by the Barcan formula (cf. Barcan, 1946; 1947), and states that possibly if there is something which satisfies a condition, then there is something such that it possibly satisfies that condition: \(\Diamond\exists x \phi x \rightarrow \exists x \Diamond\phi x\). I argue that Hale and Wright’s objections to Necessitism as a requirement on admissible abstraction can be answered; and I examine both the role of the higher-order Necessitist pro-

\(^3\)Compare Reinhardt (1974) on the imaginative exercises taking the form of counterfactuals concerning the truth of undecidable formulas. See Maddy (1988b), for critical discussion.

\(^4\)The epistemic interpretation of two-dimensional intensional semantics is first advanced in Chalmers (1996; 2004).

\(^5\)See Section 4, for further discussion. Gödel (1951: 11-12) anticipates a similar distinction between epistemic and metaphysical readings of the determinacy of mathematical truths, by distinguishing between mathematics in its subjective and objective senses. The former targets all "demonstrable mathematical propositions", and the latter includes "all true mathematical propositions".
posal in their endorsement of an abundant conception of properties, as well as cardinality issues that arise depending on whether Necessitism is accepted at first- and higher-order. In Section 4, I provide an account of the role of epistemic and metaphysical modality in explaining the prima facie justification to believe the truth of admissible abstraction principles, and demonstrate how it converges with both Hale and Wright’s (op. cit.) and Wright’s (2012; 2014) preferred theory of default entitlement rationally to trust the truth of admissible abstraction. Section 5 provides concluding remarks.

8.1 The Abstractionist Foundations of Mathematics

The abstractionist foundations of mathematics are inspired by Frege’s (1884/1980; 1893/2013) proposal that cardinal numbers can be explained by specifying an equivalence relation, expressible in the signature of second-order logic and identity, on first- or higher-order entities. Thus, e.g., in Frege (1884/1980: 64), the direction of the line, a, is identical to the direction of the line, b, if and only if lines a and b are parallel. In Frege (op. cit.: 68) and Wright (1983: 104-105), the cardinal number of the concept, A, is identical to the cardinal number of the concept, B, if and only if there is a one-to-one correspondence between A and B, i.e., there is an injective and surjective (bijective) mapping, R, from A to B. With Nx: a numerical term-forming operator,

\[ \forall A \forall B \exists R [N x : A = N x : B \equiv \exists R \forall x [A x \rightarrow \exists y (B y \land R x y \land \forall z (B z \land R x z \rightarrow y = z))] \land \forall y [B y \rightarrow \exists x (A x \land R x y \land \forall z (A z \land R z y \rightarrow x = z))]]. \]

The foregoing is referred to as ‘Hume’s Principle’. Frege’s Theorem states

\[ \forall A \forall B \exists R [N x : A = N x : B \equiv \exists R \forall x [A x \rightarrow \exists y (B y \land R x y \land \forall z (B z \land R x z \rightarrow y = z))] \land \forall y [B y \rightarrow \exists x (A x \land R x y \land \forall z (A z \land R z y \rightarrow x = z))]]. \]

---

6 Frege (1884/1980: 68) writes: ‘the Number which belongs to the concept F is the extension of the concept “equinumerous” to the concept F’ (cf. op. cit.: 72-73). Boolos (1987/1998: 186) coins the name, ‘Hume’s Principle’, for Frege’s abstraction principle for cardinals, because Frege (op. cit.: 63) attributes equinumerosity as a condition on the concept of number to Hume (1739-1740/2007: Book 1, Part 3, Sec. 1, SB71), who writes: ‘When two numbers are so combin’d, as that the one has always an unite answering to every unite of the other, we pronounce them equal . . . ’. Frege notes that identity of number via bijections is anticipated by the mathematicians, Ernst Schröder and Ernst Kossak, as well Cantor (1883/1996: Sec. 1), who writes: ‘[E]very well-defined set has a determinate power; two sets have the same power if they can be, element for element, correlated with one another reciprocally and one-to-one’, where the power [Anzahl] of a set corresponds to its cardinality (cf. Cantor, 1895/2007: 481).
that the Dedekind-Peano axioms for the language of arithmetic can be derived from Hume’s Principle, as augmented to the signature of second-order logic and identity.\textsuperscript{7} Abstraction principles have further been specified both for the real numbers (cf. Hale, 2000a; Shapiro, 2000; and Wright, 2000), and for sets (cf. Wright, 1997; Shapiro and Weir, 1999; Hale, 2000b; and Walsh, 2016).

The philosophical significance of the abstractionist program consists primarily in its provision of a neo-logicist foundation for classical mathematics, and in its further providing a setting in which to examine constraints on the identity conditions constitutive of mathematical concept possession.\textsuperscript{8} The philosophical significance of the abstractionist program consists, furthermore, in its circumvention of Benacerraf’s (1973) challenge to the effect that our knowledge of mathematical truths is in potential jeopardy, because of the absence of naturalistic, in particular causal, conditions thereon. Both Wright (1983: 13-15) and Hale (1987: 10-15) argue that the abstraction principles are epistemically tractable, only if (i) the surface syntax of the principles – e.g., the term-forming operators referring to objects – are a perspicuous guide to their logical form; and (ii) the principles satisfy Frege’s (1884/1980: X) context principle, such that the truth of the principles is secured prior to the reference of the terms figuring therein.

\textbf{8.2 Abstraction and Necessitism}

\textbf{8.2.1 Hale and Wright’s Arguments against Necessitism}

One crucial objection to the abstractionist program is that – while abstraction principles might provide a necessary and sufficient truth-condition for our grasp of the concepts of mathematical objects – an explanation of the actual truth of the principles has yet to be advanced (cf. Eklund, 2006; 2016). In response, Hale and Wright (2009: 197-198) proffer a tentative endorsement

\textsuperscript{7}Cf. Dedekind (1888/1996) and Peano (1889/1967). See Wright (1983: 154-169) for a proof sketch of Frege’s theorem; Boolos (1987) for the formal proof thereof; and Parsons (1964) for an incipient conjecture of the theorem’s validity.

\textsuperscript{8}Shapiro and Linnebo (2015) prove that Heyting arithmetic can be recovered from Frege’s Theorem. Criteria for consistent abstraction principles are examined in, inter alia, Hodes (1984a); Hazen (1985); Boolos (1990/1998); Heck (1992); Fine (2002); Weir (2003); Cook and Ebert (2005); Linnebo and Uzquiano (2009); Linnebo (2010); and Walsh (op. cit.).
of an ‘abundant’ conception of properties, according to which fixing the sense of a predicate will be sufficient for predicate reference. Eklund (2006: 102) suggests, by contrast, that one way for the truth of the abstraction principles to be explained is by presupposing what he refers to as a ‘Maximalist’ position concerning the target ontology. According to the ontological Maximalist position, if it is possible that a term has a certain extension, then actually the term does have the designated extension.

Hale and Wright (op. cit.) raise two issues for the ontological Maximalist proposal. The first is that ontological Maximalism is committed to a proposal that they take to be independently objectionable, namely ontological Necessitism (185). They write: "Most obviously, maximalism denies the possibility of contingent non-existence, to which there are obvious objections" (op. cit.) Hale and Wright (op. cit.) raise a similar contention to the effect that actual, and not merely possible, reference is what the abstractionist program intends to target; and that Maximalism and Necessitism, so construed, are purportedly silent on the status of ascertaining when the possibilities at issue are actual.

The second issue that Hale and Wright find with Maximalism is that it misconstrues the demands that the abstractionist program is required to address. The abstractionist program is supposed to be committed to ontological Maximalism, because the possibility that a term has a certain extension will otherwise not be sufficient for the success of the term’s reference. It is further thought that, without an appeal to Maximalism, and despite the actuality of successful mathematical predicate reference, there are yet possible situations in which the mathematical predicates still do not refer (193). In response, they note that no ‘collateral metaphysical assistance’ – such as ontological Maximalism would be intended to provide – is necessary in order to explain the truth of abstraction principles (op. cit.). Rather, there is prima facie, default entitlement rationally to trust that the abstraction principles are actually true, and such entitlement is sufficient to foreclose upon the risk that possibly the mathematical terms therein do not refer (192).

In the remainder of this section, I will argue that Hale and Wright’s objections to Necessitism and the ontological Maximalist approach to admissible abstraction both can be answered, and in any case are implicit in their endorsement of the abundant conception of properties. In the following section,

\footnote{For further discussion of ontological Maximalism, see Hawley (2007) and Sider (2007: IV).}
I address their second contention, and I argue for the fundamental role that Maximalism and Necessitism can play in warranting the truth of candidate abstractions and mathematical platonism.

The principle of the necessary necessity of being (NNE) can be derived from the Barcan formula.\(^\text{10}\) NNE states that necessarily everything is necessarily such that there is something to which it is identical; \(\square \forall x \exists y (x = y)\). Informally, necessarily everything has necessary being, i.e. everything is something, even if contingently non-concrete. Williamson (2013a: 6.1-6.4) targets issues for haecceity comprehension, if the negations of the Barcan formula and NNE are true at first-order, and thus for objects. With regard to properties and relations at higher-order, Williamson’s arguments have targeted closure conditions, given a modalized interpretation of comprehension principles (op. cit.). The latter take the form, \(\exists X \square \forall x (X x \iff \phi)\), with \(x\) an individual variable which may occur free in \(\phi\) and \(X\) a monadic first-order predicate variable which does not occur free in \(\phi\) (262).\(^\text{11}\) He targets, in particular, the principle of mathematical induction – with \(s\) a successor function and the quantifier ranging over the natural numbers: \(\forall X [[X 0 \land \forall n (X n \rightarrow X sn)] \rightarrow \forall n (X n)]\) – and notes that instances of mathematical induction – e.g., for \(a\) an individual constant, \(\exists R [[Ra 0 \land \forall n (Ran \rightarrow Rasn)] \rightarrow \forall n (Ran)]\) – presuppose, for their derivation, the validity of instances of the higher-order modal comprehension scheme: e.g., \(\exists X \square \forall n (X n \iff \text{Ran})\) (283-284). The foregoing provides prima facie abductive support for the requirement of Necessitism in the practice of mathematics. The constitutive role of the Necessitist modal comprehension scheme in the principle of mathematical induction answers Hale and Wright’s first contention against the Necessitist commitments of ontological Maximalism.

Williamson refers to the assignments for models in the metaphysical setting as universal interpretations (59). The analogue for logical truth occurs when a truth is metaphysically universal, i.e., if and only if its second-order universal generalization is true on the intended interpretation of the metalanguage (200). The connection between truth-in-a-model and truth simpliciter is then that – as Williamson puts it laconically – when ‘the framework at least delivers a condition for a modal sentence to be true in a universal interpretation, we can derive the condition for it to be true in the intended

\(^{10}\text{Cf. Williamson (2013a: 38).}\)

\(^{11}\text{The contingentist, by contrast, can – by rejecting the Barcan formula – countenance only ‘intra-world’ comprehension principles in which the modal operators and iterations thereof take scope over the entire formula; e.g. }\diamond \forall x \exists X (X x \iff \phi)\text{ (cf. Sider, 2016: 686).}\)
universal interpretation, which is the condition for it to be true simpliciter’ (op. cit.).

One of the crucial interests of the metaphysical universality of formulas is that the models in the class need not be pointed, in order to countenance the actuality of the possible formulas defined therein. Rather, the class of true propositions generated by the metaphysically universal formulas is sufficient for the formulas actually to be true (268-269). Williamson writes that 'since whatever is is, whatever is actually is: if there is something, then there actually is such a thing’ (23). Thus, the foregoing definition of actuality can explain why the metaphysically universal formulas which are true simpliciter are actual.

The constitutive role of metaphysical universality in bridging the necessary necessity of being with the actuality thereof answers Hale and Wright’s contention that the interaction between the possible and actual truth of abstraction principles has yet to be accounted for.

8.2.2 Hale on the Necessary Being of Purely General Properties and Objects

Note, further, that the abundant conception of properties endorsed by Hale and Wright depends upon the Necessitist Thesis, and the truth of ontological Maximalism thereby. Hale writes: ‘[I]t is sufficient for the actual existence of a property or relation that there could be a predicate with appropriate satisfaction conditions ... purely general properties and relations exist as a matter of (absolute) necessity’, where a property is purely general if and only if there is a predicate for which, and it embeds no singular terms (Hale, 2013b: 133, 135; see also 2013a: 99-100).  

\[\text{\textsuperscript{12}} \text{That the models are unpointed is noted in Williamson (2013: 100).}\]

\[\text{\textsuperscript{13}} \text{Thanks here to Bruno Jacinto for discussion.}\]

\[\text{\textsuperscript{14}} \text{Cook (2016: 398) demonstrates how formally to define modal operators within Hume’s Principle, i.e. the consistent abstraction principle for cardinal numbers. Necessitist Hume’s Principle takes the form: } \square X, Y \#(X) = \square Y \iff X \approx Y, \text{ where } X \text{ and } Y \text{ are second-order variables, } \# \text{ is a numerical term-forming operator, } \approx \text{ is a bijection, and for variables, } x, y, \text{ of arbitrary type ' } x = \square y \iff \exists z [z = x \land z = y \land \square \exists w (w = z)][. \text{ See Cook (op. cit.) for further discussion.}\]

\[\text{\textsuperscript{15}} \text{Cook (op. cit.: 388) notes the requirement of Necessitism in the abundant conception of properties, although does not discuss points at which Williamson’s and Hale’s Necessitist proposals might be inconsistent. The points of divergence between the two variations on the proposal are examined below.}\]
Hale argues for the necessary necessity of being for properties and propositions as follows (op. cit.: 135; 2013b: 167). Suppose that p refers to the proposition that a property exists, and that q refers to the proposition that a predicate for the property exists. Let the necessity operator be defined as a counterfactual with an unrestricted, universally quantified antecedent, such that, for all propositions, \( \psi: [\Box \psi \iff \forall \phi (\phi \rightarrow \psi)] \) (135). On the abundant conception of properties, \( \Box [p \iff \Diamond q] \). Intuitively: Necessarily, there is a property if and only if possibly there is a predicate for that property. Given the counterfactual analysis of the modal operator: For all propositions about a property, if there were a proposition specifying a predicate s.t. the property is in the predicate’s extension, then there would be that property.

From ‘\( \Box [p \iff \Diamond q] \)’, one can derive both ‘\( p \iff \Diamond q \)’, and – by the rule, RK – the necessitation thereof, ‘\( \Box p \iff \Box \Diamond q \)’ (op. cit.). By the B axiom in S5, \( \Diamond q \iff \Box \Diamond q \) (op. cit.). So, ‘\( \Box \Diamond q \iff \Diamond q \)’; ‘\( \Diamond q \iff p \)’; and ‘\( \Box \Diamond q \iff \Box p \)’. Thus – by transitivity – ‘\( p \iff \Box p \)’ (op. cit.); i.e., all propositions about properties are necessarily true, such that the corresponding properties have necessary being. By the 4 axiom in S5, \( \Box p \iff \Box \Box p \); so, the necessary being of properties and propositions is itself necessary. Given the endorsement of the abundant conception of properties – Hale and Wright are thus committed to higher-order necessitism, i.e., the necessary necessity of being.

Hale (2013a) endeavors to block the ontological commitments of the Barcan formula and its converse by endorsing a negative free logic. Thus, in the derivation:

Assumption,
1. \( \Box \forall x [F(x)] \).
   By \( \Box \)-elimination,
2. \( \forall x [F(x)] \).
   By \( \forall \)-elimination,
3. \( F(x) \).
   By \( \Box \)-introduction,
4. \( \Box [F(x)] \).
   By \( \forall \)-introduction,
5. \( \forall \Box [F(x)] \).
   By \( \rightarrow \)-introduction,

---

16 Proponents of the translation from modal operators into counterfactual form include Stalnaker (1968/1975), McFetridge (1990: 138), and Williamson (2007).
6. $\Box \forall x[F(x) \rightarrow \forall \Box[F(x)]]$,

Hale imposes an existence-entailing assumption in the inference from lines (2) to (3), i.e.

'(Free∀-Elimination) From $\forall x[A(x)]$, together with an existence-entailing premise $F(t)$, we may infer $A(t)$ where $t$ can be any term' (op. cit.: 208-209).

Because the concept of, e.g., cardinal number is defined by abstraction principles which are purely general because they embed no singular terms, the properties – e.g., the concepts – of numbers are argued to have necessary being. The necessary being of the essential properties of number – i.e., higher-order Necessitism about purely general properties – is argued then to explain in virtue of what abstract objects such as numbers and functions have themselves necessary being (176-177). Thus the necessary being of predicate sense for the concept of number can both suffice for and explain the necessary being of predicate reference, i.e. the necessary existence of numbers.

By contrast, essential properties defined by theoretical identity statements, which if true are necessarily so, do embed singular terms and are thus not purely general. So, the essential nature of water, i.e., the property ‘being comprised of one oxygen and two hydrogen molecules’, has contingent being, explaining in virtue of what samples of water have contingent being (216-217).

**Objections**

One objection to the foregoing is that Hale takes the modal status of the being with which both logical and non-logical properties exist to be equivalent. Following Fine (2005), he notes that there is a distinction between ‘unworldly’ or ‘transcendental’ truths about individuals, which are true, in Fine’s phrasing, ‘on the basis of [their] logical form alone and without regard to the circumstances’ (Fine, op. cit.: 324), and ‘worldly’ or necessary truths. An example of a transcendental truth is that ‘Hypatia is self-identical’, i.e., $\Box \exists x(x = H \land \Box x = x)$. A ‘worldly’ or necessary truth is, by contrast, one whose truth-value is defined in a world; e.g., that ‘water = H20’ or – if one were to augment one’s language with an existence predicate beyond the quantifiers – that ‘Socrates exists or does not exist’ (op. cit.).

However, Hale draws, as noted, no similar distinction between the modal status of a sentence true in virtue of its logical form, and a worldly sentence whose truth depends on non-logical values of its constituent variables, e.g., a truth of physics (Hale, op. cit.: 215). Hale (2000/2001: 415) notes Wright’s
argument that the conjunction of two predicates, e.g., being blue and being self-identical, is equivalent to one of the conjuncts, e.g., being blue. Thus, the predicate for the property, being self-identical, cannot be purely general. He argues, thus, that the truth of ‘this star is self-identical’ depends on the concrete existence of the star, such that the logical property, being self-identical, has contingent being. Thus, the status of the being of logical properties is contingent, in the same manner that the essential property, being, e.g., H20, depends on concrete instances of water.

A problematic consequence of the foregoing is that the logical necessity of a formula will thus depend on whether the predicates therein are purely general by embedding no singular terms, rather than on whether the formula at issue is a logical truth. It might be replied that Hale is following Frege in defining one of the constitutive marks of logical truths as consisting in their generality – e.g., the generality of their application (1893/2013: XV; 1897/1997), as well as whether the formula is a true universal generalization (op. cit.: §8-9) – rather than Tarski’s (1936/1983: 415-417) definition of a logical truth as a formula true in virtue of its logical form and thus whose truth is invariant under permutation of the values of the variables which replace the non-logical constants therein. However – even if not a purely general property because it embeds singular terms – the reflexivity of identity is a logical law, because – as Frege himself writes of reflexivity – ‘the value of this function is always the True, whatever we take as argument’ (Frege, 1891/1997: 23).17

A second objection concerns the necessary being of different types of numbers. While an abstraction principle for cardinal numbers can be specified using only purely general predicates – i.e., Hume’s Principle – abstraction principles for imaginary and complex numbers have yet to be specified. Shapiro (2000) provides an abstraction principle for the concepts of the reals by simulating Dedekind cuts, where abstraction principles are provided for the concepts of the cardinals, natural numbers, integers, and rational numbers, from which the reals are thence defined: Letting \( F, G, \) and \( R \) denote rational numbers, \( \forall F, G [\mathcal{C}(F) = \mathcal{C}(G) \iff \forall R (F \leq R \iff G \leq R)] \).

Hale’s (2000/2001) own definition of the concept of the reals is provided relative to a domain of quantities. The quantities are themselves taken to be abstract, rather than physical, entities (409). The quantitative domain can thus be comprised of both rational numbers as well as the abstracts for lengths, masses, and points. The reals are then argued not to be numbers, but rather quantities defined via an abstraction principle which states that a set of rational numbers in one quantitative domain is identical to a set of rational numbers in a second quantitative domain if and only if the two domains are isomorphic (407). Hale argues, then, that it is innocuous for the real abstraction principle to be conditional on the existence of at least one quantitative domain, because the rational numbers can be defined, similarly as on Shapiro’s approach, via cut-abstractions and abstractions on the integers, naturals, and cardinals. Thus, the reals can be treated as abstracts derived from purely general abstraction principles, and are thus possessed of necessary being. However, abstraction principles for imaginary numbers such as \( i = \sqrt{-1} \), and complex numbers which are defined as the sum of a real number and a second real multiplied by \( i \), have yet to be accounted for. The provision of an abstraction principle for complex numbers would, in any case, leave open the inquiry into how, e.g., complex-valued wave functions might interact with physical ontology; e.g., whether such functions might be metaphysically fundamental entities which serve to represent physical fields in higher-dimensional spacetime, and whether or how the domain of the functions, i.e., a real-valued configuration space for particles, might relate to the higher-dimensional, complex-valued wave function (cf. Simons, 2016; Ney, 2013; Maudlin, 2013).

The modality in the Barcan-induced Necessitist proposal at first- and higher-order is, as noted, interpreted metaphysically rather than logically, and thus incurs no similar issues with regard to the interaction between purely general properties, logical properties, and concrete entities. Further, because true on its second-order universal generalization on its intended, metaphysical interpretation, the possible truth-in-a-model of the relevant

---

18 See Dedekind (1872/1996: Sec. 4), for the cut method for the definition of the reals.  
19 An abstraction principle for lengths, based on the equivalence property of congruence relations on intervals of a line, or regions of a space, is defined in Shapiro and Hellman (2015: 5, 9). Shapiro and Hellman provide, further, an abstraction principle for points, defined as comprising, respectively, the left- and right-ends of intervals (op. cit.: 5, 10-12).  
20 Cf. Hale (op. cit.: 406-407), for the further conditions that the domains are required to satisfy.
class of formulas is, as discussed in Section 3.1, thus sufficient for entraining the actual truth of the relevant formulas.

### 8.2.3 Cardinality and Intensionality

An interesting residual question concerns the status of the worlds, upon the translation of modal first-order logic into the non-modal first-order language.\(^{21}\) Fritz (op. cit.) notes that a world can be represented by a predicate, in the latter.\(^{22}\) However, whether objects satisfy the predicate can vary from point to point, in the non-modal first-order class of points.\(^{23}\) Another issue is that modal propositional logic is equivalent only to the bisimulation-invariant fragments of both first-order logic and fixed-point monadic second-order logic, rather than to the full variants of either logic (cf. van Benthem, 1983; Janin and Walukiewicz, 1996). Thus, there cannot be a faithful translation from each modal operator in modal propositional logic into a predicate of full first- or monadic higher-order logic.

One way to mitigate the foregoing issues might be by arguing that the language satisfies real-world rather than general validity, such that necessarily the predicate will be satisfied only at a designated point in a model – intuitively, the analogue of the concrete rather than some merely possible world, simulating thereby the translation from possibilist to actualist discourse (cf. Fine, op. cit.: 211,135-136, 139-140, 154, 166-168, 170-171) – by contrast to holding of necessity as interpreted as satisfaction at every point in the model. The reply would be consistent with what Williamson refers to as ‘chunky-style necessitism’ which validates the following theorems: where the predicate \(C(x)\) denotes the property of being grounded in the concrete and \(P(x)\) is an arbitrary predicate, (a) ‘\(\forall x \diamond C(x)\)’, yet (b) ‘\(\Box \forall x [P(x_1, \ldots, x_n) \rightarrow (Cx_1, \ldots, Cx_1)]\)’ (325-332). Williamson (33, fn.5) argues, however, in favor of general, rather than real-world validity. A second issue for the reply is that principle (b), in the foregoing, is inconsistent with Williamson’s protracted defense of the ‘being constraint’, according to which \(\Box \forall x \Box [P(x_1, \ldots, x_n) \rightarrow \exists y (x = y)]\), i.e. if \(x\) satisfies a predicate, then \(x\) is something.

\(^{21}\) Thanks here to Alessandro Rossi, for discussion.

\(^{22}\) For further discussion of the standard translation between propositional modal and first-order non-modal logics, see Blackburn et al. (2001: 84).

\(^{23}\) Suppose that the model is defined over the language of second-order arithmetic, such that the points in the model are the ordinals. A uniquely designated point might then be a cardinal number whose height is accordingly indexed by the ordinals.

123
even if possibly non-concrete (148).

A related issue concerns the translation of modalized, variable-binding, generalized quantifiers of the form:

‘there are \( n \) objects such that \( \ldots \)’,

‘there are countably infinite objects such that \( \ldots \)’,

‘there are uncountably infinite objects such that \( \ldots \)’ (Fritz and Goodman, 2017).

The generalized quantifiers at issue are modalized and consistent with first-order Necessitism, because the quantifier domains include all possible – including contingently non-concrete – objects. It might be argued that the translation is not of immediate pertinence to the ontology of mathematics, because the foregoing first-order quantifiers can be restricted such that they range over only uncountably infinite necessarily non-concrete objects – i.e. abstracta – by contrast to ranging unrestricedly over all modal objects, including the contingently non-concrete entities induced via the Barcan formula – i.e., the ‘mere possibilia’ that are non-concrete as a matter of contingency.

However, the Necessitist thesis can be valid even in the quantifier domain of a first-order language restricted to necessarily non-concrete entities. If, e.g., a mathematician takes, despite iterated applications of set-forming operations, the cumulative hierarchy of sets to have a fixed cardinal height, then the first-order Necessitist thesis will still be valid, because all possible objects will actually be still something.

The first-order Necessitist proposal engendered by taking the height of the cumulative hierarchy to be fixed is further consistent with the addition to the first-order language of additional intensional operators – such as those introduced by Hodes (1984b) – in order to characterize the indefinite extensibility of the concept of set; i.e., that despite unrestricted universal quantification over all of the entities in a domain, another entity can be defined with reference to, and yet beyond the scope of, that totality, over which the quantifier would have further to range.24 First-order Necessitism is further consistent with the relatively expanding domains induced by Bernays’ (1942) Theorem. Bernays’ Theorem states that class-valued functions from classes to sub-classes are not onto, where classes are non-sets (cf. Uzquiano, 2015b: 186-187). So, the cardinality of a class will always be less than the cardinal-

---

24The concept of indefinite extensibility is introduced by Dummett (1963/1978), in the setting of a discussion of the philosophical significance of Gödel’s (1931) first incompleteness theorem. See the essays in Rayo and Uzquiano (2006); Studd (op. cit.); and Dever (ms) for further discussion.
ity of its sub-classes. Suppose that there is a generalization of Bernays’ theorem, such that the non-sets are interpreted as possible objects. Thus, the cardinality of the class of possible objects will always be less than the cardinality of the sub-classes in the image of its mapping. Given iterated applications of Bernays’ theorem, the cardinality of a domain of non-sets is purported then not to have a fixed height.

In both cases, however, the addition of Hodes’ intensional operators permits there to be multiple-indexing in the array of parameters relative to which a cardinal can be defined, while the underlying logic for metaphysical modality can be S5, partitioning the space of worlds into equivalence classes. So, both the intensional characterization of indefinite extensibility and the generalization of Bernays’ Theorem to possible objects are consistent with the first-order Necessitist proposal that all possible objects are actual, and so the cardinality of the target universe is fixed.

Fritz and Goodman suggest that a necessary condition on the equivalence of propositions is that they define the same class of models (op. cit.: 1.4). The proposed translation of the modalized generalized quantifiers would be Contingentist, by taking (NNE) to be invalid, such that the domain in the translated model would be comprised of only possible concrete objects, rather than the non-concrete objects as well (op. cit.).

Because of the existence of non-standard models, the generalized quantifier that ‘there are countably infinitely many possible …’ cannot be defined in first-order logic. Fritz and Goodman note that generalized quantifiers ranging over countably infinite objects can yet be simulated by enriching one’s first-order language with countably infinite conjunctions. On the latter approach, finitary existential and universal quantifiers can be defined as the countably infinite conjunction of formulas stating that, for all natural numbers \( n \), ‘there are \( n \) possible ...’ (2.3).

Crucially, however, there are some modalized generalized quantifiers that

25 Note that the proposal that the cardinality of the cumulative hierarchy of sets is fixed, despite continued iterated applications of set-forming operations, is anticipated by Cantor (1883/1996: Endnote [1]). Cantor writes: ‘I have no doubt that, as we pursue this path ever further, we shall never reach a boundary that cannot be crossed, but that we shall also never achieve even an approximate conception of the absolute [...]. The absolutely infinite sequence of numbers thus seems to me to be an appropriate symbol of the absolute; in contrast the infinity of the first number-class (I) [i.e., the first uncountable cardinal, \( \aleph_0 \) – HK], which has hitherto sufficed, because I can consider it to be a graspable idea (not a representation), seems to me to dwindle into nothingness by comparison’ (op. cit.; cf. Cantor, 1899/1967).
cannot be similarly paraphrased – e.g., ‘there are uncountably infinite possible objects s.t. . . .’ – and there are some modalized generalized quantifiers that cannot even be defined in first-order languages – e.g. ‘most possible objects s.t. . . .’ (2.4-2.5)

In non-modal first-order logic, it is possible to define generalized quantifiers which range over an uncountably infinite domain of objects, by augmenting finitary existential and universal quantifiers with an uncountably infinite stock of variables and an uncountably infinite stock of conjunctions of formulas (2.4). Fritz and Goodman note, however, that the foregoing would require that the quantifiers bind the uncountable variables ‘at once’, s.t. they must have the same scope. The issue with the proposal is that, in the setting of modalized existential quantification over an uncountably infinite domain, the Contingentist paraphrase requires that bound variables take different scopes, in order to countenance the different possible sets that can be defined in virtue of the indefinite extensibility of cardinal number (op. cit.).

In order to induce the Contingentist paraphrase, Fritz and Goodman suggest defining ‘strings of infinitely many existential and universal quantifiers’, such that a modalized, i.e. Necessitist, generalized quantifier of the form, ‘there are uncountably infinite possible . . .’ can be redefined by an uncountably infinite sequence of finitary quantifiers with infinite variables and conjunction symbols of the form:

‘Possibly for some $x_1$, possibly for some $x_2$, etc.: $x_1, x_2, etc.$ are pairwise distinct and are each possibly . . .’,

where etc. denotes an uncountable sequence of, respectively, ‘an uncount-

---

26Uncountable cardinals can be defined as follows. For cardinals, $\kappa, \lambda, \theta, \alpha$ let $\kappa \subseteq a$ be closed unbounded in $a$, if it is closed [if $\kappa < \lambda$ and $\bigcup (\lambda \cap \kappa) = \lambda$, then $\alpha \in \kappa$] and unbounded ($\bigcup \kappa = a$) (Kanamori, 2012b: 360). A cardinal, $\kappa$, is stationary in $a$, if, for any closed unbounded $\kappa \subseteq a$, $\kappa \cap S \neq \emptyset$ (op. cit.). An ideal is a subset of a set closed under countable unions, whereas filters are subsets closed under countable intersections. A cardinal $\kappa$ is regular if the cofinality of $\kappa$ – comprised of the unions of sets with cardinality less than $\kappa$ – is identical to $\kappa$. For models $A, B$, and conditions $\phi$, an elementary embedding, $j$: $A \rightarrow B$, is such that $\phi(a_1, \ldots, a_n)$ in $A$ if and only if $\phi(j(a_1), \ldots, j(a_n))$ in $B$ (363). A measurable cardinal is defined as the ordinal denoted by the critical point of $j$, crit($j$) (Koellner and Woodin, 2010: 7). Measurable cardinals are inaccessible (Kanamori, op. cit.). Uncountable regular limit cardinals are weakly inaccessible (op. cit.). A strongly inaccessible cardinal is regular and has a strong limit, such that if $\lambda < \kappa$, then $2^\lambda < \kappa$ (op. cit.). For the foregoing and further definitions, see Koellner and Woodin (op.cit.); Kanamori (op. cit.), and Woodin (2009, 2010).
able string of interwoven possibility operators and existential quantifiers’, and an ‘uncountable string of variables’ (op. cit.).

An argument against the proposed translation of the quantifier for there being uncountably infinite possible objects is that it is contentious whether an uncountable sequence of operators or quantifiers has a definite meaning [cf. Williamson (2013a: 7.7)]. Thus, e.g., while negation can have a determinate truth condition which specifies its meaning, a string of uncountably infinite negation operators will similarly have determinate truth conditions and yet not have an intuitive, definite meaning (357). One can also define a positive or negative integer, x, such that $s^x$ is interpreted as the successor function, x+1, and $p^x$ is interpreted as the inverse function, x-1. However, an infinitary expression consisting in uncountable, alternating iterations of the successor and inverse functions – $spspspspsp...^x$ – will similarly not have a definite meaning (op. cit.). Finally, one can define an operator $O^i$ mapping truth conditions for an arbitrary formula $A$ to the truth condition, $p^i$, of the formula $\Diamond \exists x_i (Cx_i \land A)$, with $Cx$ being the predicate for being concrete (258). Let the operators commute s.t. such that $O^iO^j$ iff $O^jO^i$, and be idempotent such that $O^iO^i$ iff $O^i$ (op. cit.). A total ordering of truth conditions defined by an infinite sequence of the operators can be defined, s.t. that the relation is reflexive, anti-symmetric, transitive, and connected $[\forall x,y(x \leq y \lor y \leq x)]$ (op. cit.). However, total orders need not have a least upper bound; and the sequence, $O^iO^iO^i...^i(p)$, would thus not have a non-arbitrary, unique value (op. cit.). The foregoing might sufficiently adduce against Fritz and Goodman’s Contingentist paraphrase of the uncountable infinitary modalized quantifier.

The philosophical significance of the barrier to a faithful translation from modal first-order to extensional full first-order languages, as well as a faithful translation from modalized, i.e. Necessitist, generalized quantifiers to Contingentist quantification, is arguably that the modal resources availed of in the abstractionist program might then be ineliminable.
8.3 Epistemic Modality, Metaphysical Modality, and Epistemic Utility and Entitlement

In this section, I address, finally, Hale and Wright’s second issue with the role of Necessitism in guaranteeing that the possible truth of abstraction principles provides warrant for the belief in their actual truth. As noted, Hale and Wright argue, against the foregoing approach, that there is non-evidential entitlement rationally to trust that acceptable abstraction principles are true, and thus that the terms defined therein actually refer. In response, I will proceed by targeting the explanation in virtue of which there is such epistemic, default entitlement. I will outline two proposals concerning the foregoing grounding claim – advanced, respectively, in chapter 10 and by Wright (2012b; 2014) – and I will argue that the approaches converge.

Wright’s elaboration of the notion of rational trust, which is intended to subserve epistemic entitlement, appeals to a notion of ‘expected epistemic utility’ in the setting of decision theory (2014: 226, 241). In order better to understand this notion of expected epistemic utility, we must be more precise.

There are two, major interpretations of (classical) expected utility.\(^{27}\) A model of decision theory is a tuple \(\langle A, O, K, V \rangle\), where \(A\) is a set of acts; \(O\) is a set of outcomes; \(K\) encodes a set of counterfactual conditionals, where an act from \(A\) figures in the antecedent of the conditional and \(O\) figures in the conditional’s consequent; and \(V\) is a function assigning a real number to each outcome. The real number is a representation of the value of the outcome. In evidential decision theory, the expected utility of an outcome is calculated as the product of the agent’s credence, conditional on her action, by the utility of the outcome. In causal decision theory, the expected utility of an outcome is calculated as the product of the agent’s credence, conditional on both her action and background knowledge of the causal efficacy thereof, by the utility

\(^{27}\) For an examination of non-classical utility measures, see Buchak (2014). Non-classical utility measures are intended to describe the innocuous rationality with which an agent’s expected utility might diminish with the order of the bets she might pursue. In the latter case, her expected utility will then be sensitive to her propensity to take risks relative to the total ordering of the gambles, such that she can have a preference for a sure-gain of .5 units of value, rather than prefer a bet with a 50 percent chance of winning either 0 or 1 units of value.
of the outcome.

First, because background knowledge concerning the causal efficacy of one’s choice of acts is presumably orthogonal to the non-evidential rational trust to believe that mathematical abstraction principles are true, I will assume that the notion of expected epistemic utility theory that Wright (op. cit.) avails of relies only on the subjective credence of the agent, multiplied by the utility that she assigns to the outcome of the proposition in which she’s placing her rational trust. Thus expected epistemic utility in the setting of decision-theory will be calculated within the (so-called) evidential, rather than causal, interpretation of the latter.

Second, there are two, major interpretations concerning how to measure the subjective credences of an agent. The philosophical significance of this choice point is that it bears directly on the very notion of the epistemic utility that an agent’s beliefs will possess. So, e.g., according to pragmatic accounts of the accuracy of one’s partial beliefs, one begins by defining a preference ordering on the agent’s space of acts and outcomes. If the preference ordering is consistent with the Kolmogorov axioms\textsuperscript{28}, then one can set up a representation theorem from which the agent’s subjective probability and utility measures (i.e., their expected utility measure) can be derived.\textsuperscript{29} The epistemic utility associated with the pragmatic approach is, generally, utility maximization.

By contrast to the pragmatic approach, the epistemic approach to measuring the accuracy of one’s beliefs is grounded in the notion of dominance (cf. Joyce, 1998; 2009). According to the epistemic approach, there is an ideal, or vindicated, probability concerning a proposition’s obtaining, and if an agent’s subjective probability measure does not satisfy the Kolmogorov axioms, then one can prove that it will always be dominated by a distinct measure; i.e. it will always be the case that a distinct subjective probability measure will be closer to the vindicated world than one’s own. The epistemic utility associated with the epistemic approach is thus the minimization of in-

\textsuperscript{28}Namely: normality (which states that the probability of a tautology maps to 1); non-negativity (which states that the probability operator must take a non-negative value); additivity (which states that for all disjoint probability densities, the probability of their union is equal to the probability of the first density added to the probability of the second); and conditionalization (which states that the probability of \( \phi \) conditional on \( \psi \) equals the probability of the intersection of \( \phi \) and \( \psi \), divided by the probability of \( \psi \)).

\textsuperscript{29}Cf. Ramsey (1926); Savage (1954); and Jeffrey (1965).
accuracy (cf. Pettigrew, 2014).30

Wright notes that rational trust subserving epistemic entitlement will be
pragmatic, and makes the intriguing point that ‘pragmatic reasons are not
a special genre of reason, to be contrasted with e.g. epistemic, prudential,
and moral reasons’ (2012: 484). He provides an example according to which
one might be impelled to prefer the ‘alleviation of Third world suffering’ to
one’s own ‘eternal bliss’ (op. cit.); and so presumably has the pragmatic
approach to expected utility in mind. The intriguing point to note, however,
is that epistemic utility is variegated; one’s epistemic utility might consist,
e.g., in both the reduction of epistemic inaccuracy and in the satisfaction
of one’s preferences. Wright concludes that there is thus ‘no good cause to
deny certain kinds of pragmatic reason the title ‘epistemic’. This will be the
case where, in the slot in the structure of the reasons for an action that is
to be filled by the desires of the agent, the relevant desires are focused on
epistemic goods and goals’ (op. cit.).

Third, and most crucially: The very idea of expected epistemic utility in
the setting of decision theory makes implicit appeal to the notion of possible
worlds. The full and partial beliefs of an agent will have to be defined on
a probability distribution, i.e. a set of epistemically possible worlds. The
philosophical significance of this point is that it demonstrates how Hale and
Wright’s appeal to default, rational entitlement to trust that abstraction
principles are true converges with the modal approach to the epistemology
of mathematics advanced in chapter 10. The latter proceeds by examining
undecidable sentences via the epistemic interpretation of two-dimensional inten-
tensional semantics. The latter can be understood as recording the thought
that the semantic value of a proposition relative to a first parameter (a con-
text) which ranges over epistemically possible worlds, will constrain the se-
monic value of the proposition relative to a second parameter (an index)

30 The distinction between the epistemic (also referred to as the alethic) and the prag-
matic approaches to epistemic utility is anticipated by Clifford (1877) and James (1896),
with Clifford endorsing the epistemic approach, and James the pragmatic. The distance
measures comprising the scoring rules for the minimization of inaccuracy are examined in,
inter alia, Fitelson (2001); Leitgeb and Pettigrew (2010); and Moss (2011). A generaliza-
tion of Joyce’s argument for probabilism to models of non-classical logic is examined in
Paris (2001) and Williams (2012). A dominance-based approach to decision theory is ex-
amined in Easwaran (2014), and a dominance-based approach to the notion of coherence –
which can accommodate phenomena such as the preface paradox, and is thus weaker than
the notion of consistency in an agent’s belief set – is examined in Easwaran and Fitelson
which ranges over metaphysically possible worlds. The formal clauses for epistemic and metaphysical mathematical modalities are as follows:

Let $C$ denote a set of epistemically possibilities, such that $[\phi]^c \subseteq C$;
($\phi$ is a formula encoding a state of information at an epistemically possible world).

- $\text{pri}(x) = \lambda c.[x]^c$;
  (This is an epistemic intension, such that the two parameters relative to which $x$ – a propositional variable – obtains its value are epistemically possible worlds).

- $\text{sec}(x) = \lambda w.[x]^w$;
  (This is a metaphysical intension, such that the two parameters relative to which $x$ obtains its value are metaphysically possible worlds).

Then:

- **Epistemic Mathematical Necessity**
  
  
  $[\Box \phi]^{c,w} = 1 \iff \forall c'[\phi]^{c,c'} = 1$
  
  ($\phi$ is true at all points in epistemic modal space).

- **Epistemic Mathematical Possibility**
  
  $[\Diamond \phi] \neq \emptyset \iff [\neg \Box \neg \phi] = 1$
  
  ($\phi$ might be true if and only if it is not epistemically necessary for $\phi$ to be false).

Epistemic mathematical modality is constrained by consistency, and the formal techniques of provability and forcing. A mathematical formula is metaphysically impossible, if it can be disproved or induces inconsistency in a model.

- **Convergence**
  
  $\forall c \exists w[\phi]^{c,w} = 1$
  
  (the value of $x$ is relative to a parameter for the space of epistemically possible worlds. The value of $x$ relative to the first parameter determines the value of $x$ relative to the second parameter for the space of metaphysical possibilities).
According, then, to the latter, the possibility of deciding mathematical propositions which are currently undecidable relative to a background mathematical language such as ZFC should be two-dimensional. The epistemic possibility of deciding Orey sentences can thus be a guide to the metaphysical possibility thereof.\footnote{See Kanamori (2008) and Woodin (2010), for further discussion of the mathematical properties at issue.} Further, both the numerical term-forming operator, $N_x$, in abstraction principles, as well as entire abstraction principles themselves, can receive a two-dimensional treatment, such that the value of numerical terms relative to epistemic possibilities considered as actual can determine the value of numerical terms relative to metaphysical possibilities, and the epistemic possibility of an abstraction principle’s truth can determine the metaphysical possibility thereof.

The convergence between Wright’s and my approaches consists, then, in that – on both approaches – there is a set of epistemically possible worlds. In the former case, the epistemically possible worlds subserve the preference rankings for the definability of expected epistemic utility. Epistemic mathematical modality is thus constitutive of the notion of rational entitlement to which Hale and Wright appeal, and – in virtue of its convergence with the two-dimensional intensional semantics here proffered – epistemically possible worlds can serve as a guide to the metaphysical mathematical possibility that mathematical propositions, such as abstraction principles for cardinals, reals, and sets, are true.

8.3.1 Epistemic Two-dimensional Truthmaker Semantics

If one prefers hyperintensional semantics to possible worlds semantics – in order e.g. to avoid the situation in intensional semantics according to which all necessary formulas express the same proposition because they are true at all possible worlds – one can avail of the following epistemic two-dimensional truthmaker semantics, which specifies a notion of exact verification in a state space and where states are parts of whole worlds (Fine 2017a,b; Hawke and Özgün, forthcoming). According to truthmaker semantics for epistemic logic, a modalized state space model is a tuple $\langle S, P, \leq, v \rangle$, where $S$ is a non-empty set of states, i.e. parts of the elements in $A$ in the foregoing epistemic modal algebra $U$, $P$ is the subspace of possible states where states $s$ and $t$ are
compatible when \( s \sqcup t \in P \), \( \leq \) is a partial order, and \( v: \text{Prop} \to (2^S \times 2^S) \) assigns a bilateral proposition \( \langle p^+, p^- \rangle \) to each atom \( p \in \text{Prop} \) with \( p^+ \) and \( p^- \) incompatible (Hawke and Özgün, forthcoming: 10-11). Exact verification \( (\vdash) \) and exact falsification \( (\dashv) \) are recursively defined as follows (Fine, 2017a: 19; Hawke and Özgün, forthcoming: 11):

\[
\begin{align*}
& s \vdash p \text{ if } s \in \llbracket p \rrbracket^+ \\
& \quad (s \text{ verifies } p, \text{ if } s \text{ is a truthmaker for } p \text{ i.e. if } s \text{ is in } p\text{'s extension}); \\
& s \dashv p \text{ if } s \in \llbracket p \rrbracket^- \\
& \quad (s \text{ falsifies } p, \text{ if } s \text{ is a falsifier for } p \text{ i.e. if } s \text{ is in } p\text{'s anti-extension}); \\
& s \vdash \neg p \text{ if } s \dashv p \\
& \quad (s \text{ verifies not } p, \text{ if } s \text{ falsifies } p); \\
& s \dashv \neg p \text{ if } s \vdash p \\
& \quad (s \text{ falsifies not } p, \text{ if } s \text{ verifies } p); \\
& s \vdash p \land q \text{ if } \exists t,u, t \vdash p, u \vdash q, \text{ and } s = t \sqcap u \\
& \quad (s \text{ verifies } p \text{ and } q, \text{ if } s \text{ is the fusion of states, } t \text{ and } u, t \text{ verifies } p, \text{ and } u \text{ verifies } q); \\
& s \dashv p \land q \text{ if } s \vdash p \text{ or } s \dashv q \\
& \quad (s \text{ falsifies } p \text{ or } q, \text{ if } s \text{ falsifies } p \text{ or } s \text{ falsifies } q); \\
& s \vdash p \lor q \text{ if } s \vdash p \text{ or } s \vdash q \\
& \quad (s \text{ verifies } p \text{ or } q, \text{ if } s \text{ verifies } p \text{ or } s \text{ verifies } q); \\
& s \dashv p \lor q \text{ if } \exists t,u, t \dashv p, u \dashv q, \text{ and } s = t \sqcap u \\
& \quad (s \text{ falsifies } p \text{ or } q, \text{ if } s \text{ is the state overlapping the states } t \text{ and } u, t \text{ falsifies } p, \text{ and } u \text{ falsifies } p, \text{ or } u \text{ falsifies } q); \\
& s \text{ exactly verifies } p \text{ if and only if } s \vdash p \text{ if } s \in \llbracket p \rrbracket; \\
& s \text{ inexactly verifies } p \text{ if and only if } s \triangleright p \text{ if } \exists s' \sqsubseteq S, s' \vdash p; \text{ and } \\
& s \text{ loosely verifies } p \text{ if and only if, } \forall t, s.t. s \sqcup t, s \sqcup t \vdash p \text{ (35-36)}; \\
& s \vdash A \phi \text{ if and only if for all } t \in P \text{ there is a } t' \in P \text{ such that } t' \sqcup t \in P \text{ and } t' \vdash \phi; \\
& s \dashv A \phi \text{ if and only if there is a } t \in P \text{ such that for all } u \in P \text{ either } t \sqcup u \in P \text{ or } u \dashv \phi, \text{ where } A \phi \text{ or } \Box \phi \text{ denotes the apriority of } \phi.
\end{align*}
\]

In order to account for two-dimensional indexing, we augment the model, \( M \), with a second state space, \( S^* \), on which we define both a new parthood relation, \( \leq^* \), and partial function, \( V^* \), which serves to map propositions in a domain, \( D \), to pairs of subsets of \( S^* \), \( \{1,0\} \), i.e. the verifier and falsifier of \( p \), such that \( \llbracket P \rrbracket^+ = 1 \text{ and } \llbracket p \rrbracket^- = 0 \). Thus, \( M = \langle S, S^*, D, \leq, \leq^*, V, V^* \rangle \). The two-dimensional hyperintensional profile of propositions may then be recorded by defining the value of \( p \) relative to two parameters, \( c,i \): \( c \) ranges over subsets of \( S \), and \( i \) ranges over subsets of \( S^* \).
(*) $M, s \in S, s^* \in S^* \vdash p$ iff:

(i) $\exists c_s [[p]]^{c, c} = 1$ if $s \in [[p]]^+$; and

(ii) $\exists i_{s^*} [[p]]^{c, i} = 1$ if $s^* \in [[p]]^+$

(Distinct states, $s, s^*$, from distinct state spaces, $S, S^*$, provide a multidimensional verification for a proposition, $p$, if the value of $p$ is provided a truthmaker by $s$. The value of $p$ as verified by $s$ determines the value of $p$ as verified by $s^*$).

We say that $p$ is hyper-rigid iff:

(*) $M, s \in S, s^* \in S^* \vdash p$ iff:

(i) $\forall c', s \in [[p]]^{c, c'} = 1$ if $s \in [[p]]^+$; and

(ii) $\forall i_{s^*} [[p]]^{c, i} = 1$ if $s^* \in [[p]]^+$

The foregoing provides a two-dimensional hyperintensional semantic framework within which to interpret the values of a proposition. In order to account for partial contents, we define the values of subpropositional entities relative again to tuples of states from the distinct state spaces in our model:

$s$ is a two-dimensional exact truthmaker of $p$ if and only if (*)

$s$ is a two-dimensional inexact truthmaker of $p$ if and only if $\exists s' \sqsubseteq S, s \to s', s^* \vdash p$ and such that

$\exists c', s' \in [[p]]^{c, c'} = 1$ if $s' \in [[p]]^+$, and

$\exists i_{s^*} [[p]]^{c, i} = 1$ if $s^* \in [[p]]^+$;

$s$ is a two-dimensional loose truthmaker of $p$ if and only if, $\exists t, s.t. s \sqcup t, s \sqcup t \vdash p$:

$\exists c_{s,t} [[p]]^{c, c} = 1$ if $s' \in [[p]]^+$, and

$\exists i_{s^*} [[p]]^{c, i} = 1$ if $s^* \in [[p]]^+$.

Epistemic (primary), subjunctive (secondary), and 2D hyperintensions can be defined as follows, where hyperintensions are functions from states to extensions, and intensions are functions from worlds to extensions:

- Epistemic Hyperintension:
  $$\text{pri}(x) = \lambda s. [[x]]^{s, s}, \text{ with } s \text{ a state in the state space defined over the foregoing epistemic modal algebra, } U;$$

- Subjunctive Hyperintension:
  $$\text{sec}_{i_\theta}(x) = \lambda i. [[x]]^{\theta, i}, \text{ with } i \text{ a state in metaphysical state space } I;$$

- 2D-Hyperintension:
  $$\text{2D}(x) = \lambda s \lambda w [[x]]^{s, i} = 1.$$
8.4 Concluding Remarks

In this essay, I have endeavored to provide an account of the modal foundations of mathematical platonism. Hale and Wright’s objections to the idea that Necessitism cannot account for how possibility and actuality might converge were shown to be readily answered. In response, further, to Hale and Wright’s objections to the role of epistemic and metaphysical modalities in countenancing the truth of abstraction principles and the success of mathematical predicate reference, I demonstrated how my two-dimensional intensional and hyperintensional approaches to the epistemology of mathematics, augmented with Necessitism, are consistent with Hale and Wright’s conception of the epistemic entitlement rationally to trust that abstraction principles are true. Epistemic and metaphysical states and possibilities may thus be shown to play a constitutive role in vindicating the reality of mathematical objects and truth, and in explaining our possible knowledge thereof.
Chapter 9

Ω-Logicism: Automata, Neo-Logicism, and Set-theoretic Realism

This essay examines the philosophical significance of the consequence relation defined in the Ω-logic for set-theoretic languages. I argue that, as with second-order logic, the modal profile of validity in Ω-Logic enables the property to be epistemically tractable. Because of the duality between coalgebras and algebras, Boolean-valued models of set theory can be interpreted as coalgebras. In Section 2, I demonstrate how the modal profile of Ω-logical validity can be countenanced within a coalgebraic logic, and how Ω-logical validity can further be defined via automata. In Section 3, the philosophical significance of the characterization of the modal profile of Ω-logical validity for the philosophy of mathematics is examined. I argue (i) that it vindicates a type of neo-logicism with regard to mathematical truth in the set-theoretic multiverse, and (ii) that it provides a modal and computational account of formal grasp of the concept of ‘set’, adducing in favor of a realist conception of the cumulative hierarchy of sets. Section 4 provides concluding remarks.

9.1 Definitions

In this section, I define the axioms of Zermelo-Fraenkel set theory with choice. I define the mathematical properties of the large cardinal axioms to which ZFC can be adjoined, and I provide a detailed characterization of the prop-
roperties of Ω-logic for ZFC. Because coalgebras are dual to Boolean-valued algebraic models of Ω-logic, a category of coalgebraic logic is then characterized which models both modal logic and deterministic automata. Modal coalgebraic models of automata are then argued to provide a precise characterization of the modal and computational profiles of Ω-logical validity.

9.1.1 Axioms

- **Extensionality**
  \[ \forall x,y. (\forall z. z \in x \iff z \in y) \rightarrow x = y \]

- **Empty Set**
  \[ \exists x. \forall y. y \notin x \]

- **Pairing**
  \[ \forall x, y. \exists z. \forall w. w \in z \iff w = x \lor w = y \]

- **Union**
  \[ \forall x. \exists y. \forall z. z \in y \iff \exists w. w \in x \land z \in w \]

- **Powerset**
  \[ \forall x. \exists y. \forall z. z \in y \iff z \subseteq x \]

- **Separation (with \( \vec{x} \) a parameter)**
  \[ \forall \vec{x}, y. \exists z. \forall w. w \in z \iff w \in y \land A(w, \vec{x}) \]

- **Infinity**
  \[ \exists x. \emptyset \in x \land \forall y. y \in x \rightarrow y \cup \{y\} \in x \]

- **Foundation**
  \[ \forall x. (\exists y. y \in x) \rightarrow \exists y. \forall z. z \notin y \]

- **Replacement**
  \[ \forall x, y, [\forall z. \exists! w. A(z, w, y)] \rightarrow \exists u. \forall w. w \in u \iff \exists z. A(z, w, y) \]

- **Choice**
  \[ \forall x. \emptyset \notin x \rightarrow \exists f. (x \rightarrow \cup x). \forall y. y \in x. f(y) \in y \]

\(^1\)For a standard presentation, see Jech (2003). The presentation here follows Avigad (2021). For detailed, historical discussion, see Maddy (1988,a).
9.1.2 Large Cardinals

Borel sets of reals are subsets of $\omega^\omega$ or $\mathbb{R}$, closed under countable intersections and unions.\(^2\) For all ordinals, $a$, such that $0 < a < \omega_1$, and $b < a$, $\Sigma^0_a$ denotes the open subsets of $\omega^\omega$ formed under countable unions of sets in $\Pi^0_b$, and $\Pi^0_a$ denotes the closed subsets of $\omega^\omega$ formed under countable intersections of $\Sigma^0_b$.

Projective sets of reals are subsets of $\omega^\omega$, formed by complementations ($\omega^\omega - u$, for $u \subseteq \omega^\omega$) and projections $[p(u) = \{\langle x_1, \ldots, x_n \rangle \in \omega^\omega \mid \exists y (\langle x_1, \ldots, x_n, y \rangle \in u)\}]$. For all ordinals $a$, such that $0 < a < \omega$, $\Pi^1_0$ denotes closed subsets of $\omega^\omega$; $\Pi^1_a$ is formed by taking complements of the open subsets of $\omega^\omega$, $\Sigma^1_a$; and $\Sigma^1_{a+1}$ is formed by taking projections of sets in $\Pi^1_a$.

The full power set operation defines the cumulative hierarchy of sets, $V$, such that $V_0 = \emptyset$; $V_{a+1} = P(V_a)$; and $V_\lambda = \bigcup_{a < \lambda} V_a$.

In the inner model program (cf. Woodin, 2001a,b, 2010, 2011; Kanamori, 2012,a,b), the definable power set operation defines the constructible universe, $L(\mathbb{R})$, in the universe of sets $V$, where the sets are transitive such that $a \in C \iff a \subseteq C$; $L(\mathbb{R}) = V_{\omega+1}$; $L_{a+1}(\mathbb{R}) = \text{Def}(L_{a}(\mathbb{R}))$; and $L_\lambda(\mathbb{R}) = \bigcup_{a < \lambda} (L_{a}(\mathbb{R}))$.

Via inner models, Gödel (1940) proves the consistency of the generalized continuum hypothesis, $\aleph_\alpha = \aleph_{\alpha+1}$, as well as the axiom of choice, relative to the axioms of ZFC. However, for a countable transitive set of ordinals, $M$, in a model of ZF without choice, one can define a generic set, $G$, such that, for all formulas, $\phi$, either $\phi$ or $\neg \phi$ is forced by a condition, $f$, in $G$. Let $M[G] = \bigcup_{a < \kappa} M_a[G]$, such that $M_0[G] = \{G\}$; with $\lambda < \kappa$, $M_\lambda[G] = \bigcup_{a < \lambda} M_a[G]$; and $M_{a+1}[G] = V_a \cap M_a[G]$.\(^3\) $G$ is a Cohen real over $M$, and comprises a set-forcing extension of $M$. The relation of set-forcing, $\Vdash$, can then be defined in the ground model, $M$, such that the forcing condition, $f$, is a function from a finite subset of $\omega$ into $\{0,1\}$, and $f \Vdash u \in G$ if $f(u) = 1$ and $f \Vdash u \notin G$ if $f(u) = 0$. The cardinalities of an open dense ground model, $M$, and a generic extension, $G$, are identical, only if the countable chain condition (c.c.c.) is satisfied, such that, given a chain -- i.e., a linearly ordered subset of a partially ordered (reflexive, antisymmetric, transitive) set -- there is a countable, maximal antichain consisting of pairwise incompatible forcing conditions. Via set-forcing extensions, Cohen (1963, 1964) constructs

---

\(^2\)See Koellner (2013), for the presentation, and for further discussion, of the definitions in this and the subsequent paragraph.

\(^3\)See Kanamori (2012,a: 2.1; 2012,b: 4.1), for further discussion.
a model of ZF which negates the generalized continuum hypothesis, and thus proves the independence thereof relative to the axioms of ZF.\textsuperscript{4}

Gödel (1946/1990: 1-2) proposes that the value of Orey sentences such as the GCH might yet be decidable, if one avails of stronger theories to which new axioms of infinity – i.e., large cardinal axioms – are adjoined.\textsuperscript{5} He writes that: ‘In set theory, e.g., the successive extensions can be represented by stronger and stronger axioms of infinity. It is certainly impossible to give a combinatorial and decidable characterization of what an axiom of infinity is; but there might exist, e.g., a characterization of the following sort: An axiom of infinity is a proposition which has a certain (decidable) formal structure and which in addition is true. Such a concept of demonstrability might have the required closure property, i.e. the following could be true: Any proof for a set-theoretic theorem in the next higher system above set theory \ldots is replaceable by a proof from such an axiom of infinity. It is not impossible that for such a concept of demonstrability some completeness theorem would hold which would say that every proposition expressible in set theory is decidable from present axioms plus some true assertion about the largeness of the universe of sets.'

For cardinals, $x,a,C, C \subseteq a$ is closed unbounded in $a$, if it is closed \textit{if} $x < C$ and $\bigcup (C \cap a) = a$, \textit{then} $a \in C$ and unbounded ($\bigcup C = a$) (Kanamori, op. cit.: 360). A cardinal, $S$, is stationary in $a$, if, for any closed unbounded $C \subseteq a$, $C \cap S \neq \emptyset$ (op. cit.). An ideal is a subset of a set closed under countable unions, whereas filters are subsets closed under countable intersections (361). A cardinal $\kappa$ is regular if the cofinality of $\kappa$ – comprised of the unions of sets with cardinality less than $\kappa$ – is identical to $\kappa$. Uncountable regular limit cardinals are weakly inaccessible (op. cit.). A strongly inaccessible cardinal is regular and has a strong limit, such that if $\lambda < \kappa$, then $2^\lambda < \kappa$ (op. cit.).

Large cardinal axioms are defined by elementary embeddings.\textsuperscript{6} Elementary embeddings can be defined thus. For models $A,B$, and conditions $\phi$, $j$: $A \rightarrow B$, $\phi(a_1, \ldots, a_n)$ in $A$ if and only if $\phi(j(a_1), \ldots, j(a_n))$ in $B$ (363). A measurable cardinal is defined as the ordinal denoted by the critical point of $j$,
Let $\kappa$ be a cardinal, and $\eta > \kappa$ an ordinal. $\kappa$ is then $\eta$-strong, if there is a transitive class $M$ and an elementary embedding, $j: V \rightarrow M$, such that $\text{crit}(j) = \kappa$, $j(\kappa) > \eta$, and $V_\eta \subseteq M$ (Koellner and Woodin, op. cit.).

$\kappa$ is strong if and only if, for all $\eta$, it is $\eta$-strong (op. cit.).

If $A$ is a class, $\kappa$ is $\eta$-$A$-strong, if there is a $j: V \rightarrow M$, such that $\kappa$ is $\eta$-strong and $j(A \cap V_\kappa) \cap V_\eta = A \cap V_\eta$ (op. cit.).

$\kappa$ is a Woodin cardinal, if $\kappa$ is strongly inaccessible, and for all $A \subseteq V_\kappa$, there is a cardinal $\kappa_A < \kappa$, such that $\kappa_A$ is $\eta$-$A$-strong, for all $\eta$ such that $\kappa_\eta$, $\eta < \kappa$ (Koellner and Woodin, op. cit.: 8).

$\kappa$ is superstrong, if $j: V \rightarrow M$, such that $\text{crit}(j) = \kappa$ and $V_{j(\kappa)} \subseteq M$, which entails that there are arbitrarily large Woodin cardinals below $\kappa$ (op. cit.).

Large cardinal axioms can then be defined as follows.

(i) $\exists x \Phi$ is a large cardinal axiom, because:

(ii) if $\kappa$ is a cardinal, such that $V \models \Phi(\kappa)$, then $\kappa$ is strongly inaccessible; and

(iii) for all generic partial orders $P \in V_\kappa$, $V^P \models \Phi(\kappa)$; $I_{NS}$ is a non-stationary ideal; $A^G$ is the canonical representation of reals in $L(\mathbb{R})$, i.e. the interpretation of $A$ in $M[G]$; $H(\kappa)$ is comprised of all of the sets whose transitive closure is $< \kappa$ (cf. Woodin, 2001a: 569); and $L(\mathbb{R})^P_{\text{max}} \models \langle H(\omega_2), \in, I_{NS}, A^G \rangle \models '\phi'$. $P$ is a homogeneous partial order in $L(\mathbb{R})$, such that the generic extension of $L(\mathbb{R})^P$ inherits the generic invariance, i.e., the absoluteness, of $L(\mathbb{R})$. Thus, $L(\mathbb{R})^P_{\text{max}}$ is (i) effectively complete, i.e. invariant under set-forcing extensions; and (ii) maximal, i.e. satisfies all $\Pi_2$-sentences and is thus consistent by set-forcing over ground models (Woodin, ms: 28).

Assume ZFC and that there is a proper class of Woodin cardinals; $A \in \mathbb{P}(\mathbb{R}) \cap L(\mathbb{R})$; $\phi$ is a $\Pi_2$-sentence; and $V(G)$, s.t. $\langle H(\omega_2), \in, I_{NS}, A^G \rangle \models '\phi'$. Then, it can be proven that $L(\mathbb{R})^P_{\text{max}} \models \langle H(\omega_2), \in, I_{NS}, A^G \rangle \models '\phi'$, where $'\phi' := \exists A \in \Gamma^\infty \langle H(\omega_1), \in, A \rangle \models \psi$.

The axiom of determinacy (AD) states that every set of reals, $a \subseteq \omega^\omega$ is determined, where $\kappa$ is determined if it is decidable.

Woodin’s (1999) Axiom (*) can be thus countenanced:

AD$^L(\mathbb{R})$ and $L[(\mathbb{P}(\omega_1)]$ is a $P_{\text{max}}$-generic extension of $L(\mathbb{R})$,

from which it can be derived that $2^{\aleph_0} = \aleph_2$. Thus, $\neg \text{CH}$; and so $\text{CH}$ is absolutely decidable.
9.1.3 $\Omega$-Logic

For partial orders, $\mathbb{P}$, let $V^\mathbb{P} = V^\mathbb{B}$, where $\mathbb{B}$ is the regular open completion of $(\mathbb{P})$. $\mathbb{M}_a = (V_a)^M$ and $\mathbb{M}^\mathbb{B}_a = (V^\mathbb{B}_a)^M = (V_a)^{\mathbb{M}^\mathbb{B}_a}$. $\text{Sent}$ denotes a set of sentences in a first-order language of set theory. $\text{Tu}\{\phi\}$ is a set of sentences extending ZFC. $c.t.m$ abbreviates the notion of a countable transitive $\epsilon$-model. $c.B.a.$ abbreviates the notion of a complete Boolean algebra.

Define a $c.B.a.$ in $V$, such that $V^\mathbb{B}_a$. Let $V^\mathbb{B}_0 = \emptyset$; $V^\mathbb{B}_\lambda = \bigcup_{b<\lambda} V^\mathbb{B}_b$, with $\lambda$ a limit ordinal; $V^\mathbb{B}_{a+1} = \{f : X \to B \mid X \subseteq V^\mathbb{B}_a\}$; and $V^\mathbb{B} = \bigcup_{a \in \text{On}} V^\mathbb{B}_a$. $\phi$ is true in $V^\mathbb{B}$, if its Boolean-value is $1^\mathbb{B}$, if and only if $V^\mathbb{B}_a \models \phi$ iff $\llbracket \phi \rrbracket^\mathbb{B} = 1$.

Thus, for all ordinals, $a$, and every $c.B.a.$ $\mathbb{B}$, $V^\mathbb{B}_a \equiv (V_a)^{V^\mathbb{B}_a}$ iff for all $x \in V^\mathbb{B}$, $\exists y \in V^\mathbb{B}[x = y]^\mathbb{B} = 1^\mathbb{B}$ iff $\llbracket x \in V^\mathbb{B} \rrbracket^\mathbb{B} = 1^\mathbb{B}$.

Then, $V^\mathbb{B}_a \models \phi$ iff $V^\mathbb{B} \models 'V_a \models \phi'$. $\Omega$-logical validity can then be defined as follows:

For $\text{Tu}\{\phi\} \subseteq \text{Sent}$,

$T \models_\Omega \phi$, if for all ordinals, $a$, and $c.B.a.$ $\mathbb{B}$, if $V^\mathbb{B}_a \models T$, then $V^\mathbb{B}_a \models \phi$.

Supposing that there exists a proper class of Woodin cardinals and if $\text{Tu}\{\phi\} \subseteq \text{Sent}$, then for all set-forcing conditions, $\mathbb{P}$:

$T \models_\Omega \phi$ iff $V^T \models 'T \models_\Omega \phi'$,

where $T \models_\Omega \phi \equiv \emptyset \models 'T \models_\Omega \phi'$.

The $\Omega$-Conjecture states that $V \models_\Omega \phi$ iff $V^\mathbb{B} \models_\Omega \phi$ (Woodin, ms). Thus, $\Omega$-logical validity is invariant in all set-forcing extensions of ground models in the set-theoretic multiverse.

The soundness of $\Omega$-Logic is defined by universally Baire sets of reals. For a cardinal, $e$, let a set $A$ be $e$-universally Baire, if for all partial orders $\mathbb{P}$ of cardinality $e$, there exist trees, $S$ and $T$ on $\omega \times \lambda$, such that $A = p[T]$ and if $G \subseteq \mathbb{P}$ is generic, then $p[T]^G = R^G - p[S]^G$ (Koellner, 2013). $A$ is universally Baire, if it is $e$-universally Baire for all $e$ (op. cit.).

$\Omega$-Logic is sound, such that $V \models_\Omega \phi \to V \models_\Omega \phi$. However, the completeness of $\Omega$-Logic has yet to be resolved.

Finally, in category theory, a category $C$ is comprised of a class $\text{Ob}(C)$ of objects a family of arrows for each pair of objects $C(A,B)$ (Venema, 2007: 421). A functor from a category $C$ to a category $D$, $E : C \to D$, is an operation mapping objects and arrows of $C$ to objects and arrows of $D$ (422). An endofunctor on $C$ is a functor, $E : C \to C$ (op. cit.).

---

7 The definitions in this section follow the presentation in Bagaria et al. (2006).
A $E$-coalgebra is a pair $A = (A, \mu)$, with $A$ an object of $C$ referred to as the carrier of $A$, and $\mu: A \rightarrow E(A)$ is an arrow in $C$, referred to as the transition map of $A$ (390).

$A = (A, \mu: A \rightarrow E(A))$ is dual to the category of algebras over the functor $\mu$ (417-418). If $\mu$ is a functor on categories of sets, then coalgebraic models are dual to Boolean-algebraic models of $\Omega$-logical validity.

The significance of the foregoing is that coalgebraic models may themselves be availed of in order to define modal logic and automata. Coalgebras provide therefore a setting in which the Boolean-valued models of set theory, the modal profile of $\Omega$-logical validity, and automata can be interdefined. In what follows, $A$ will comprise the coalgebraic model – dual to the complete Boolean-valued algebras defined in the $\Omega$-Logic of $ZFC$ – in which modal similarity types and automata are definable. As a coalgebraic model of modal logic, $A$ can be defined as follows (407):

For a set of formulas, $\Phi$, let $\nabla \Phi := \Box \lor \Phi \land \Diamond \Phi$, where $\Diamond \Phi$ denotes the set $\{\Diamond \phi \mid \phi \in \Phi\}$ (op. cit.). Then,

$\Diamond \phi \equiv \nabla \{\phi, T\}$,

$\Box \phi \equiv \nabla \emptyset \lor \nabla \phi$ (op. cit.)

$[\nabla \Phi] = \{w \in W \mid R[w] \subseteq \bigcup \{[\phi] \mid \phi \in \Phi\}$ and $\forall \phi \in \Phi, [\phi] \cap R[w] \neq \emptyset\}$ (Fontaine, 2010: 17).

Let an $E$-coalgebraic modal model, $A = (S, \lambda, R[\cdot])$, such that $S, s \vdash \nabla \Phi$ if and only if, for all (some) successors $\sigma$ of $s \in S$, $[\Phi, \sigma(s) \in E(\vdash_{A})]$ (Venema, 2007: 407), with $E(\vdash_{A})$ a relation lifting of the satisfaction relation $\vdash_{A} \subseteq S \times \Phi$. Let a functor, $K$, be such that there is a relation $K! \subseteq K(A) \times K(A')$ (17). Let $Z$ be a binary relation s.t. $Z \subseteq A \times A$ and $\varphi!Z \subseteq \varphi(A) \times \varphi(A')$, with

$\varphi!Z := \{(X,X') \mid \forall x \in X \exists x' \in X' \text{ with } (x,x') \in Z \land \forall x' \in X' \exists x \in X \text{ with } (x,x') \in Z\}$ (op. cit.). Then, we can define the relation lifting, $K!$, as follows:

$K! := \{([\pi, X], ([\pi', X']) \mid \pi = \pi' \text{ and } (X,X') \in \varphi!Z\}$ (Venema, 2012: 17).

A coalgebraic model of deterministic automata can be thus defined (391).

An automaton is a tuple, $A = (A, a_I, C, \delta, F)$, such that $A$ is the state space of the automaton $A; a_I \in A$ is the automaton’s initial state; $C$ is the coding for the automaton’s alphabet, mapping numerals to properties of the natural numbers; $\delta$: $A \times C \rightarrow A$ is a transition function, and $F \subseteq A$ is the collection of admissible states, where $F$ maps $A$ to $\{1,0\}$, such that $F$: $A \rightarrow 1$ if $a \in F$ and $A \rightarrow 0$ if $a \notin F$ (op. cit.). The determinacy of coalgebraic automata, the category of which is dual to the Set category satisfying $\Omega$-logical consequence, is secured by the existence of Woodin cardinals: Assuming $ZFC$, that $\lambda$ is
a limit of Woodin cardinals, that there is a generic, set-forcing extension $G \subseteq$ the collapse of $\omega < \lambda$, and that $\mathbb{R}^* = \bigcup\{R^G[a] \mid a < \lambda\}$, then $\mathbb{R}^* \models$ the axiom of determinacy (AD) (Koellner and Woodin, op. cit.: 10).

Modal automata are defined over a modal one-step language (Fontaine and Venema, 2018: 3.1-3.2; Venema, 2020: 7.2). With $A$ being a set of propositional variables the set, $\text{Latt}(X)$, of lattice terms over $X$ has the following grammar:

$$\pi ::= \bot \mid \top \mid x \mid \pi \land \pi \mid \pi \lor \pi,$$

with $x \in X$ and $\pi \in \text{Latt}(A)$ (op. cit.).

The set, $\text{1ML}(A)$, of modal one-step formulas over $A$ has the following grammar:

$$\alpha \in A ::= \bot \mid \top \mid \lozenge \pi \mid \Box \pi \mid \alpha \land \alpha \mid \alpha \lor \alpha \ (\text{op. cit.}).$$

A modal $P$-automaton $A$ is a triple, $(A, \Theta, a_I)$, with $A$ a non-empty finite set of states, $a_I \in A$ an initial state, and the transition map

$$\Theta: A \times \wp(P) \to \text{1ML}(A)$$

maps states to modal one-step formulas, with $\wp(P)$ the powerset of the set of proposition letters, $P$ (op. cit.: 7.3).

Finally, $A = \langle A, \alpha:A \to E(A) \rangle$ is dual to the category of algebras over the functor $\alpha$ (417-418). For a category $C$, object $A$, and endofunctor $E$, define a new arrow, $\alpha$, s.t. $\alpha:EA \to A$. A homomorphism, $f$, can further be defined between algebras $\langle A, \alpha \rangle$, and $\langle B, \beta \rangle$. Then, for the category of algebras, the following commutative square can be defined: (i) $EA \to EB (Ef)$; (ii) $EA \to A (\alpha)$; (iii) $EB \to B (\beta)$; and (iv) $A \to B (f)$ (cf. Hughes, 2001: 7-8). The same commutative square holds for the category of coalgebras, such that the latter are defined by inverting the direction of the morphisms in both (ii) $[A \to EA (\alpha)]$, and (iii) $[B \to EB (\beta)]$ (op. cit.).

Thus, $A$ is the coalgebraic category for modal, deterministic automata, dual to the complete Boolean-valued algebraic models of $\Omega$-logical validity, as defined in the category of sets.

Leach-Krouse (ms) defines the modal logic of $\Omega$-consequence as satisfying the following axioms:

For a theory $T$ and with $\Box \phi := T^{\Box} \alpha \vDash ZFC \Rightarrow T^{\Box} \alpha \vDash \phi$,

$ZFC \vdash \phi \Rightarrow ZFC \vdash \Box \phi$

$ZFC \vdash \Box (\phi \to \psi) \to (\Box \phi \to \Box \psi)$

143
\[
\begin{align*}
\text{ZFC} \vdash \Box \phi \rightarrow \phi & \Rightarrow \text{ZFC} \vdash \phi \\
\text{ZFC} \vdash \Box \phi \rightarrow \Box \Box \phi & \\
\text{ZFC} \vdash \Box (\Box \phi \rightarrow \phi) & \rightarrow \Box \phi \\
\Box (\Box \phi \rightarrow \psi) & \lor \Box (\Box \psi \land \psi \rightarrow \phi), \text{ where this clause added to GL is the logic of } ^*\text{true in all } V_\kappa \text{ for all } \kappa \text{ strongly inaccessible}^*_\text{ in ZFC.}
\end{align*}
\]

9.2 Modal Coalgebraic Automata and the Philosophy of Mathematics

This section examines the philosophical significance of modal coalgebraic automata and the Boolean-valued models of set-theoretic languages to which they are dual. I argue that, similarly to second-order logical consequence, (i) the ‘mathematical entanglement’ of $\Omega$-logical validity does not undermine its status as a relation of pure logic; and (ii) both the modal profile and model-theoretic characterization of $\Omega$-logical consequence provide a guide to its epistemic tractability. I argue, then, that there are several considerations adducing in favor of the claim that the interpretation of the concept of set constitutively involves modal notions. The role of the category of modal coalgebraic deterministic automata in (i) characterizing the modal profile of $\Omega$-logical consequence, and (ii) being constitutive of the formal understanding-conditions for the concept of set, provides, then, support for a realist conception of the cumulative hierarchy.

9.2.1 Neo-Logicism

Frege’s (1884/1980; 1893/2013) proposal – that cardinal numbers can be explained by specifying a biconditional between the identity of, and an equivalence relation on, concepts, expressible in the signature of second-order logic – is the first attempt to provide a foundation for mathematics on the basis of logical axioms rather than rational or empirical intuition. In Frege (1884/1980. cit.: 68) and Wright (1983: 104-105), the number of the concept, $A$, is argued to be identical to the number of the concept, $B$, if and only if there is a one-to-one correspondence between $A$ and $B$, i.e., there is a bijective mapping, $R$, from $A$ to $B$. With $N_x$: a numerical term-forming operator,

\[\text{The phrase, } ^*\text{‘mathematical entanglement’}^*_\text{, is owing to Koellner (2010: 2).}\]
\[\forall A \forall B \exists R[Nx : A = Nx : B \equiv \exists y (By \land Rx y \land \forall z (Bz \land Rx z \rightarrow y = z))] \land \forall y [By \rightarrow \exists x (Ax \land Rx y \land \forall z (Az \land Rzy \rightarrow x = z))].\]

Frege’s Theorem states that the Dedekind-Peano axioms for the language of arithmetic can be derived from the foregoing abstraction principle, as augmented to the signature of second-order logic and identity.\(^9\) Thus, if second-order logic may be counted as pure logic, despite that domains of second-order models are definable via power set operations, then one aspect of the philosophical significance of the abstractionist program consists in its provision of a foundation for classical mathematics on the basis of pure logic as augmented with non-logical implicit definitions expressed by abstraction principles.

There are at least three reasons for which a logic defined in ZFC might not undermine the status of its consequence relation as being logical. The first reason for which the mathematical entanglement of \(\Omega\)-logical validity might be innocuous is that, as Shapiro (1991: 5.1.4) notes, many mathematical properties cannot be defined within first-order logic, and instead require the expressive resources of second-order logic. For example, the notion of well-foundedness cannot be expressed in a first-order framework, as evinced by considerations of compactness. Let \(E\) be a binary relation. Let \(m\) be a well-founded model, if there is no infinite sequence, \(a_0, \ldots, a_i\), such that \(Ea_0, \ldots, Ea_{i+1}\) are all true. If \(m\) is well-founded, then there are no infinite-descending \(E\)-chains. Suppose that \(T\) is a first-order theory containing \(m\), and that, for all natural numbers, \(n\), there is a \(T\) with \(n + 1\) elements, \(a_0, \ldots, a_n\), such that \(\langle a_0, a_1\rangle, \ldots, \langle a_n, a_{n-1}\rangle\) are in the extension of \(E\). By compactness, there is an infinite sequence such that that \(a_0 \ldots a_i\), s.t. \(Ea_0, \ldots, Ea_{i+1}\) are all true. So, \(m\) is not well-founded.

By contrast, however, well-foundedness can be expressed in a second-order framework:

\[\forall X[\exists x X x \rightarrow \exists x [X x \land \forall y (X y \rightarrow \neg Ey x)]],\]

such that \(m\) is well-founded iff every non-empty subset \(X\) has an element \(x\), s.t. nothing in \(X\) bears \(E\) to \(x\).

One aspect of the philosophical significance of well-foundedness is that it provides a distinctively second-order constraint on when the membership relation in a given model is intended. This contrasts with Putnam’s (1980)

claim, that first-order models \(mod\) can be intended, if every set \(s\) of reals in \(mod\) is such that an \(\omega\)-model in \(mod\) contains \(s\) and is constructible, such that – given the Downward Lowenheim-Skolem theorem\(^{10}\) – if \(mod\) is non-constructible but has a submodel satisfying \(\forall \text{ is constructible}\), then the model is non-well-founded and yet must be intended. The claim depends on the assumption that general understanding-conditions and conditions on intendedness must be co-extensive, to which I will return in Section 4.2.

A second reason for which \(\Omega\)-logic’s mathematical entanglement might not be pernicious, such that the consequence relation specified in the \(\Omega\)-logic might be genuinely logical, may again be appreciated by its comparison with second-order logic. Shapiro (1998) defines the model-theoretic characterization of logical consequence as follows:

'\(\Phi\) is a logical consequence of \([\text{a model}]\ \Gamma\) if \(\Phi\) holds in all possibilities under every interpretation of the nonlogical terminology which holds in \(\Gamma\)' (148).

A condition on the foregoing is referred to as the ‘isomorphism property’, according to which if two models \(M, M'\) are isomorphic vis-a-vis the nonlogical items in a formula \(\Phi\), then \(M\) satisfies \(\Phi\) if and only if \(M'\) satisfies \(\Phi\)' (151).

Shapiro argues, then, that the consequence relation specified using second-order resources is logical, because of its modal and epistemic profiles. The epistemic tractability of second-order validity consists in ‘typical soundness theorems, where one shows that a given deductive system is ‘truth-preserving’ (154). He writes that: ‘[I]f we know that a model is a good mathematical model of logical consequence (10), then we know that we won’t go wrong using a sound deductive system. Also, we can know that an argument is a logical consequence ... via a set-theoretic proof in the metatheory’ (154-155).

The modal profile of second-order validity provides a second means of accounting for the property’s epistemic tractability. Shapiro argues, e.g., that: ‘If the isomorphism property holds, then in evaluating sentences and arguments, the only ‘possibility’ we need to ‘vary’ is the size of the universe. If enough sizes are represented in the universe of models, then the modal nature of logical consequence will be registered ... [T]he only ‘modality’ we keep is ‘possible size’, which is relegated to the set-theoretic metatheory’

\(^{10}\)For any first-order model \(M, M\) has a submodel \(M'\) whose domain is at most denumerably infinite, s.t. for all assignments \(s\) on, and formulas \(\phi(x)\) in, \(M', M,s \models \phi(x) \iff M',s \models \phi(x)\).
Shapiro's remarks about the considerations adducing in favor of the logicality of non-effective, second-order validity generalize to Ω-logical validity. In the previous section, the modal profile of Ω-logical validity was codified by the duality between the category, $\mathcal{A}$, of coalgebraic modal logics and complete Boolean-valued algebraic models of Ω-logic. As with Shapiro's definition of logical consequence, where $\Phi$ holds in all possibilities in the universe of models and the possibilities concern the 'possible size' in the set-theoretic metatheory, the Ω-Conjecture states that $V \models_\Omega \phi$ iff $V^B \models_\Omega \phi$, such that Ω-logical validity is invariant in all set-forcing extensions of ground models in the set-theoretic multiverse.

Finally, the epistemic tractability of Ω-logical validity is secured, both – as on Shapiro’s account of second-order logical consequence – by its soundness, but also by its being the dual of coalgebraic category of deterministic automata, where the determinacy thereof is again secured by the existence of Woodin cardinals.

9.2.2 Set-theoretic Realism

In this section, I argue, finally, that the modal profile of Ω-logic can be availed of in order to account for the understanding-conditions of the concept of set, and thus crucially serve as part of the argument for set-theoretic realism.

Putnam (op. cit.: 473-474) argues that defining models of first-order theories is sufficient for both understanding and specifying an intended interpretation of the latter. Wright (1985: 124-125) argues, by contrast, that understanding-conditions for mathematical concepts cannot be exhausted by the axioms for the theories thereof, even on the intended interpretations of the theories. He suggests, e.g., that:

'If there really were uncountable sets, their existence would surely have to flow from the concept of set, as intuitively satisfactorily explained. Here, there is, as it seems to me, no assumption that the content of the ZF-axioms cannot exceed what is invariant under all their classical models. [Benacerraf] writes, e.g., that: 'It is granted that they are to have their 'intended interpretation': 'e' is to mean set-membership. Even so, and conceived as encoding the intuitive concept of set, they fail to entail the existence of uncountable sets. So how can it be true that there are such sets? Benacerraf’s reply is that the ZF-axioms are indeed faithful to the relevant informal notions only if, in addition to ensuring that 'E' means set-membership, we inter-
pret them so as to observe the constraint that 'the universal quantifier has to mean all or at least all sets' (p. 103). It follows, of course, that if the concept of set does determine a background against which Cantor’s theorem, under its intended interpretation, is sound, there is more to the concept of set that can be explained by communication of the intended sense of ‘e’ and the stipulation that the ZF-axioms are to hold. And the residue is contained, presumably, in the informal explanations to which, Benacerraf reminds us, Zermelo intended his formalization to answer. At least, this must be so if the ‘intuitive concept of set’ is capable of being explained at all. Yet it is notable that Benacerraf nowhere ventures to supply the missing informal explanation – the story which will pack enough into the extension of ‘all sets’ to yield Cantor’s theorem, under its intended interpretation, as a highly non-trivial corollary’ (op. cit).

In order to provide the foregoing explanation in virtue of which the concept of set can be shown to be associated with a realistic notion of the cumulative hierarchy, I will argue that there are several points in the model theory and epistemology of set-theoretic languages at which the interpretation of the concept of set constitutively involves modal notions. The aim of the section will thus be to provide a modal foundation for mathematical platonism.

One point is in the coding of the signature of the theory, T, in which Gödel’s incompleteness theorems are proved (cf. Halbach and Visser, 2014). Relative to,

(i) a choice of coding for an ω-complete, recursively axiomatizable language, L, of T – i.e. a mapping between properties of numbers and properties of terms and formulas in L;
(ii) a predicate, phi; and
(iii) a fixed-point construction:

Let phi express the property of ‘being provable’, and define (iii) such that, for all consistent theories T of L, there are sentences, p_{phi}, corresponding to each formula, phi(x), in T, s.t. for ‘m’ := p_{phi},

\[ \vdash_T p_{phi} \iff \phi(m) \]

One can then construct a sentence, ‘m’ := ¬phi(m), such that L is incomplete (the first incompleteness theorem).

Moreover, L cannot prove its own consistency:

If:

\[ \vdash_T 'm' \iff \neg\phi(m) \]

Then:
Thus, $L$ is consistent only if $L$ is inconsistent (the second incompleteness theorem).

In the foregoing, the choice of coding bridges the numerals in the language with the properties of the target numbers. The choice of coding is therefore intensional, and has been marshalled in order to argue that the very notion of syntactic computability – via the equivalence class of partial recursive functions, $\lambda$-definable terms, and the transition functions of discrete-state automata such as Turing machines – is constitutively semantic (cf. Rescorla, 2015). Further points at which intensionality can be witnessed in the phenomenon of self-reference in arithmetic are introduced by Reinhardt (1986). Reinhardt (op. cit.: 470-472) argues that the provability predicate can be defined relative to the minds of particular agents – similarly to Quine’s (1968) and Lewis’ (1979) suggestion that possible worlds can be centered by defining them relative to parameters ranging over tuples of spacetime coordinates or agents and locations – and that a theoretical identity statement can be established for the concept of the foregoing minds and the concept of a computable system.

In the previous section, intensional computational properties were defined via modal coalgebraic deterministic automata, where the coalgebraic categories are dual to the category of sets in which $\Omega$-logical validity was defined. Coalgebraic modal logic was shown to elucidate the modal profile of $\Omega$-logical consequence in the Boolean-valued algebraic models of set theory. The intensionality witnessed by the choice of coding may therefore be further witnessed by the modal automata specified in the foregoing coalgebraic logic.

A second point at which understanding-conditions may be shown to be constitutively modal can be witnessed by the conditions on the epistemic entitlement to assume that the language in which Gödel’s second incompleteness theorem is proved is consistent (cf. Dummett, 1963/1978; Wright, 1985). Wright (op. cit.: 91, fn.9) suggests that ‘[T]o treat [a] proof as establishing consistency is implicitly to exclude any doubt ... about the consistency of first-order number theory’. Wright’s elaboration of the notion of epistemic entitlement, appeals to a notion of rational ‘trust’, which he argues is recorded by the calculation of ‘expected epistemic utility’ in the setting of decision theory (2004; 2014: 226, 241). Wright notes that the rational trust subserving epistemic entitlement will be pragmatic, and makes the intriguing point that ‘pragmatic reasons are not a special genre of reason, to be contrasted with e.g. epistemic, prudential, and moral reasons’ (2012: 484). Crucially,
however, the very idea of expected epistemic utility in the setting of decision theory makes implicit appeal to the notion of possible worlds, where the latter can again be determined by the coalgebraic logic for modal automata.

A third consideration adducing in favor of the thought that grasp of the concept of set might constitutively possess a modal profile is that the concept can be defined as an intension – i.e., a function from possible worlds to extensions. The modal similarity types in the coalgebraic modal logic may then be interpreted as dynamic-interpretational modalities, where the dynamic-interpretational modal operator has been argued to entrain the possible reinterpretations both of the domains of the theory’s quantifiers (cf. Fine, 2005, 2006), as well as of the intensions of non-logical concepts, such as the membership relation (cf. Uzquiano, 2015).

The fourth consideration avails directly of the modal profile of $\Omega$-logical consequence. While the above dynamic-interpretational modality will suffice for possible reinterpretations of mathematical terms, the absoluteness and generic invariance of the consequence relation is such that, if the $\Omega$-conjecture is true, then $\Omega$-logical validity is invariant in all possible set-forcing extensions of ground models in the set-theoretic multiverse. The truth of the $\Omega$-conjecture would thereby place an indefeasible necessary condition on a formal understanding of the intension for the concept of set.

### 9.3 Concluding Remarks

In this essay I have examined the philosophical significance of the duality between modal coalgebraic models of automata and Boolean-valued algebraic

---

11For an examination of the philosophical significance of modal coalgebraic automata beyond the philosophy of mathematics, see Baltag (2003). Baltag (op. cit.) proffers a colagebraic semantics for dynamic-epistemic logic, where coalgebraic functors are intended to record the informational dynamics of single- and multi-agent systems. For an algebraic characterization of dynamic-epistemic logic, see Kurz and Palmigiano (2013). The latter proceeds by examining undecidable sentences via the epistemic interpretation of multi-dimensional intensional semantics. See Reinhardt (1974), for a similar epistemic interpretation of set-theoretic languages, in order to examine the reduction of the incompleteness of undecidable sentences on the counterfactual supposition that the language is augmented by stronger axioms of infinity; and Maddy (1988,b), for critical discussion. Chihara (2004) argues, as well, that conceptual possibilities can be treated as imaginary situations with regard to the construction of open-sentence tokens, where the latter can then be availed of in order to define nominalistically adequate arithmetic properties.
models of modal $\Omega$-logic. I argued that – as with the property of validity in second-order logic – $\Omega$-logical validity is genuinely logical, and thus entails a type of neo-logicism in the foundations of mathematics. I argued, then, that modal coalegebraic deterministic automata, which characterize the modal profile of $\Omega$-logical consequence, are constitutive of the interpretation of mathematical concepts such as the membership relation. The philosophical significance of modal $\Omega$-logic is thus that it can be availed of to vindicate both a neo-logicist foundation for set theory and a realist interpretation of the cumulative hierarchy of sets.
Chapter 10

Epistemic Modality and Absolute Decidability

10.1 Introduction

This essay aims to contribute to the analysis of the nature of mathematical modality, and to the applications of the latter to unrestricted quantification and absolute decidability. I argue that mathematical modality falls under at least three types; the interpretational, the metaphysical, and the logical. The interpretational type of mathematical modality has traditionally been taken to concern the interpretation of the quantifiers (cf. Linnebo, 2009, 2010, 2013; Studd, 2013); the possible reinterpretations of the intensions of the concept of set (Uzquiano, 2015,a); and the possibility of reinterpreting the domain over which the quantifiers range, in order to avoid inconsistency (cf. Fine, 2006, 2007). The metaphysical type of modality concerns the ontological profile of abstracta and mathematical truth. Abstracta are thus argued to have metaphysically necessary being, and mathematical truths hold of metaphysical necessity, if at all (cf. Fine, 1981). Instances, finally, of the logical type of mathematical modality might concern the properties of consistency (cf. Field, 1989: 249-250, 257-260; Rayo, 2013: 50; Leng: 2007; 2010: 258), and can perhaps be further witnessed by the logic of provability (cf. Boolos, 1993) and the modal profile of forcing (cf. Kripke 1965; Hamkins and Löwe, 2008).

The significance of the present contribution is as follows. (i) Rather than countenancing the interpretational type of mathematical modality as a prim-
itive, I argue that the interpretational type of mathematical modality is a species of epistemic modality.\(^1\) (ii) I argue, then, that the framework of two-dimensional semantics ought to be applied to the mathematical setting. The framework permits of a formally precise account of the priority and relation between epistemic mathematical modality and metaphysical mathematical modality. I target, in particular, the modal axioms that the respective interpretations of the modal operator ought to satisfy. The discrepancy between the modal systems governing the parameters in the two-dimensional intensional setting provides an explanation of the difference between the metaphysical possibility of absolute decidability and our knowledge thereof. (iii) Finally, I examine the application of the mathematical modalities beyond the issues of unrestricted quantification and indefinite extensibility. As a test case for the two-dimensional approach, I investigate the interaction between the epistemic and metaphysical mathematical modalities and large cardinal axioms. The two-dimensional framework permits of a formally precise means of demonstrating how the metaphysical possibility of absolute decidability and the continuum hypothesis can be accessed by their epistemic-modal-mathematical profile. The logical mathematical modalities – of consistency, provability, and forcing – provide the means for discerning whether mathematical truths are themselves epistemically possible. I argue that, in the absence of disproof, large cardinal axioms are epistemically possible, and thereby provide a sufficient guide to the metaphysical mathematical possibility of determinacy claims and the continuum hypothesis.

In Section 2, I define the formal clauses and modal axioms governing the epistemic and metaphysical types of mathematical modality. In Section 3, I discuss how the properties of the epistemic mathematical modality and metaphysical mathematical modality converge and depart from previous at-

\(^1\)A precedent to the current approach is Parsons (1979-1980; 1983: p. 25, chs.10-11; 2008: 176), who argues that intuition is both a species of the imagination and can be formalized by a mathematical modality. The mathematical modality is governed by S4.2, and concerns possible iterations of the successor operation in arithmetic and possible extensions of the set-theoretic cumulative hierarchy. Among the differences between Parsons’ approach and the one here outlined is (i) that, by contrast to the current proposal, Parsons notes that his notion of mathematical modality is not epistemic (2008: 81fn1); and (ii) that Parsons (1997: 348-351; 2008: 98-100) suggests that the intuitional mathematical modality concerning computable functions is an idealization insensitive to distinctions such as those captured by computational complexity theory, rather than being defined relative to an epistemic modal space comprising the computational theory of mind. (See chapters 2-3, for further discussion.)
tempts to delineate the contours of similar notions. Section 4 extends the
two-dimensional intensional framework to the issue of mathematical knowl-
edge; in particular, to the modal profile of large cardinal axioms and to the
absolute decidability of the continuum hypothesis. Section 5 provides con-
cluding remarks.

10.2 Mathematical Modality

10.2.1 Metaphysical Mathematical Modality

A formula is a logical truth if and only if the formula is true in an intended
model structure, $M = <W, D, R, V>$, where $W$ designates a space of meta-
physically possible worlds; $D$ designates a domain of entities, constant across
worlds; $R$ designates an accessibility relation on worlds; and $V$ is an assign-
ment function mapping elements in $D$ to subsets of $W$.

**Metaphysical Mathematical Possibility**

\[ [\Diamond \phi]^{v,w} = 1 \iff \exists w' [\phi]^{v,w'} = 1 \]

**Metaphysical Mathematical Necessity**

\[ [\Box \phi]^{v,w} = 1 \iff \forall w' [\phi]^{v,w'} = 1, \]

with $\Diamond := \neg \Box \neg$.

10.2.2 Epistemic Mathematical Modality

In order to accommodate the notion of epistemic possibility, we enrich $M$
with the following conditions: $M = <C, W, D, R, V>$, where $C$, a set of
epistemically possibilities, is constrained as follows:

Let $[\phi]^c \subseteq C$;

($\phi$ is a formula encoding a state of information at an epistemically possible
world).

**Intensions**

- $\text{pri}(x) = \lambda c. [x]^{c,c}$;
  (the two parameters relative to which $x$ – a propositional variable – ob-
tains its value are epistemically possible worlds).
- $\text{sec}(x) = \lambda c. [x]^{u,u}$
  (the two parameters relative to which $x$ obtains its value are metaphysi-
cally possible worlds).
Then:

- **Epistemic Mathematical Necessity**
  \[\llbracket \Box \phi \rrbracket^{c,w} = 1 \iff \forall c' [\llbracket \phi \rrbracket^{c,c'} = 1]\]
  \((\phi \text{ is true at all points in epistemic modal space}).\)

- **Epistemic Mathematical Possibility**
  \[\llbracket \Diamond \phi \rrbracket \neq \emptyset \iff \llbracket \neg \Box \neg \phi \rrbracket = 1\]
  \((\phi \text{ might be true if and only if it is not epistemically necessary for } \phi \text{ to be false}).\)

Crucially, epistemic mathematical modality is constrained by consistency, and the formal techniques of provability and forcing. A mathematical formula is false, and therefore metaphysically impossible, if it can be disproved or induces inconsistency in a model.

### 10.2.3 Interaction

- **Convergence**
  \[\forall c \exists w [\llbracket \phi \rrbracket^{c,w} = 1]\]
  \((\text{the value of } x \text{ is relative to a parameter for the space of epistemically possible worlds. The value of } x \text{ relative to the first parameter determines the value of } x \text{ relative to the second parameter for the space of metaphysical possibility}).\)

- **Super-rigidity**
  \[\llbracket \phi \rrbracket^{c,w} = 1 \iff \forall w', c' [\llbracket \phi \rrbracket^{c',w'} = 1]\]
  \((\text{the intension of } \phi \text{ is rigid in all points in epistemic and metaphysical modal space}).\)

### 10.2.4 Modal Axioms

- **Metaphysical mathematical modality is governed by the modal system KTE, as augmented by the Barcan formula and its Converse (cf. Fine, 1981).**
\( K: \Box[\phi \rightarrow \psi] \rightarrow [\Box\phi \rightarrow \Box\psi] \)

\( T: \Box\phi \rightarrow \phi \)

\( E: \neg\Box\phi \rightarrow \Box\neg\Box\phi \)

Barcan: \( \diamond\exists xFx \rightarrow \exists x\diamond Fx \)

Converse Barcan: \( \exists x\diamond Fx \rightarrow \diamond\exists xFx \)

- Epistemic mathematical modality is governed by the modal system, \( KT4 \), as augmented by the Barcan formula and the Converse Barcan formula.\(^2\)

\( K: \square[\phi \rightarrow \psi] \rightarrow [\square\phi \rightarrow \square\psi] \)

\( T: \square\phi \rightarrow \phi \)

\( 4: \square\phi \rightarrow \square\square\phi \)

Barcan: \( \Diamond\exists xFx \rightarrow \exists x\Diamond Fx \)

Converse Barcan: \( \exists x\Diamond Fx \rightarrow \Diamond\exists xFx \)

Note that, if one prefers a hyperintensional semantics to an intensional semantics, one can avail of the definitions of hyperintensions as functions from states in a state space to extensions instead of from whole epistemically and metaphysically possible worlds. See chapters 4 and 8 for the relevant models and definitions.

### 10.3 Departures from Precedent

The approach to mathematical modality, according to which it yields a representation of the cumulative universe of sets, has been examined by Fine

\(^2\)Reasons adducing against including the Smiley-Gödel-Löb provability formula among the axioms of epistemic mathematical modality are examined in Section 5. GL states that \( \square\square\phi \rightarrow \phi \rightarrow \square\phi \). For further discussion of the properties of GL, see Löb (1955); Smiley (1963); Kripke (1965); and Boolos (1993). Löb’s provability formula was formulated in response to Henkin’s (1952) problem concerning whether a sentence which ascribes the property of being provable to itself is provable. (Cf. Halbach and Visser, 2014, for further discussion.) For an anticipation of the provability formula, see Wittgenstein (1933-1937/2005: 378). Wittgenstein writes: ‘If we prove that a problem can be solved, the concept ‘solution’ must somehow occur in the proof. (There must be something in the mechanism of the proof that corresponds to this concept.) But the concept mustn’t be represented by an external description; it must really be demonstrated. / The proof of the provability of a proposition is the proof of the proposition itself’ (op. cit.). Wittgenstein contrasts the foregoing type of proof with ‘proofs of relevance’ which are akin to the mathematical, rather than empirical, propositions, discussed in Wittgenstein (2001: IV, 4-13, 30-31).
(2006) and Uzquiano (op. cit.). Fine argues that the mathematical modality should be interpretational; and thus taken to concern the reinterpretation of the domain over which the quantifiers range, in order to avoid inconsistency. Uzquiano argues similarly for an interpretational construal of mathematical modality, where the cumulative hierarchy of sets is fixed, yet what is possibly reinterpreted is the non-logical vocabulary of the language, in particular the membership relation.3

On Fine’s approach, the interpretational modality is both postulational, and ‘prescriptive’ or imperatival. The prescriptive element consists in the rule:

‘Introduction: !x.C(x),’

such that one is enjoined to postulate, i.e. to ‘introduce an object x conforming to the condition C(x)’ (2005: 91; 2006: 38).

In the setting of unrestricted quantification, suppose, e.g., that there is an interpretation for the domain over which a quantifier ranges. Fine writes that an interpretation ‘I is exten[s]ible – in symbols, E(I) – if possibly some interpretation extends it, i.e. ◻∃J(I ⊂ J)’ (2006: 30). Then, the interpretation of the domain over which the quantifier ranges is extensible, if ‘∀I.E(I)’. The interpretation of the domain over which the quantifier ranges is indefinitely extensible, if ‘◻∀I.E(I)’ iff ‘◻∀I ◻∃J(I ⊂ J)’, where the reinterpretation is induced via the prescriptive imperative to postulate the existence of a new object by the foregoing ‘Introduction’ rule (2006: 30-31; 38). Fine clarifies that the interpretational approach is consistent with a ‘realist ontology’ of the set of reals. He refers to the imperative to postulate new objects, and thereby reinterpret the domain for the quantifier, as the ‘mechanism’ by which epistemically to track the cumulative hierarchy of sets (2007: 124-125).

In accord with Fine’s approach, the epistemic mathematical modality defined in the previous section was taken to have a similarly representational interpretation, and perhaps the postulational property is an optimal means of inducing a reinterpretation of the domain of the quantifier. However, the present approach avoids a potential issue with Fine’s account, with regard to the the introduction of deontic modal properties of the prescriptive and imperatival rules that he mentions. It is sufficient that the interpretational modalities are a species of epistemic modality, i.e. possibilities that are relative to agents’ spaces of states of information.

Developing Parsons’ (1983) program, Linnebo (2013) outlines a modalized

3Compare Gödel, 1947; Williamson, 1998; and Fine, 2005c.
version of ZF. Similarly to the modal axioms for the epistemic mathematical modality specified in the previous section, Linnebo argues that his modal set theory ought to be governed by the system S4.2, the Converse Barcan formula, and (at least a restricted version of) the Barcan formula. However – rather than being either interpretational or epistemic – Linnebo deploys the mathematical modality in order to account for the notion of ‘potential infinity’, as anticipated by Aristotle. The mathematical modality is thereby intended to provide a formally precise answer to the inquiry into the extent of the cumulative set-theoretic hierarchy; i.e., in order to precisify the answer that the hierarchy extends ‘as far as possible’ (2013: 205).

Thus, Linnebo takes the modality to be constitutive of the actual ontology of sets; and the quantifiers ranging over the actual ontology of sets are claimed to have an ‘implicitly modal’ profile (2010: 146; 2013: 225). He suggests, e.g., that: ’As science progresses, we formulate set theories that characterize larger and larger initial segments of the universe of sets. At any one time, precisely those sets are actual whose existence follows from our strongest, well-established set theory’ (2010: 159n21). However – despite his claim that the modality is constitutive of the actual ontology of sets – Linnebo concedes that the mathematical modality at issue cannot be interpreted metaphysically, because sets exist of metaphysical necessity if at all (2010: 158; 2013: 207). In order partly to allay the tension, Linnebo re-

---

4Linnebo (2018) discusses the differences between Putnam’s and Parsons’ accounts of the role of modality in mathematics. Berry (forthcoming) also discusses the differences between the foregoing. Linnebo (op. cit.: 265-266) avails of two-dimensional indexing for the relation between interpretational and circumstantial modalities. The appeal to epistemic two-dimensional semantics in order to account for interpretational as epistemic and circumstantial as metaphysical modalities and their interaction in this essay was written in 2015 and pursued prior to knowledge of Linnebo’s account. My approach differs, as well, by countenancing a hyperintensional, epistemic two-dimensional truthmaker semantics and applying it to the epistemology of mathematics, as in chapters 8 and 10.

5Cf. Aristotle, Physics, Book III, Ch. 6.

6Precursors to the view that modal operators can be availed of in order to countenance the potential hierarchy of sets include Hodes (1984b). Intensional constructions of set theory are further developed by Reinhardt (1974); Parsons (op. cit.); Myhill (1985); Scedrov (1985); Flagg (1985); Goodman (1985); Hellman (1990); Nolan (2002); and Studd (2013). (See Shapiro (1985) for an intensional construction of arithmetic.) Chihara (2004: 171-198) argues that 'broadly logical' conceptual possibilities can be used to represent imaginary situations relevant to the construction of open-sentence tokens. The open-sentences can then be used to define the properties of natural and cardinal numbers and the axioms of Peano arithmetic.
marks, then, that set theorists 'do not regard themselves as located at some particular stage of the process of forming sets' (2010: 159); and this might provide evidence that the inquiry – concerning at which stage in the process of set-individuation we happen to be, at present – can be avoided.

Another distinction to note is that both Linnebo (op. cit.) and Uzquiano (op. cit.) avail of second-order plural quantification, in developing their primitivist and interpretational accounts of mathematical modality. By contrast to their approaches, the epistemic and metaphysical modalities defined in the previous section are defined with second-order singular quantification over sets.

Linnebo and Uzquiano both suggest that their mathematical modalities ought to be governed by the G axiom; i.e. $\lozenge \Box \phi \rightarrow \Box \lozenge \phi$. The present approach eschews, however, of the G axiom, in virtue of the following. Williamson (2009) demonstrates that – because KT4G is a sublogic of S5 – an epistemic operator which validates the conjunction of the 4 axiom of positive introspection and the E axiom of negative introspection will be inconsistent with the condition of 'recursively enumerable conservativeness' (30). 'If a modal logic is r.e. (quasi-)conservative then every (consistent) r.e. theory in the language without $\Box$ [interpreted as "I know that..."] is conservatively extended by an r.e. theory in the language with $\Box$ such that it is consistent in the modal logic for a recursively enumerable theory$R$ to be exactly what the agent cognizes in the language without $\Box$ while what the agent cognizes in the language with $\Box$ constitutes an r.e. theory' (12). As axioms of an agent’s consistent, recursively axiomatizable theorizing about the theory of its own states of knowledge and belief, the conjunction of 4 and E would entail that the agent’s theory is both consistent and decidable, in conflict with Gödel’s (1931) second incompleteness theorem. The modal system, KT4, avoids the foregoing result. In the present setting, the circumvention is innocuous, because the undecidability – yet recursively enumerable quasi-conservativeness – of an epistemic agent’s consistent theorizing about its epistemic states is consistent with the epistemic mathematical possibility that large cardinal axioms are absolutely decidable.

Finally, my application of epistemic two-dimensional semantics to the epistemology of mathematics departs from full-blooded platonism, as well. According to full-blooded platonism, whatever mathematical objects can exist, do exist, and every consistent mathematical theory describes either a different part of the mathematical universe or distinct mathematical universes altogether (Balaguer, 1998). Thus, ZFC+CH and ZFC+¬CH both *truly de-
scribe collections of mathematical objects", holding in distinct albeit equally real mathematical universes (Balaguer, 2001: 97: see also Hamkins, 2012).

Epistemic two-dimensionalism and full-blooded platonism differ on both the nature of their target possibilities and on the status of the actuality of the possibilities. Epistemic two-dimensionalism avails of epistemic possibilities, whereas full-blooded platonism avails of logical possibilities. Further, not all epistemic possibilities are actual according to epistemic two-dimensionalism, whereas the objects of any logically consistent theory actually exist according to full-blooded platonism. One reason to prefer epistemic two-dimensionalism to full-blooded platonism is that the former can be formalized, whereas Restall (2003) has shown that there are significant challenges to formalizing the latter. Another reason to prefer epistemic two-dimensionalism is that – unlike full-blooded platonism – it avoids commitment to the existence of inconsistent universes of sets where e.g. both ZFC+V=L and ZFC+V≠L would obtain.

10.4 Knowledge of Absolute Decidability

Williamson (2016) examines the extension of the metaphysically modal profile of mathematical truths to the question of absolute decidability. A statement is decidable if and only if there is a mechanical procedure for deciding it or its negation. Statements are absolutely undecidable if and only if they are "undecidable relative to any set of axioms that are justified" rather than just relative to a system (Koellner, 2006: 153), and they are absolute decidable if and only if they are not absolutely undecidable. In this section, I aim to extend Williamson’s analysis to the notion of epistemic mathematical modality that has been developed in the foregoing sections. The extension provides a crucial means of witnessing the significance of the two-dimensional intensional approach for the epistemology of mathematics.

Williamson proceeds by suggesting the following line of thought. Suppose that A is a true interpreted mathematical formula which eludes present human techniques of provability; e.g. the continuum hypothesis (op. cit.). Williamson argues that mathematical truths are metaphysically necessary (op. cit.). Williamson then enjoins one to consider the following scenario: It is metaphysically possible that there is a species which can prove that A. Therefore, A is absolutely provable; that is, A 'can in principle be known by a normal mathematical process' such as derivation in an axiomatizable formal
system with quantification and identity. He proposes a safety condition on knowledge. He writes: "In current epistemological terms, their knowledge of A meets the condition of safety: they could not easily have been wrong in a relevantly similar case. Here the relevantly similar cases include cases in which the creatures are presented with sentences that are similar to, but still discriminably different from, A, and express different and false propositions; by hypothesis, the creatures refuse to accept such other sentences, although they may also refuse to accept their negations" (11). Williamson writes then that: 'The claim is not just that A would be absolutely provable if there were such creatures. The point is the stronger one that A is absolutely provable because there could in principle be such creatures.'

Williamson’s scenario evinces one issue for the 'back-tracking' approach to modal epistemology, at least as it might be applied to the issue of possible mathematical knowledge. On the back-tracking approach, the method of modal epistemology is taken to proceed by first discerning the metaphysical modal truths – normally by natural-scientific means – and then working backward to the exigent incompleteness of an individual’s epistemic states concerning such truths (cf. Stalnaker, 2003; Vetter, 2013).

The issue for the back-tracking method that Williamson’s scenario illuminates is that the metaphysical mathematical possibility that CH is absolutely decidable must in some way converge with the epistemic possibility thereof. The normal mathematical techniques that Williamson specifies – i.e. proof and forcing – have both an epistemic and a metaphysical dimension. Thus, whether CH is metaphysically necessary – and thus, as Williamson claims, metaphysically possible and absolutely decidable thereby – can only be witnessed by the epistemic means of demonstrating that its absolute decidability is not impossible. Nevertheless, the epistemic mathematical possibility of the decidability of CH is a guide to its metaphysical mathematical possibility.

The significance of the two-dimensional intensional framework outlined in the foregoing is that it provides an explanation of the discrepancy between metaphysical mathematical modality and epistemic mathematical modality. Metaphysical mathematical modality is governed by the system S5, the Barcan formula, and its Converse, whereas epistemic mathematical modality is governed by KT4, the Barcan formula, and its Converse. Thus, epistemic mathematical modality figures as the mechanism, which enables the tracking of metaphysically possible mathematical truth.7

7A provisional definition of large cardinal axioms is as follows.
Leitgeb (2009) endeavors similarly to argue for the convergence between the notion of informal provability – countenanced as an epistemic modal operator, K – and mathematical truth. Availing of Hilbert’s (1923/1996: ¶18-42) epsilon terms for propositions, such that, for an arbitrary predicate, C(x), with x a propositional variable, the term ‘εp.C(p)’ is intuitively interpreted as stating that ‘there is a proposition, x(/p), s.t. the formula, that p satisfies C, obtains’ (op. cit.: 290). Leitgeb purports to demonstrate that ∀p(p → Kp), i.e. that informal provability is absolute; i.e. truth and provability are co-extensive. He argues as follows. Let A(p) abbreviate the formula ‘p ∧ ¬K(p)’, i.e., that the proposition, p, is true while yet being unprovable. Let K be the informal provability operator reflecting knowability or epistemic necessity, with ⟨K⟩ its dual. Then:


By necessitation,


Applying modal axioms, KT, to (1), however,


Thus,


Leitgeb suggests that (4) be rewritten

5. (K)∀p(p → Kp).

Abbreviate (5) by B. By existential introduction and modal axiom K, both

6. B → ∃p[K(p → B) ∨ K(p → ¬B) ∧ p], and

7. ¬B → ∃p[K(p → B) ∨ K(p → ¬B) ∧ p].

Thus,

8. ∃p[K(p → B) ∨ K(p → ¬B) ∧ p].

Abbreviate (8) by C(p). Introducing epsilon notation,


∃xΦ is a large cardinal axiom, because:
(i) Φx is a Σ₂-formula;
(ii) if κ is a cardinal, such that V |= Φ(κ), then κ is strongly inaccessible, where a cardinal κ is regular if the cofinality of κ – comprised of the unions of sets with cardinality less than κ – is identical to κ, and a strongly inaccessible cardinal is regular and has a strong limit, such that if λ < κ, then 2^λ < κ (Cf. Kanamori, 2012b: 360); and
(iii) for all generic partial orders P ∈ V_κ, and all V-generics G ⊆ P, V[G] |= Φx (Koellner, 2006: 180).

8See Section 5, for further discussion of the duality of knowledge, and its relation to doxastic operators.
By K,
10. \[K(\epsilon p.C(p) \rightarrow KB) \lor K(\epsilon p.C(p) \rightarrow K\neg B)\].

From (9) and necessitation, one can further derive
11. K\(\epsilon p.C(p)\).

By (10) and (11),
12. KB \lor K\neg B.

From (5), (12), and K, Leitgeb derives
13. KB.

By, then, the T axiom,
14. \(\forall p(p \rightarrow Kp)\) (291-292).

Rather than accounting for the coextensiveness of epistemic provability and truth, Leitgeb interprets the foregoing result as cause for pessimism with regard to whether the formulas contended in epistemic logic and via epsilon terms are genuinely logical truths if true at all (292).

In response to the attending pressure on the status of epistemic logic as concerning truths of logic, one can challenge the derivation, in the above proof, from lines (12) to (13). The inference depends on line (5), i.e., the epistemic possibility of completeness: \(\langle K \rangle \forall p(p \rightarrow Kp)\). Assume that line (5) is valid. Then, the validity of the inference from (12) to (13) can be challenged by the restriction on the quantifier on worlds in the Knowability Principle expressed by (5). The epistemic operator in lines (12) and (13) records, by contrast, the epistemic necessity, rather than the possibility, of the truth of the formulas and subformulas therein. Thus, from (12) either the provability of the provability of propositions or the provability of the unprovability of propositions, one cannot derive (13) the provability of the provability of propositions, because – by (5) – it is only epistemically possible that all true propositions are provable.

A final question is whether Orey sentences such as the Continuum Hypothesis (CH) have a determinate epistemic intension given that there are currently models in which CH is true and models in which CH is false, such that it isn’t determinate which epistemic possibility is actual. In response to this worry, the epistemic intension is arguably indeterminate for non-ideal reasoners, yet determinate for ideal ones, such that the epistemic mathematical modality at issue can be divided into non-ideal and ideal varieties.\(^9\)

\(^9\)For the distinction between ideal and prima facie (i.e., non-ideal) conceivable, see Chalmers (2002).
10.5 Concluding Remarks

In this essay, I have endeavored to delineate the types of mathematical modality, and to argue that the epistemic interpretation of two-dimensional semantics can be applied in order to explain, in part, the epistemic status of large cardinal axioms and the decidability of Orey sentences. The formal constraints on mathematical conceivability adumbrated in the foregoing can therefore be considered a guide to our possible knowledge of unknown mathematical truth.
Chapter 11

Grothendieck Universes, and Indefinite Extensibility

This essay endeavors to provide a characterization of the notion of definiteness, in order to provide a non-circular definition of the concept of indefinite extensibility. The concept of indefinite extensibility is introduced by Dummett (1963/1978), in the setting of a discussion of the philosophical significance of Gödel’s (1931) first incompleteness theorem. Gödel’s theorem can be characterized as stating that – relative to a coding defined over the signature of first-order arithmetic, a predicate expressing the property of provability, and a fixed point construction which is non-trivial, such that the formula in which the above predicate figures precludes interpretations such as ‘0=1’ – the formula can be defined as not satisfying the provability predicate. Dummett’s concern is with the conditions on our grasp of the concept of natural number, given that the latter figures in a formula whose truth appears to be satisfied despite the unprovability – and thus non-constructivist profile – thereof (186). His conclusion is that the concept of natural number ‘exhibits a particular variety of inherent vagueness, namely indefinite extensibility’, where a ‘concept is indefinitely extensible if, for any definite characterisation of it, there is a natural extension of this characterisation, which yields a more inclusive concept; this extension will be made according to some general principle for generating such extensions, and, typically, the extended characterisation will be formulated by reference to the previous, unextended, characterisation’ (195-196). Elaborating on the notion of indefinite extensibility, Dummett (1996: 441) redefines the concept as follows: an ‘indefinitely extensible concept is one such that, if we can form a definite conception of
a totality all of whose members fall under the concept, we can, by reference to that totality, characterize a larger totality all of whose members fall under it. Subsequent approaches to the notion have endeavored to provide a more precise elucidation thereof, either by providing an explanation of the property which generalizes to an array of examples in number theory and set theory (cf. Wright and Shapiro, 2006), or by availing of modal notions in order to capture the properties of definiteness and extendability which are constitutive of the concept (cf. Fine, 2006; Linnebo, 2013; Uzquiano, 2015). However, the foregoing modal characterizations of indefinite extensibility have similarly been restricted to set-theoretic languages. Furthermore, the modal notions that the approaches avail of are taken to belong to a proprietary type which is irreducible to either the metaphysical or the logical interpretations of the operator.

The aim of this essay is to redress the foregoing, by providing a modal characterization of indefinite extensibility in the setting of category theory, rather than number or set theory. One virtue of the category-theoretic, modal definition of indefinite extensibility is that it provides for a robust account of the epistemological foundations of modal-structuralist approaches to the ontology of mathematics. A second aspect of the philosophical significance of the examination is that it can serve to redress the lacuna noted in the appeal to an irreducible type of mathematical modality, which is argued (i) to be representational, (ii) still to bear on the ontological expansion of domains of sets, and yet (iii) not to range over metaphysical possibilities. By contrast to the latter approach, the category-theoretic characterization of indefinite extensibility is able to identify the functors of coalgebraic nondeterministic automata with elementary embeddings and the modal properties of set-theoretic, Ω-logical consequence. The functors are interpreted both epistemically and metaphysically, such that the functors receive their mappings relative to two parameters, the first ranging over epistemically possible worlds and the second ranging over metaphysically possible worlds. The functors thus receive their values in an epistemic two-dimensional semantics. The semantics can be either intensional or hyperintensional, as set-out in chapter 8.

In Section 2, I examine the extant approaches to explaining both the property and the understanding-conditions on the concept of indefinite extensibility. In Section 3, I outline the elements of the category theory of sets and define Grothendieck Universes. In Section 4, modal coalgebraic automata are availed of to model Grothendieck Universes, and I define the no-
tion of indefinite extensibility in the category-theoretic setting. I argue that the category-theoretic definition of indefinite extensibility, via Grothendieck Universes as modal coalgebraic automata, yields an explanation of the generative property of indefinite extensibility, as well as of the notion of definiteness which figures in the definition. I argue that the generative property of indefinite extensibility can be captured by identifying Kripke functors of colagebras with elementary embeddings. I argue, then, that the notion of definiteness can be captured by the role of Grothendieck Universes-as-modal coalgebraic automata in characterizing the modal profile of Ω-logical consequence, where the latter accounts for the absoluteness of mathematical truths throughout the set-theoretic multiverse. The category-theoretic definition is shown to circumvent the issues faced by rival attempts to define indefinite extensibility via extensional and intensional notions within the setting of set theory. Section 5 provides concluding remarks.

11.1 Indefinite Extensibility in Set Theory: Modal and Extensional Approaches

Characterizations of indefinite extensibility have so far occurred in the language of set theory, and have availed of both extensional and intensional resources. In an attempt to define the notion of definiteness, Wright and Shapiro (op. cit.) argue, for example, that indefinite extensibility may be intuitively characterized as occurring when there is a function which falls under a first-order concept; for a sub-concept of the first-order concept, an application of the function on the sub-concept does not fall within that sub-concept’s range; however, a new sub-concept can be formed, and defined as the set-theoretic union of the initial sub-concept and the function applied thereon (266).

Formally, let Π be a higher-order concept of type τ. Let P be a first-order concept falling under Π of type τ. Let f be a function from entities to entities of the same type as P. Finally, let X be a sub-concept of P. P is indefinitely extensible with respect to Π, if and only if:

\[
\epsilon(P) = f(X), \quad \epsilon(X) = \neg[f(X)], \quad \exists X'[\Pi(X') = (X \cup \{fX\})] \quad (\text{op. cit.}).
\]

The notion of definiteness is then defined as the limitless preservation of
'II-hood' by sub-concepts thereof 'under iteration of the relevant operation', f (269).

The foregoing impresses as a necessary condition on the property of indefinite extensibility. Wright and Shapiro note, e.g., that the above formalization generalizes to an array of concepts countenanced in first-order number theory and analysis, including concepts of the finite ordinals (defined by iterations of the successor function); of countable ordinals (defined by countable order-types of well-orderings); of regular cardinals (defined as occurring when the cofinality of a cardinal, \( \kappa \) – comprised of the unions of sets with cardinality less than \( \kappa \) – is identical to \( \kappa \)); of large cardinals (defined by elementary embeddings from the universe of sets into proper subsets thereof, which specify critical points measured by the ordinals); of real numbers (defined as cuts of sets of rational numbers); and of Gödel numbers (defined as natural numbers of a sequence of recursively enumerable truths of arithmetic) (266-267).

As it stands, however, the definition might not be sufficient for the definition of indefinite extensibility, by being laconic about the reasons for which new sub-concepts – comprised as the union of preceding sub-concepts with a target operation defined thereon – are presumed interminably to generate. In response to the above desideratum, concerning the reasons for which indefinite extensibility might be engendered, philosophers have recently appealed to modal properties of the formation of sets. Fine (2006) argues, e.g., that – in order to avoid the Russell property when quantifying over all sets – there are interpretational modalities which induce a reinterpretation of quantifier domains, and serve as a mechanism for tracking the ontological inflation of the hierarchy of sets via, e.g., the power-set operation (2007). Fine (2005) suggests that the interpretational modality at issue might be a species of dynamic modality, which defines modalities as concerning the information entrained by program executions. Reinhardt (1974) and Williamson (2007) argue that modalities are inter-definable with counterfactuals. While Williamson (2016) argues both that imaginative exercises take the form of counterfactual presuppositions and that it is metaphysically possible to decide propositions which are undecidable relative to the current axioms of extensional mathematical languages such as ZF – Reinhardt (op. cit.) argues that large cardinal axioms and undecidable sentences in extensional ZF can similarly be imagined as obtaining via counterfactual presupposition. In an examination of the iterative hierarchy of sets, Parsons (1977/1983) notes that the notion of potential infinity, as anticipated in Book 3, ch. 6 of Aristotle’s *Physics*, may be codified in a modal set theory by both a principle which
is an instance of the Barcan formula (namely, for predicates P and rigidifying predicates Q, \( \forall x(Px \iff Qx) \land \Box \{ \forall x(\Box Qx \lor \Box \neg Qx) \land \forall R[\forall x(\Box Qx \rightarrow Rx) \rightarrow \Box \forall x(Qx \rightarrow Rx)] \} \) (fn. 24), as well as a principle for definable set-forming operations (e.g., unions) for Borel sets of reals \( \Box (\forall x) \diamond (\exists y)[y = x \cup \{ x \}] \) (528). The modal extension is argued to be a property of the imagination, or intuition, and to apply further to iterations of the successor function in an intensional variant of arithmetic (1979-1980).

Hellman (1990) develops the program intimated in Putnam (1967), and thus argues for an eliminativist, modal approach to mathematical structuralism as applied to second-order plural ZF. The possibilities at issue are taken to be logical – concerning both the consistency of a set of formulas as well as the possible satisfaction of existential formulas – and he specifies, further, an “extendability principle”, according to which “every natural model [of ZF] has a proper extension” (421).

Extending Parsons’ and Fine’s projects, Linnebo (2009, 2013) avails of a second-order, plural modal set theory in order to account for both the notion of potential infinity as well as the notion of definiteness. Similarly to Parsons’ use of the Barcan formula (i.e., \( \forall \Box \phi \rightarrow \Box \forall \phi \)), Linnebo’s principle for the foregoing is as follows: \( \forall u(u \prec xx \rightarrow \Box \phi) \rightarrow \Box \forall u(u \prec \rightarrow \phi) \) (2013: 211). He argues, further, that the logic for the modal operator is S4.2, i.e. K [\( \Box (\phi \rightarrow \psi) \rightarrow (\Box \phi \rightarrow \Box \psi) \)], T (\( \Box \phi \rightarrow \phi \)), 4 (\( (\Box \phi \rightarrow \Box \Box \phi) \)), and G (\( \diamond \Box \phi \rightarrow \Box \diamond \phi \)). Studd (2013) examines the notion of indefinite extensibility by availing of a bimodal temporal logic. Uzquiano’s (2015) approach to defining the concept of indefinite extensibility argues that the height of the cumulative hierarchy is in fact fixed, and that indefinite extensibility can similarly be captured via the use of modal operators in second-order plural modal set theory. The modalities are taken to concern the possible reinterpretations of the intensions of the non-logical vocabulary – e.g., the set-membership relation – which figures in the augmentation of the theory with new axioms and the subsequent climb up the fixed hierarchy of sets (cf. Gödel, 1947/1964).

Chapters 8 and 10 proffer a novel epistemology of mathematics, based on an application of the epistemic interpretation of two-dimensional semantics in set-theoretic languages to the values of large cardinal axioms and undecidable sentences. Modulo logical constraints such as consistency and generic absoluteness in the extensions of ground models of the set-theoretic multiverse, the epistemic possibility that an undecidable proposition receives a value may serve, then, as a guide to the metaphysical possibility thereof. Finally, chapter 9 argues that the modal profile of the consequence relation, in
the $\Omega$-logic defined in Boolean-valued models of set-theory, can be captured by coalgebraic modal automata, and provides a necessary condition on the formal grasp of the concept of ‘set’.

The foregoing accounts of the metaphysics and epistemology of indefinite extensibility are each defined in the languages of number and set theory. In the following section, I examine the nature of indefinite extensibility in the setting of category theory, instead. One aspect of the philosophical significance of the examination is that it can serve to provide an analysis of the mathematical modality at issue, by availing only of model-theoretic resources. By contrast to Hellman’s approach, which takes the mathematical modality at issue to be logical (cf. Field, 1989: 37; Rayo, 2013), and Fine’s (op. cit.) approach, which takes the mathematical modality to be dynamic, I argue in the following sections that the mathematical modality can be captured by the functors of coalgebraic modal automata, where the latter can model Grothendieck Universes, and the functors receive their mappings relative to two parameters, the first ranging over epistemically possible worlds and the second ranging over metaphysically possible worlds.

11.2 Grothendieck Universes

We work within a two-sorted language in which the Eilenberg-Mac Lane Axioms of category theory are specified. Types are labeled $A,B,C$ for objects and $x,y,z$ for arrows. The relevant operators are the domain operator, $\text{Dom}$, which takes arrows to objects; the codomain operator, $\text{Cod}$, which operates similarly, and the identity operator, $1_\alpha$, which takes objects to arrows. Finally, a composition relation, $C(x,y; z)$, is defined on arrows, where the open formula reads $z$ is the composite of $x$ and $y$ (McLarty, 2008: 13). The Eilenberg-Mac Lane axioms can then be defined as follows:

- **Axioms of Domain and Codomain:**
  \[ \forall f,g,h, \text{ if } C(f,g,h), \text{ then } \text{Dom} f = \text{Dom} h \text{ and } \text{Cod} f = \text{Dom} g \text{ and } \text{Cod} g = \text{Cod} h \]

- **Axioms of Existence and Uniqueness of Composites:**
  \[ \forall f,g, \text{ if } \text{Cod} f = \text{Dom} g, \text{ then } \exists! h, \text{ s.t. } C(f,g; h) \]

- **Axioms for Identity Arrows:**
  \[ \forall A, \text{Dom} 1_A = \text{Cod} 1_A = A \]
\( \forall f, C(1_{\text{Dom} f}, f; f) \)
\( \forall f, C(f, 1_{\text{Dom} f}; f) \)

- **Axiom of Associativity of Composition:**
  \( \forall f, g, h, i, j, k, \) if \( C(f, g; i) \) and \( C(g, h; j) \) and \( C(f, j; k) \), then \( C(i, h; k) \) (op. cit.).

Categorical Set Theory is defined by augmenting the Eilenberg-Mac Lane axioms with the axioms of Lawvere’s Elementary Theory for the Category of Sets (ETCS) (op. cit.; Lawvere, 2005). Following McLarty, we define the singleton of a set as one for which ‘every set has exactly one function to it’ (op. cit.: 25). An element of a set \( A, x \in A, \) is a function \( x: 1 \to A \) (op. cit.). Composition occurs if and only if, for two arrows, \( f, g, \) and object \( x, (gf)(x) = g(f(x)) \) (26). Finally, an equalizer \( e: E \to A \) for a pair of functions \( f, g: A \to B \) is defined as ‘a universal solution to the equation \( fe = ge \)’ (29). The axioms are then defined as follows (op. cit.):

- **Every pair of sets,** \( A, B, \) **has a product:**
  \( \forall T, f, g, \) with \( f: T \to A, g: T \to B, \exists !(f, g): T \to A \times B \)

- **Every parallel pair of functions,** \( f, g: A \to B, \) **has an equalizer:**
  \( \forall T, h, \) with \( fh = gh, \exists ! u: T \to E \)

- **There is a function set from each set** \( A \) **to each set** \( B:**
  \( \forall C \) and \( g: C \times A \to B, \exists ! g': C \to B^A \)

- **There is a truth value** \( \text{true} \): \( 1 \to 2:**
  \( \forall A \) and monic \( S \hookrightarrow A, \exists ! \chi_i, \) such that \( S \) is an equalizer

- **There is a natural number triple,** \( N, 0, s:**
  \( \forall T \) and \( x: 1 \to T \) and \( f: T \to T, \exists ! u: x \to T \)

- **Extensionality**
  \( \forall f \neq g: A \to B, \exists x: 1 \to A, \) with \( f(x) \neq g(x) \)

- **Non-triviality**
  \( \exists false: 1 \to 2, \) s.t. \( false \neq true \)
• Choice
\[
\forall \text{ onto functions } f: A \to B, \exists h: B \to A, \text{ s.t. } fh = 1_A.
\]

From the axioms of ETCS, only a version of the ZF separation axiom with bounded quantifiers can be recovered (37). The axiom of separation states that \( \exists x \forall u[ u \in x \iff u \in a \land \phi(u)] \). In order to redress the restriction to bounded quantifiers, we work within a stronger category theory for sets, i.e. the 'category of categories as foundation' (CCAF). The axioms of the CCAF build upon those of both ETCS and Eilenberg-Mac Lane category theory, by augmenting them with the following (53):

• Every category C has a unique functor, \( C \to 1 \)
• The category 2 has exactly two functors from 1 and 3 to itself
• Let a pushout be defined such that if \( f: A \to C \) and \( g: B \to C \), then \( a: C \to A \) and \( b: C \to B \) (Pettigrew, ms: 19). The category 3 is a pushout, and there is a functor \( \gamma: 2 \to 3 \), with \( \gamma_0 = \alpha_0 \) and \( \gamma_1 = \beta_1 \)
• Arrow Extensionality
\[ \forall F,G: A \to B, \text{ if } F \neq G \text{ then } \exists f: 2 \to A \text{ with } Ff \neq Gf. \]

A Grothendieck Universe may finally be defined as a set, U, which satisfies the axioms of ZF set theory without choice, yet as augmented by at least strongly inaccessible large cardinals. The axioms of ZF are:

• Extensionality
\[ \forall x,y. (\forall z. z \in x \iff z \in y) \to x = y \]
• Empty Set
\[ \exists x. \forall y. y \notin x \]
• Pairing
\[ \forall x,y. \exists z. \forall w. w \in z \iff w = x \lor w = y \]
• Union
\[ \forall x. \exists y. \forall z. z \in y \iff \exists w. w \in x \land z \in w \]

172
• Powerset
  \[ \forall x. \exists y. \forall z. z \in y \iff z \subseteq x \]

• Separation (with \( \overrightarrow{x} \) a parameter)
  \[ \forall \overrightarrow{x}, y. \exists z. \forall w. w \in z \iff w \in y \land A(w, \overrightarrow{x}) \]

• Infinity
  \[ \exists x. \emptyset \in x \land \forall y. y \in x \rightarrow y \cup \{y\} \in x \]

• Foundation
  \[ \forall x. (\exists y. y \in x) \rightarrow \exists y. \forall z. z \notin y \]

• Replacement
  \[ \forall x, \overrightarrow{y}. [\forall z. \exists! w. A(z, w, \overrightarrow{y})] \rightarrow \exists u. \forall w. w \in u \iff \exists z. A(z, w, \overrightarrow{y}) \]

• Choice
  \[ \forall x. \emptyset \notin x \rightarrow \exists f \in (x \rightarrow \bigcup x). \forall y. f(y) \in y \]

Large cardinal axioms are defined by elementary embeddings.\(^1\) Elementary embeddings can be defined thus. For models A, B, and conditions \( \phi, j: A \rightarrow B \), \( \phi(a_1, \ldots, a_n) \) in A if and only if \( \phi(j(a_1), \ldots, j(a_n)) \) in B (Kanamori, 2012: 363). A cardinal \( \kappa \) is regular if the cofinality of \( \kappa \) – comprised of the unions of sets with cardinality less than \( \kappa \) – is identical to \( \kappa \) (op. cit.: 360). Uncountable regular limit cardinals are weakly inaccessible (op. cit.). A strongly inaccessible cardinal is regular and has a strong limit, such that if \( \lambda < \kappa \), then \( 2^\lambda < \kappa \) (op. cit.).

By augmenting languages of the theory of CCAF with Grothendieck Universes, U, CCAF proves thereby that:

CCAF \( \vdash \forall n \in \mathbb{N}, \exists \{\aleph_0, \aleph_1, \ldots, \aleph_n\} \), in the category of Sets, U-Set (37-38).

### 11.3 Modal Coalgebraic Automata and Indefinite Extensibility

This section examines, finally, the reasons for which category theory provides a more theoretically adequate setting in which to define indefinite extensibility than do competing approaches such as the Neo-Fregean epistemology

\(^1\)Cf. Koellner and Woodin (2010); Woodin (2010).
of mathematics. According, e.g., to the Neo-Fregean program, concepts of number in arithmetic and analysis are definable via implicit definitions which take the form of abstraction principles. Abstraction principles specify biconditionals in which – on the left-hand side of the formula – an identity is taken to hold between numerical term-forming operators from entities of a type to abstract objects, and – on the right-hand side of the formula – an equivalence relation on such entities is assumed to hold.

In the case of cardinal numbers, the relevant abstraction principle is referred to as Hume’s principle, and states that, for all x and y, the number of the x’s is identical to the number of the y’s if and only if the x’s and the y’s can be put into a one-to-one correspondence, i.e., there is a bijection from the x’s onto the y’s. Abstraction principles for the concepts of other numbers have further been specified. Thus, e.g., Shapiro (2000: 337-340) specifies an abstraction principle for real numbers, which proceeds along the method of Dedekind’s definition of the reals (cf. Wright, 2007: 172). According to the latter method, one proceeds by specifying an abstraction principle which avails of the natural numbers, in order to define pairs of finite cardinals:

$$\forall x,y,z,w \left[ \langle x,y \rangle = \langle z,w \rangle \iff x = z \land y = w \right].$$

A second abstraction principle is defined which takes the differences of the foregoing pairs of cardinals, identifying the differences with integers:

$$\text{Diff}(\langle x,y \rangle) = \text{Diff}(\langle z,w \rangle) \iff x + w = y + z.$$  

One specifies, then, a principle for quotients of the integers, identifying them subsequently with the rational numbers:

$$Q \langle m,n \rangle = Q \langle p,q \rangle \iff n = 0 \land q = 0 \lor n \neq 0 \land q \neq 0 \land m \times q = n \times p.$$  

Finally, one specifies sets of rational numbers, i.e. the Dedekind cuts thereof, and identifies them with the reals:

$$\forall F,G \left[ \text{Cut}(F) = \text{Cut}(G) \iff \forall r (F \leq r \iff G \leq r) \right].$$

The abstractionist program faces several challenges, including whether conditions can be delineated for the abstraction principles, in order for the principles to avoid entraining inconsistency\(^2\); whether unions of abstraction principles can avoid the problem of generating more abstractions than concepts (Fine, 2002); and whether abstraction principles can be specified for mathematical entities in branches of mathematics beyond first and second-order arithmetic (cf. Boolos, 1997; Hale, 2000; Shapiro, op. cit.; and Wright, 2000).

I will argue that the last issue – i.e., being able to countenance definitions for the entities and structures in branches of mathematics beyond first and second-order arithmetic – is a crucial desideratum, the satisfaction of which

\(^2\)Cf. Hodes (1984); Hazen (1985); Boolos (1990); Heck (1992); Fine (2002); Weir (2003); Cook and Ebert (2005); Linnebo and Uzquiano (2009); Linnebo (2010); and Walsh (2016).
remains elusive for the Neo-Fregean program while yet being satisfiable and thus adducing in favor of the modal platonist approach that is outlined in what follows.

One issue for the attempt, along abstractionist lines, to provide an implicit definition for the concept of set is that doing so with an unrestricted comprehension principle yields a principle identical to Frege’s (1893/2013) Basic Law V; and thus – in virtue of Russell’s paradox – entrains inconsistency. However, two alternative formulas can be defined, in order to provide a suitable restriction to the inconsistent abstraction principle. The first, conditional principle states that $\forall F,G[[\text{Good}(F) \lor \text{Good}(G)] \rightarrow [{x|Fx} = {Gx} \iff \forall x(Fx \iff Gx)]$. The second principle is an unconditional version of the foregoing, and states that $\forall F,G[{x|Fx} = {Gx} \iff [\text{Good}(F) \lor \text{Good}(G) \rightarrow \forall x(Fx \iff Gx)]$. Following von Neumann’s (1925/1967: 401-402) suggestion that Russell’s paradox can be avoided with a restriction of the set comprehension principle to one which satisfies a constraint on the limitation of its size, Boolos (1997) suggests that the ‘Good’ predicate in the above principles is intensionally isomorphic to the notion of smallness in set size, and refers to the principle as New V. However, New V is insufficient for deriving all of the axioms of ZF set theory, precluding, in particular, both the axioms of infinity and the power-set axiom (cf. Wright and Hale, 2005: 193). Further, there are other branches of number theory for which it is unclear whether acceptable abstraction principles can be specified. Wiles’ proof of Fermat’s Last Theorem (i.e., that, save for when one of the variables is 0, the Diophantine equation, $x^n = y^n = z^n$, has no solutions when $n > 2$; cf. Hardy and Wright, 1979: 190) relies, e.g., on both invariants and Grothendieck Universes in cohomological number theory (cf. McLarty, 2009: 4).

The foregoing issues with regard to the definability of abstracta in number theory, algebraic geometry (McLarty, op. cit.: 6-8), set theory, et al., can be circumvented in the category-theoretic setting; and in particular by colagebras. In the remainder of this section, I endeavor to demonstrate how modal coalgebraic automata are able to countenance two, fundamental mathematical concepts. The first is the target concept in this essay, namely indefinite extensibility. The second concerns the epistemic and modal properties of the concept of logical consequence, in the $\Omega$-logic in axiomatic set theory.

A labeled transition system is a tuple, $\text{LTS}$, comprised of a set of worlds, $M$; a valuation, $V$, from $M$ to its powerset, $\wp(M)$; and a family of accessibility relations, $R$. So $\text{LTS} = \langle M, V, R \rangle$ (cf. Venema, 2012: 7). A Kripke coalgebra combines $V$ and $R$ into a Kripke functor, $\sigma_R$; i.e. the set of binary morphisms
from M to P(M) (op. cit.: 7-8). Thus, for an s ∈ M, σ(s) := [σ_V(s), σ_R(s)] (op. cit.). σ(s) can be interpreted both epistemically and metaphysically. Satisfaction for the system is defined inductively as follows: For a formula φ defined at a state, s, in M,

\[ [\phi]^M = V(s) \]
\[ [\neg \phi]^M = S - V(s) \]
\[ [\bot]^M = \emptyset \]
\[ [T]^M = M \]
\[ [\phi \lor \psi]^M = [\phi]^M \cup [\psi]^M \]
\[ [\phi \land \psi]^M = [\phi]^M \cap [\psi]^M \]
\[ [\diamond s \phi]^M = \langle R_s \rangle [\phi]^M \]
\[ [\Box s \phi]^M = [R_s][\phi]^M, \]
with
\[ \langle R_s \rangle (\phi) := \{ s' \in S \mid R_s[s'] \cap \phi \neq \emptyset \} \]
\[ [R_s][\phi) := \{ s' \in S \mid R_s[s'] \subseteq \phi \} \]

Modal coalgebraic automata can be thus characterized. Let a category C be comprised of a class Ob(C) of objects and a family of arrows for each pair of objects C(A,B) (Venema, 2007: 421). A functor from a category C to a category D, E: C → D, is an operation mapping objects and arrows of C to objects and arrows of D (422). An endofunctor on C is a functor, E: C → C (op. cit.).

A E-coalgebra is a pair \( A = (A, \mu) \), with A an object of C referred to as the carrier of A, and \( \mu: A \to E(A) \) is an arrow in C, referred to as the transition map of A (390).

Kripke Coalgebras can be availed of in order to model Grothendieck Universes. In CCAF, the elementary embeddings which are jointly necessary and sufficient for positing the existence of large cardinals can further be identified with the transition functions of modal coalgebraic automata. A coalgebraic model of deterministic automata can finally be thus defined (391). An automaton is a tuple, \( A = (A, a_I, C, \delta, F) \), such that A is the state space of the automaton A; \( a_I \in A \) is the automaton’s initial state; C is the coding for the automaton’s alphabet, mapping numerals to properties of the natural numbers; \( \delta: A \times C \to A \) is a transition function, and \( F \subseteq A \) is the collection of admissible states, where F maps A to \{1,0\}, such that F: A → 1 if a∈F and A → 0 if a∉F (op. cit.).

Modal automata are defined over a modal one-step language (Fontaine and Venema, 2018: 3.1-3.2; Venema, 2020: 7.2). With A being a set of

\[ 3 \text{Equivalently, } M,s \vDash \phi \text{ if } s \in V(\phi) \]
propositional variables the set, \( \text{Latt}(X) \), of lattice terms over \( X \) has the following grammar:

\[
\pi ::= \bot \mid \top \mid x \mid \pi \land \pi \mid \pi \lor \pi,
\]

with \( x \in X \) and \( \pi \in \text{Latt}(A) \) (op. cit.).

The set, \( 1\text{ML}(A) \), of modal one-step formulas over \( A \) has the following grammar:

\[
\alpha \in A ::= \bot \mid \top \mid \diamond \pi \mid \square \pi \mid \alpha \land \alpha \mid \alpha \lor \alpha \text{ (op. cit.)}.
\]

A modal P-automaton \( \mathcal{A} \) is a triple, \((A, \Theta, a_I)\), with \( A \) a non-empty finite set of states, \( a_I \in A \) an initial state, and the transition map

\[
\Theta: A \times \wp(P) \to 1\text{ML}(A)
\]

maps states to modal one-step formulas, with \( \wp(P) \) the powerset of the set of proposition letters, \( P \) (op. cit.: 7.3).

Finally, Kripke coalgebras are the dual representations of Boolean-valued models of the \( \Omega \)-logic of set theory (cf. Venema, 2007). Modal coalgebraic automata are able, then, to countenance the constitutive conditions of indefinite extensibility. Modal coalgebraic automata are capable, e.g., of defining both the generative property of indefinite extensibility, as well as the notion of definiteness which figures therein. Further, the category-theoretic definition of indefinite extensibility is arguably preferable to those advanced in the set-theoretic setting, because modal coalgebraic automata can account for both the modal profile and the epistemic tractability of \( \Omega \)-logical consequence.

The \textit{generative} property of indefinite extensibility is captured by the foregoing Kripke functor, \( \sigma_R \), and which we have identified with elementary embeddings, \( j: A \to B, \phi(a_1, \ldots, a_n) \in A \text{ if and only if } \phi(j(a_1), \ldots, j(a_n)) \in B \).

The notion of \textit{definiteness} is captured by the role of modal coalgebraic automata in characterizing the modal profile of \( \Omega \)-logical validity. \( \Omega \)-logical validity can be defined as follows:

For \( T, \{ \phi \} \subseteq \text{Sent} \),

\( T \models_\Omega \phi \), if for all ordinals \( a \) and countable Boolean algebras \( B \), if \( V^B_a \models T \), then \( V^B_a \models \phi \) (Bagaria et al., 2006). The \( \Omega \)-Conjecture states that \( V \models_\Omega \phi \text{ iff } V^B \models_\Omega \phi \) (Woodin, ms). Thus, \( \Omega \)-logical validity is invariant in all set-forcing extensions of ground models in the set-theoretic multiverse.

The invariance property of \( \Omega \)-logical consequence can then be characterized by modal coalgebraic automata. Mathematical truths are thus said to
be definite in virtue of holding of necessity, as recorded by the functors of
the modal colagebraic automata which are dually isomorphic to the Boolean-
valued algebraic models for the $\Omega$-logic of set theory.\footnote{See chapter 9 for further details.}

Thus, whereas the Neo-Fregean approach to comprehension for the con-
cept of set relies on an unprincipled restriction of the size of the universe
in order to avoid inconsistency, and one according to which the axioms of
ZF still cannot all be recovered, modal coalgebraic automata provide a na-
tural means for defining the minimal conditions necessary for formal grasp
of the concept set. The category-theoretic definition of indefinite extensi-
bility is sufficient for uniquely capturing both the generative property as
well as the notion of definiteness which are constitutive of the concept. The
category-theoretic definition of indefinite extensibility avails of a notion of
mathematical modality which captures both the epistemic property of pos-
sible interpretations of quantifiers, as well as the metaphysical property of
set-theoretic ontological expansion.

\section{11.4 Concluding Remarks}

In this essay, I outlined a number of approaches to defining the notion of
indefinite extensibility, each of which restricts the scope of their character-
ization to set-theoretic languages. I endeavored, then, to define indefinite
extensibility in the setting of category-theoretic languages, and examined
the benefits accruing to the approach, by contrast to the extensional and
modal approaches pursued in ZF.

The extensional definition of indefinite extensibility in ZF was shown to
be insufficient for characterizing the generative property in virtue of which
number-theoretic concepts are indefinitely extensible. The generative prop-
erty of indefinite extensibility in the category-theoretic setting was argued,
by contrast, to be identifiable with the Kripke functors of modal coalgebraic
automata, where the automata model Grothendieck Universes, and Kripke
functors are further identifiable with the elementary embeddings by which
large cardinal axioms can be specified. The modal definitions of indefinite
extensibility in ZF were argued to be independently problematic, in virtue
of endeavoring simultaneously to account for the epistemic properties of in-
definite extensibility – e.g., possible reinterpretations of quantifier domains

178
and mathematical vocabulary – as well as the metaphysical properties of indefinite extensibility – i.e., the ontological expansion of the target domains, without providing an account of how this might be achieved. The Kripke functors definable in Grothendieck universes-as-modal coalgebraic automata can secure these two dimensions, by having both epistemic and metaphysical interpretations. The functors are interpreted both epistemically and metaphysically, such that the functors receive their mappings relative to two parameters, the first ranging over epistemically possible worlds and the second ranging over metaphysically possible worlds. The functors thus receive their values in an epistemic two-dimensional semantics.

Finally, against the Neo-Fregean approach to defining concepts of number, and the limits thereof in the attempt to define concepts of mathematical objects in other branches of mathematics beyond arithmetic, I demonstrated how – by characterizing the modal profile of Ω-logical validity and thus the generic invariance and absoluteness of mathematical truths concerning large cardinals throughout the set-theoretic multiverse – modal coalgebraic automata are capable of capturing the notion of definiteness within the concept of indefinite extensibility.
Chapter 12

A Modal Logic for Gödelian Intuition

'The incompleteness results do not rule out the possibility that there is a theorem-proving computer which is in fact equivalent to mathematical intuition' – Gödel, quoted in Wang (1986: 186).

12.1 Introduction

In his remarks on the epistemology of mathematics, Gödel avails of a notion of non-sensory intuition – alternatively, 'consciousness', or 'phenomenology' (cf. Gödel, 1961: 383) – as a fundamental, epistemic conduit into mathematical truths.\(^1\) According to Gödel, the defining properties of mathematical intuition include (i) that it either is, or is analogous to, a type of perception (1951: 323; 1953,V: 359; 1964: 268); (ii) that it enables subjects to alight

\(^1\)Note however that, in the next subsequent sentence, Gödel records scepticism about the foregoing. He remarks: 'But they imply that, in such a – highly unlikely for other reasons – case, either we do not know the exact specification of the computer or we do not know that it works correctly' [Gödel, quoted in Wang (op. cit.)].

\(^2\)Another topic that Gödel suggests as being of epistemological significance is the notion of 'formalism freeness', according to which the concepts of computability, demonstrability (i.e., absolute provability), and ordinal definability can be specified independently of a background formal language (cf. Gödel 1946, and Kennedy 2013 for further discussion). Kennedy notes however that, in his characterizations of demonstrability and definability, Gödel assumes ZFC as his metatheory (op. cit.: 383). Further examination of the foregoing is beyond the scope of the present essay.
upon new axioms which are possibly true (1953,III: 353,fn.43; 1953,V: 361; 1961: 383, 385; 1964: 268); (iii) that it is associated with modal properties, such as provability and necessity (1933: 301; 1964: 261); and (iv) that the non-sensory intuition of abstractions such as concepts entrains greater conceptual 'clarification' (1953,III: 353,fn.43; 1961: 383). Such intuitions are purported to be both of abstracta and formulas, as well as to the effect that the formulas are true.3

In this essay, I aim to outline the logical foundations for rational intuition, by examining the nature of property (iii). The primary objection to Gödel’s approach to mathematical knowledge is that the very idea of rational intuition is insufficiently constrained.4 Subsequent research has thus endeavored to expand upon the notion, and to elaborate on intuition’s roles. Chudnoff (2013) suggests, e.g., that intuitions are non-sensory experiences which represent non-sensory entities, and that the justificatory role of intuition is that it enables subjects to be aware of the truth-makers for propositions (p. 3; ch. 7). He argues, further, that intuitions both provide evidence for beliefs as well as serve to guide actions (145). Bengson (2015: 718-723) suggests that rational intuition can be identified with the 'presentational', i.e., phenomenal, properties of representational mental states – namely, cognitions – where the phenomenal properties at issue are similarly non-sensory; are not the product of a subject’s mental acts, and so are 'non-voluntary'; are qualitatively gradational; and they both 'dispose or incline assent to their contents' and further 'rationalize' assent thereof.5

3The distinction between 'intuition-of' and 'intuition-that' is explicitly delineated in Parsons (1980: 145-146), and will be further discussed in Section 2.

4See, e.g., Hale and Wright (2002). Wright (2004) provides a vivid articulation of the issue: 'A major — but not the only — problem is that, venerable as the tradition of postulating intuitive knowledge of first principles may be, no-one working within it has succeeded at producing even a moderately plausible account of how the claimed faculty of rational intuition is supposed to work — how exactly it might be constituted so as to be reliably responsive to basic logical validity as, under normal circumstances, vision, say, is reliably responsive to the configuration of middle-sized objects in the nearby environment of a normal human perceiver' (op. cit.: 158).

5A similar proposal concerning the justificatory import of cognitive phenomenology – i.e., the properties of consciousness unique to non-sensory mental states such as belief – can be found in Smithies (2013a,b). Smithies prescinds, however, from generalizing his approach to the epistemology of mathematics.

6Compare Kriegel (2015: 68), who stipulates that 'making a judgment that p involves a feeling of involuntariness' and 'making a judgment always involves the feeling of mobilizing a concept'.
Rather than target objections to the foregoing essays, the present discussion aims to rebut the primary objection to mathematical intuition alluded to above, by providing a logic for its defining properties. The significance of the proposal is thus that it will make the notion of intuition formally tractable, and might thus serve to redress the contention that the notion is mysterious and ad hoc.

In his (1933) and (1964), Gödel suggests that intuition has a constitutively modal profile. Constructive intuitionistic logic is shown to be translatable into the modal logic, S4, with the rule of necessitation, while the modal operator is interpreted as concerning provability.\footnote{For further discussion both of provability logic and of intuitionistic systems of modal logic, see Löb (1955); Smiley (1963); Kripke (1965); and Boolos (1993). L"ob's provability formula was formulated in response to Henkin's (1952) problem concerning whether a sentence which ascribes the property of being provable to itself is provable. (Cf. Halbach and Visser, 2014, for further discussion.) For an anticipation of the provability formula, see Wittgenstein (1933-1937/2005: 378), where Wittgenstein writes: 'If we prove that a problem can be solved, the concept 'solution' must somehow occur in the proof. (There must be something in the mechanism of the proof that corresponds to this concept.) But the concept mustn't be represented by an external description; it must really be demonstrated. / The proof of the provability of a proposition is the proof of the proposition itself' (op. cit.). Wittgenstein distinguishes the foregoing type of proof with 'proofs of relevance' which are akin to the mathematical, rather than empirical, propositions, discussed in Wittgenstein (2001: IV, 4-13, 30-31).} Mathematical intuition of set-theoretic axioms is, further, purported both to entrain 'intrinsic' justification, and to illuminate the 'intrinsic necessity' thereof.\footnote{Gödel (1964) does not define intrinsic justifications, although he does contrast the latter with the notion of extrinsic justifications, for which he provides a few defining remarks. Extrinsic justifications are associated, for example, with both the evidential probability of propositions, and the 'fruitful' consequences of a mathematical theory subsequent to adopting new axioms. See Gödel (op. cit.: 269).} Following Gödel's line of thought, I aim, in this essay, to provide a modal logic for the notion of 'intuition-that'.\footnote{Cf. Parsons (1979-1980; 1983: p. 25, chs.10-11; 2008: 176).}

If rational intuition is identical with cognitive phenomenal properties of representational states such as beliefs and judgments, then – via correspondence results between modal logic and first-order logic [cf. van Benthem (1983; 1984/2003)] – a precise translation can be provided between the notion of 'intuition-of', i.e., the cognitive phenomenal properties of thoughts whose contents can concern the axioms of mathematical languages, and the
modal operators regimenting the notion of 'intuition-that'.\footnote{This provides a precise answer to the target inquiry advanced by Parsons (1993: 233).} I argue, then, that intuition-that can further be shown to entrain property (iv), i.e. conceptual elucidation, by way of figuring as an interpretational modality which induces the reinterpretation of domains of quantification (cf. Fine, 2006). Fine (op. cit.) has suggested that the interpretational modality is imperatival, and that the deontic aspects of the modality might best be captured by a dynamic logic (p.c.). Following Fine’s suggestion, I argue that intuition-that can thus be understood to be a species of fixed point dynamic provability logic, which is equivalent to the bisimulation-invariant fragment of monadic second-order logic (cf. Janin and Walukiewicz, 1996; Venema, 2014, ms). Modalized rational intuition is therefore expressively equivalent to – and can crucially serve as a guide to the interpretation of – the entities, such as mathematical concepts, that are formalizable in monadic first- and second-order formal languages.

In Section 2, I elucidate the properties of rational intuition, by examining arguments and evidence adducing in favor of the existence of cognitive phenomenal consciousness. In Section 3, I countenance and motivate a multi-modal logic, which augments the provability logic, GL, with fixed point dynamic logic, i.e. the modal $\mu$-calculus. I argue that the dynamic properties of modalized rational intuition provide a precise means of accounting for the manner by which intuition can yield the reinterpretation of quantifier domains and mathematical vocabulary; and thus explain the role of rational intuition in entraining conceptual elucidation. In Section 4, I examine remaining objections to the viability of rational intuition and provide concluding remarks.

12.2 Rational Intuition as Cognitive Phenomenology

A property of a mental state is phenomenal only if it is the property of being aware of the state. If the mental state at issue is sensory, then sensory phenomenal properties will be properties of being aware of one’s perceptions, where the perceptions at issue will be unique to the subject’s sensory modal-
Let cognitive phenomenal consciousness refer to the properties of being aware of the non-sensory representational mental states toward which subjects can bear attitudes. Such states can be identified with, e.g., formulas, \( \phi \), in a language of thought,\(^{12}\) the syntax for which mirrors that of natural language sentences, and which can fall within the scope of various operators such as fully or partially believing that \((x, x/\phi)\), knowing that \((x/\phi)\), judging that \((x/\phi)\), asserting that \((x/\phi)\), questioning whether \((x/\phi)\) has a particular semantic value, et al.

Pitt (2004: 8) provides the following argument for the existence of cognitive phenomenal properties, which – for the purposes of this essay – I will assume to be sound:

(K1) 'It is possible immediately to identify one’s occurrent conscious thoughts (equivalently (see below): one can know by acquaintance which thought a particular occurrent conscious thought is); but

(K2) It would not be possible immediately to identify one’s conscious thoughts unless each type of conscious thought had a proprietary, distinctive, individuative phenomenology; so

(P) Each type of conscious thought – each state of consciously thinking that \(p\), for all thinkable contents \(p\) – has a proprietary, distinctive, individuative phenomenology'.

In his examination of the conditions on measuring partial beliefs, i.e., subjective probability, Ramsey (1926) records scepticism about whether subjects are aware in a non-sensory way of all of their (partial) beliefs.\(^{13}\) For the pur-

\(^{11}\)For issues concerning the taxonomy of the sense modalities, see Macpherson (2011). Bottom-up/exogenous, spatial-based, property-based, and diffuse/focal attention arguably comprise a necessary condition on the instantiation of phenomenal properties. The condition is witnessed by the phenomenon referred to as the attentional blink. The attentional blink holds if and only if shifting attentional allocation from one stimulus to another induces a lack of awareness of the first stimulus to which attention was previously distributed. See the essays in Mole et al. (eds.) (2010), for further discussion.


\(^{13}\)See Ramsey (op. cit.: 169): 'Suppose that the degree of a belief is something perceptible by its owner; for instance that beliefs differ in the intensity of a feeling by which they are accompanied, which might be called a belief-feeling or feeling of conviction, and that by the degree of belief we mean the intensity of this feeling. This view would be very inconvenient, for it is not easy to ascribe numbers to the intensities of feelings; but apart from this it seems to me observably false, for the beliefs which we hold most strongly are often accompanied by practically no feeling at all; no one feels strongly about things he takes for granted'. Contrast Koopman (1940: 271), who argues that intuition can serve as a guide to veridical judgments concerning the comparison of probability measures [cf.
poses of this note, it is sufficient that at least some cognitive states are states of which subjects can be aware – where, again, the properties of awareness at issue are non-sensory and purported to be unique to distinct cognitive attitudes, such that being aware of one’s belief that φ will qualitatively differ from one’s awareness of one’s interrogative state concerning whether φ is true.

The evidence for the claim that at least some cognitive states are associated with a unique set of non-sensory properties of awareness has proceeded via introspective reports. In the latter, subjects verbally report upon the valence of their awareness of their states, where their reports are assumed to be reliable. One phenomenon of awareness of one’s thoughts might, e.g., be that of inner speech. It is an open issue whether inner speech has a sensory basis (cf. Prinz, 2012). If not, however, and assuming that introspective report is a reliable method for discerning whether subjects are aware of their states – irrespective of their being able to ascertain a precise value thereof – then there might be properties of awareness that are unique to one’s thoughts and cognitive propositional attitudes.

---

14 The results of the method of introspection are availed of by Pitt (op. cit.), and discussed in the essays in Bayne and Montague (eds.) (2011). For an excellent survey of the experimental paradigms endeavoring to corroborate that intuition can be a source of evidence, see Nagel (2007; 2013a).

15 See, however, Schwitzgebel (2011) for an examination of a series of case studies evincing that introspective report is unreliable as a method for measuring consciousness.


17 Nagel (2013a) examines an approach to intuitions which construes the latter as a type of cognition, rather than as a phenomenal property of judgments. She distinguishes, e.g., between intuition and reflection, on the basis of experimental results which corroborate that there are distinct types of cognitive processing (op. cit.: 226-228). Intuitive and reflective cognitive processing are argued to interact differently with the phenomenal information comprising subjects’ working memory stores. Nagel notes that – by contrast to intuitive cognition – reflective cognition ‘requires the sequential use of a progression of conscious contents to generate an attitude, as in deliberation’ (231).
12.3 Modalized Rational Intuition and Conceptual Elucidation

In this section, I will outline the logic for Gödelian intuition. The motivation for providing a logic for rational intuition will perhaps be familiar from treatments of the property of knowledge in formal epistemology. The analogy between rational intuition and the property of knowledge is striking: Just as knowledge has been argued to be a mental state (Williamson, 2001; Nagel, 2013b); to be propositional (Stanley and Williamson, 2001); to be factive; and to possess modal properties (Hintikka, 1962; Nozick, 1981; Fagin et al., 1995; Meyer and van der Hoek, 1995), so rational intuition can be argued to be a property of mental states; to be propositional, as recorded by the notion of intuition-that; and to possess modal properties amenable to rigorous treatment in systems of modal logic.

I should like to suggest that the modal logic of Gödelian intuition is the bimodal logic combining GL – which is comprised of axioms K, 4, GL, and the rule of necessitation – with the modal $\mu$-calculus.

Let $M$ be a model over the Kripke frame, $\langle W, R \rangle$; so, $M = \langle W, R, V \rangle$.

$W$ is a non-empty set of possible worlds. $R$ is a binary relation on $W$. $V$ is a function assigning proposition letters, $\phi$, to subsets of $W$.

$\langle M, w \rangle \models \phi$ if and only if $w \in V(\phi)$.

$\langle M, w \rangle \models \neg \phi$ iff it is not the case that $\langle M, w \rangle \models \phi$.

$\langle M, w \rangle \models \phi \land \psi$ iff $\langle M, w \rangle \models \phi$ and $\langle M, w \rangle \models \psi$.

$\langle M, w \rangle \models \phi \lor \psi$ iff $\langle M, w \rangle \models \phi$ or $\langle M, w \rangle \models \psi$.

$\langle M, w \rangle \models \phi \rightarrow \psi$ iff, if $\langle M, w \rangle \models \phi$, then $\langle M, w \rangle \models \psi$.

$\langle M, w \rangle \models \phi \iff \psi$ iff \[\langle M, w \rangle \models \phi \iff \langle M, w \rangle \models \psi\].

$\langle M, w \rangle \models \Box \phi$ iff $\forall w'[R(w, w') \text{ and } \langle M, w' \rangle \models \phi]$.

$\langle M, w \rangle \models \mu x. \phi$ iff $\bigcap \{U \subseteq W | [\phi]_\tau[x \mapsto U] \subseteq U\}$ (Fontaine, 2010: 18)

$\langle M, w \rangle \models \nu x. \phi$ iff $\bigcup \{U \subseteq W | U \subseteq [\phi]_\tau[x \mapsto U]\}$ (op. cit.; Fontaine and Place, 2010),

*where $\tau[x \mapsto U]$ is the assignment $\tau$ s.t. $\tau'(x) = U$ and $\tau'(y) = \tau(y)$, for all variables $y \neq x$* (op. cit.).

K states that $\Box(\phi \rightarrow \psi) \rightarrow (\Box \phi \rightarrow \Box \psi)$; i.e., if one has an intuition that $\phi$ entails $\psi$, then if one has the intuition that $\phi$ then one has the intuition that $\psi$. GL states that $\Box(\Box \phi \rightarrow \phi) \rightarrow \Box \phi$; i.e., if one has the intuition that the intuition that $\phi$ entails that $\phi$ is true, then one has the intuition that
that $\phi$. 4 states that $\Box \phi \rightarrow \Box \Box \phi$; i.e., if one has the intuition that $\phi$, then one intuits that one has the intuition that $\phi$. Necessitation states that $\vdash \phi \rightarrow \vdash \Box \phi$. Because intuition—that is non-factive, we eschew in our modal system of axiom T, which states that $\Box \phi \rightarrow \phi$; i.e., one has the intuition that $\phi$ only if $\phi$ is the case [cf. BonJour (1998: 4.4); Parsons (2008: 141)].

In order to account for the role of rational intuition in entraining conceptual elucidation (cf., Gödel, 1961: 383), I propose to follow Fine (2006) and Uzquiano (2015a) in suggesting that there are interpretational modalities associated with the possibility of reinterpreting both domains of quantification (Fine, op. cit.) and the non-logical vocabulary of mathematical languages, such as the membership relation in ZF set theory (Uzquiano, op. cit.).

Fine (2005b) has taken the interpretational modality to be imperatival, and has suggested that a dynamic logic might be an optimal means of formalizing the imperative to reinterpret quantifier domains. He (op. cit.) suggests, further, that the interpretational modality might be characterized as a program, whose operations can take the form of 'simple' and 'complex' postulates which enjoin subjects to reinterpret the domains. Uzquiano’s (op. cit.) generalization of the interpretational modality, in order to target the reinterpretation of the extensions of terms such as the membership relation, can similarly be treated.

In propositional dynamic logic (PDL), there are an infinite number of diamonds, with the form $\langle \pi \rangle$. $\pi$ denotes a non-deterministic program, which in the present setting will correspond to Fine’s postulates adumbrated in the foregoing. $\langle \pi \rangle \phi$ abbreviates 'some execution of $\pi$ from the present state entrains a state bearing information $\phi$'. The dual operator is $[\pi] \phi$, which abbreviates 'all executions of $\pi$ from the present state entrain a state bearing information $\phi$'. $\pi^*$ is a program that executes a distinct program, $\pi$, a number of times $\geq 0$. This is known as the iteration principle. PDL is similarly closed under finite and infinite unions. This is referred to as

---

18 A variant strategy is pursued by Eagle (2008). Eagle suggests that the relation between rational intuition and conceptual elucidation might be witnessed via associating the fundamental properties of the entities at issue with their Ramsey sentences; i.e., existentially generalized formulas, where the theoretical terms therein are replaced by second-order variables bound by the quantifiers. However, the proposal would have to be expanded upon, if it were to accommodate Gödel’s claim that mathematical intuitions possess a modal profile.

the 'choice' principle: If \( \pi_1 \) and \( \pi_2 \) are programs, then so is \( \pi_1 \cup \pi_2 \). The forth condition is codified by the 'composition' principle: If \( \pi_1 \) and \( \pi_2 \) are programs, then \( \pi_1;\pi_2 \) is a program (intuitively: the composed program first executes \( \pi_1 \) then \( \pi_2 \)). The back condition is codified by Segerberg’s induction axiom (Blackburn et al., op. cit: p. 13): All executions of \( \pi^* \) (at \( t \)) entrain the following conditional state: If it is the case that (i) if \( \phi \), then all the executions of \( \pi \) (at \( t \)) yield \( \phi \); then (ii) if \( \phi \), then all executions of \( \pi^* \) (at \( t \)) yield \( \phi \). Formally, \( [\pi^*](\phi \rightarrow [\pi]\phi) \rightarrow (\phi \rightarrow [\pi^*]\phi) \).

We augment, then, the provability logic for Gödelian intuition with the axiom, \( \Box \phi \rightarrow [\pi]\phi \), in order to yield a bimodal, dynamic provability logic thereof.

Crucially, the iteration principle permits \( \pi^* \) to be interpreted as a fixed point for the equation: \( x \iff \phi \lor \Diamond x \). The smallest solution to the equation will be the least fixed point, \( \mu x.\phi \lor \Diamond x \), while the largest solution to the equation, when \( \pi^* \lor \infty \), will be the greatest fixed point, \( \nu x.\phi \lor \Diamond x \). Janin and Walukiewicz (op. cit.) have proven that the modal \( \mu \)-calculus is equivalent to the bisimulation-invariant fragment of second-order logic.

Fine’s simple dynamic-postulational modality takes, then, the form:

'(i) Introduction. \( !x.C(x) \)', which states the imperative to: 'Introduce an object \( x \) [to the domain] conforming to the condition \( C(x) \).'

Fine’s complex dynamic-postulational modalities are the following:

(ii) 'Composition. Where \( \beta \) and \( \gamma \) are postulates, then so is \( \beta;\gamma \). We may read \( \beta;\gamma \) as: do \( \beta \) and then do \( \gamma \); and \( \beta;\gamma \) is to be executed by first executing \( \beta \) and then executing \( \gamma \).

(iii) Conditional. Where \( \beta \) is a postulate and \( A \) an indicative sentence, then \( A \rightarrow \beta \) is a postulate. We may read \( A \rightarrow \beta \) as: if \( A \) then do \( \beta \). How \( A \rightarrow \beta \) is executed depends upon whether or not \( A \) is true: if \( A \) is true, \( A \rightarrow \beta \) is executed by executing \( \beta \); if \( A \) is false, then \( A \rightarrow \beta \) is executed by doing nothing.

(iv) Universal. Where \( \beta(x) \) is a postulate, then so is \( \forall x.\beta(x) \). We may read \( \forall x.\beta(x) \) as: do \( \beta(x) \) for each \( x \); and \( \forall x.\beta(x) \) is executed by simultaneously executing each of \( \beta(x_1), \beta(x_2), \beta(x_3), \ldots \), where \( x_1, x_2, x_3, \ldots \) are the values of \( x \) (within the current domain). Similarly for the postulate \( \forall F.\beta(F) \), where \( F \) is a second-order variable.
(v) Iterative Postulates. Where \( \beta \) is a postulate, then so is \( \beta^* \). We may read \( \beta^* \) as: iterate \( \beta \); and \( \beta^* \) is executed by executing \( \beta \), then executing \( \beta \) again, and so on ad infinitum' (op. cit.: 91-92).

Whereas Fine avails of the foregoing interpretational modalities in order both to account for the notion of indefinite extensibility and to demonstrate how unrestricted quantification can be innocuous without foundering upon Russell’s paradox (op. cit.: 26-30), the primary interest in adopting modal \( \mu \) provability logic as the logic of rational intuition is its capacity to account for reinterpretations of mathematical vocabulary and quantifier domains; and thus to illuminate how the precise mechanisms codifying modalized rational intuition might be able to entrain advances in conceptual elucidation.

Finally, the computational profile of modalized rational intuition can be outlined as follows. In category theory, a category \( C \) is comprised of a class \( \text{Ob}(C) \) of objects a family of arrows for each pair of objects \( C(A,B) \) (Venema, 2007: 421). An \( E \)-coalgebra is a pair \( A = (A, \mu) \), with \( A \) an object of \( C \) referred to as the carrier of \( A \), and \( \mu: A \to E(A) \) is an arrow in \( C \), referred to as the transition map of \( A \) (390). A coalgebraic model of deterministic automata can be thus defined (391). An automaton is a tuple, \( A = \langle A, a_I, C, \delta, F \rangle \), such that \( A \) is the state space of the automaton \( A \); \( a_I \in A \) is the automaton’s initial state; \( C \) is the coding for the automaton’s alphabet, mapping numerals to properties of the natural numbers; \( \delta: A \times C \to A \) is a transition function, and \( F \subseteq A \) is the collection of admissible states, where \( F \) maps A to \( \{1,0\} \), such that \( F: A \to 1 \) if \( a \in F \) and \( A \to 0 \) if \( a \notin F \) (op. cit.). Modal automata are defined as in chapter 2.

Let
\[
\diamond \phi \equiv \nabla \{ \phi, T \},
\]
\[
\Box \phi \equiv \nabla \emptyset \lor \nabla \phi \quad \text{(op. cit.)}
\]
\[
[\nabla \Phi] = \{ w \in W \mid R[w] \subseteq \bigcup \{ [[\phi]] \mid \phi \in \Phi \} \text{ and } \forall \phi \in \Phi, [[\phi]] \cap R[w] \neq \emptyset \}
\]
(17). Let an \( E \)-coalgebraic modal model, \( A = \langle S, \lambda, R[.] \rangle \), such that \( S, s \models \nabla \Phi \) if and only if, for all (some) successors \( s' \in S \), \( [\Phi, \sigma(s)] \in E([R_s]) \) (Venema, 2007: 407), with \( E([R_s]) \) a relation lifting of the satisfaction relation \( \models \subseteq S \times \Phi \). Let a functor, \( K \), be such that there is a relation \( K! \subseteq K(A) \times K(A') \) (17). Let \( Z \) be a binary relation s.t. \( Z \subseteq A \times A' \) and \( \varphi!Z \subseteq \varphi(A) \times \varphi(A') \), with
\[
\varphi!Z := \{ (X,X') \mid \forall x \in X \exists x' \in X' \text{ with } (x,x') \in Z \land \forall x' \in X' \exists x \in X \text{ with } (x,x') \in Z \}
\]
(op. cit.). Then, we can define the relation lifting, \( K! \), as follows:
\( K! := \{ ([\pi, X], [\pi', X']) \mid \pi = \pi' \text{ and } (X, X') \in \mathcal{P}!Z \} \) (Venema, 2012: 17).

The philosophical significance of the foregoing is that the modal logic of rational intuition can be interpretable in the category of modal coalgebraic automata. The foregoing accounts for the distinctively computational nature of the modal profile of rational intuition.

### 12.4 Concluding Remarks

In this note, I have endeavored to outline the modal logic of Gödel’s conception of intuition, in order both (i) to provide a formally tractable foundation thereof, and thus (ii) to answer the primary objection to the notion as a viable approach to the epistemology of mathematics. I have been less concerned with providing a defense of the general approach from the array of objections that have been proffered in the literature. Rather, I have sought to demonstrate how the mechanisms of rational intuition can be formally codified and thereby placed on a secure basis.

Among, e.g., the most notable remaining objections, Koellner (2009) has argued that the best candidates for satisfying Gödel’s conditions on being intrinsically justified are reflection principles, which state that the height of the hierarchy of sets, \( V \), cannot be constructed ‘from below’, because, for all true formulas in \( V \), the formulas will be true in a proper initial segment of the hierarchy. Koellner’s limitative results are, then, to the effect that reflection principles cannot exceed the use of second-order parameters without entraining inconsistency (op. cit.). Another crucial objection is that the properties of rational intuition, as a species of cognitive phenomenology, lack clear and principled criteria of individuation. Burgess (2014) notes, e.g., that the role of rational intuition in alighting upon mathematical truths might be distinct from the functions belonging to what he terms a ‘heuristic’ type of intuition. The constitutive role of the latter might be to guide a mathematician’s non-algorithmic insight as she pursues an informal proof. A similar objection is advanced in Cappelen (2012: 3.2-3.3), who argues that – by contrast to the properties picked out by theoretical terms such as ‘utility function’ – terms purporting to designate cognitive phenomenal properties both lack paradigmatic criteria of individuation and must thereby be a topic of disagreement, in virtue of the breadth of variation in the roles that the notion has been intended to satisfy.

The foregoing issues notwithstanding, I have endeavored to demonstrate
that – as with the property of knowledge – an approach to the notion of intuition—that which construes the notion as a modal operator, and the provision thereof with a philosophically defensible logic, might be sufficient to counter the objection that the very idea of rational intuition is mysterious and constitutively unconstrained.
Chapter 13

An Epistemicist Solution to Curry’s Paradox

13.1 Introduction

This essay targets a series of potential issues for the discussion of, and reso-
lution to, the alethic paradoxes advanced by Scharp, in his Replacing Truth
(2013). I aim, then, to provide a novel, epistemicist treatment to Curry’s
Paradox. The novel, epistemicist solution that I advance enables the reten-
tion of both classical logic and the traditional rules for the alethic predicate:
thruth-elimination and truth-introduction.

In Section 2, Scharp’s replacement strategy is outlined, and his semantic
model is described in detail.

In Section 3, novel extensions of Scharp’s theory to the preface paradox;
to the property version of Russell’s paradox in the setting of unrestricted
quantification; to probabilistic self-reference; and to the sorites paradox are
examined.

Section 4 examines six crucial issues for the approach and the semantic
model that Scharp proffers. The six issues target the following points of
contention:

(i) The status of revenge paradoxes in Scharp’s theory;
(ii) Whether a positive theory of validity might be forthcoming on Scharp’s
approach, given that Scharp expresses sympathy with treatments on which –
in virtue of Curry’s paradox – validity is not identical with necessary truth-
preservation;
The failure of compositionality in Scharp’s Theory of Ascending and Descending Truth (ADT) and whether the theory is not, then, in tension with natural language semantics. The foregoing might be pernicious, given Scharp’s use of consistency with natural language semantics as a condition for the success of approaches to the paradoxes. A related issue concerns whether it is sufficient to redress the failure of compositionality by availing of hybrid conditions which satisfy both Ascending and Descending Truth;

(iv) Whether ADT can generalize, in order to account for other philosophical issues that concern indeterminacy;

(v) Whether Descending Truth and Ascending Truth can countenance the manner in which truth interacts with objectivity. It is unclear, e.g., how the theorems unique to each of Descending Truth and Ascending Truth – respectively, T-Elimination and T-Introduction – can capture distinctions between the reality of the propositions mapping to 1 in mathematical inquiry, by contrast to propositions – about humor, e.g. – whose mapping to 1 might be satisfied by more deflationary conditions; and

(vi) Whether the replacement strategy in general and ADT in particular can be circumvented, in virtue of approaches to the alethic paradoxes which endeavor to resolve them by targeting constraints on the contents of propositions and the values that they signify.

Section 5 examines the alethic paradoxes in the setting of epistemic logic. I outline a novel, epistemicist solution to Curry’s paradox, and the epistemicist theory avoids the series of issues adducing against the ADT approach.

Section 6 provides concluding remarks.

13.2 Scharp’s Replacement Theory

Scharp avers that two main alethic principles target the use of the predicate as a device of endorsement and as a device of rejection. When the truth predicate is governed by (T-Out), then it can be deployed in the guise of a device of endorsement. When the truth predicate is governed by (T-In), then it can be deployed in the guise of a device of rejection.

Scharp’s theory aims to replace truth with two distinct concepts. His explicit maneuver is to delineate the two, smallest inconsistent subsets of alethic principles; and then to pair one of the subsets with one of the replacement concepts, and the other subset with the second replacement concept.

Thus, one replacement concept will be governed by (T-In) and not by
Scharp refers to one of his two, preferred replacement concepts as Descending Truth (henceforth DT). DT is governed by (T-Out).

Scharp refers to the second of his two, preferred replacement concepts as Ascending Truth (henceforth AT). AT is governed by (T-In).

In his ‘Syntactical Treatments of Modality, with Corollaries on Reflexion Principles and Finite Axiomatizability’ (1963), Montague proved that, for any predicate $H(x)$, the following conditions on the predicate are inconsistent.

Montague’s (1963) Lemma 3:

(i) All instances of $H(\phi) \rightarrow \phi$ are theorems.
(ii) All instances of $H[H(\phi) \rightarrow \phi]$ are theorems.
(iii) All instances of $H(\phi)$, where $\phi$ is a logical axiom, are theorems.
(iv) All instances of $H(\phi \rightarrow \psi) \rightarrow [H(\phi) \rightarrow H(\psi)]$ are theorems.
(v) Q – i.e., Robinson Arithmetic – is a subtheory.

Scharp notes that Montague’s conditions target only Predicate-Elimination, and are thus apt for governing DT.

Scharp argues that (v) is necessary, in order for languages that express the theory to refer to their own sentences. Condition (i) is necessary, because it captures (T-Out). Condition (ii) is necessary, because denying iterations of DT entrains a version of the revenge paradox.

Thus, either Condition (iii) or Condition (iv) must be rejected. Condition (iii) states that all tautologies are Descending True. Condition (iv) is an instance of closure. In virtue of considerations pertaining to validity (see Section 3), Scharp is impelled to reject (iv), s.t. DT cannot satisfy closure (151).

### 13.2.1 Properties of DT and AT

Scharp argues that the alethic principles, DT and AT, ought to include the following.

DT ought to satisfy:

- [¬Exc: $D(\neg\phi) \rightarrow \neg D(\phi)$];
- [$\wedge$-Exc: $D(\phi \wedge \psi) \rightarrow D(\phi) \wedge D(\psi)$]; and
- [$\lor$-Imb: $D(\phi) \vee D(\psi) \rightarrow D(\phi \vee \psi)$].

However, DT is not governed by:

- [$\neg$-Imb: $\neg T(\phi) \rightarrow T(\neg\phi)$]; nor by
\[
\lor\text{-Exc: } T(\phi \lor \psi) \rightarrow T(\phi) \lor T(\psi).
\]
AT ought to satisfy:
\[
[\neg\text{-Imb: } \neg A(\phi) \rightarrow A(\neg\phi);
\land\text{-Exc: } A(\phi \land \psi) \rightarrow A(\phi) \land A(\psi); \text{ and}
\lor\text{-Imb: } A(\phi) \lor A(\psi) \rightarrow A(\phi \lor \psi).
\]
However, AT is not governed by:
\[
[\neg\text{-Exc: } T(\neg\phi) \rightarrow \neg T(\phi)]; \text{ nor by}
[\land\text{-Imb: } T(\phi) \land T(\psi) \rightarrow T(\phi \land \psi).
\]
Scharp argues, further:
- that classical tautologies are Descending True;
- that the axioms governing the syntax of the theory are Descending True;
- that the axioms of PA are Descending True, in order to induce self-reference via Gödel-numbering; and
- that the axioms of the theories for both AT and DT are themselves Descending True (152).

DT takes classical values, and, in Scharp’s theory, there are no restrictions on the language’s expressive resources. This is problematic, because 'a' := ‘¬A(x)’ and ‘d’ := ‘¬D(x)’ can be countenanced in the language, and thereby yield contradictions:

Because A(x) is governed by (T-In), ‘a’ entails that A(a), although a states of itself that ¬A(x). Contradiction.

Because D(x) is governed by (T-Out), ‘d’ entails that replacing ‘d’ for x in ‘d’ is not a descending truth, i.e., ¬D(d)]. So – by condition (ii) – ‘D[D(x)] → x’ entails that it is not a descending truth that replacing ‘d’ for x in ‘d’ is not a descending truth [i.e., ¬D(¬d)].

Thus, Scharp concedes that there must be problematic sentences in the language for his theory, s.t. both the sentences and their negations are Ascending True, and s.t. the sentences and their negations are not Descending True (op. cit.).

Scharp endeavors to block the foregoing, by suggesting that DT can be governed by both unrestricted (T-Out), as well as a restricted version of (T-In). Similarly, AT can be governed by both unrestricted (T-In), as well as a restricted version of (T-Out).

To induce the foregoing, Scharp introduces a ‘Safety’ predicate, S(x). A sentence ϕ is safe if and only if ϕ is either (DT and AT) or not AT.

Thus,
\[
S(\phi) \land \phi \rightarrow D(\phi); \text{ and}
S(\phi) \land A(\phi) \rightarrow \phi.
\]
A sentence $\phi$ is unsafe if and only if $\phi$ is AT and not DT:

$$S(\phi) \iff D(\phi) \lor \neg A(\phi).$$

From which it follows that:

$$\neg S(\phi) \iff \neg D(\phi) \land A(\phi), \text{ s.t.}$$

$$A(\phi) \to \neg D(\phi);$$

$$D(\phi) \to A(\phi);$$

$$\neg A(\phi) \to \neg D(\phi);$$

$$\neg \exists \phi [D(\phi) \land \neg A(\phi)] \ (153).$$

Scharp avers too that AT and DT are duals. Thus,

$$D(\phi) \iff \neg A(\phi); \text{ and}$$

$$A(\phi) \iff \neg D(\phi) \ (152).$$

### 13.2.2 Scharp’s Theory: ADT

Scharp’s Theory is referred to as ADT. The necessary principles comprising ADT are as follows (cf. 154):

- **Descending Truth**

  (D1): $D(\phi) \to \phi$

  (D2): $D(\neg \phi) \to \neg D(\phi)$

  (D3): $D(\phi \land \psi) \to [D(\phi) \land D(\psi)]$

  (D4): $[D(\phi) \lor D(\psi)] \to D(\phi \lor \psi)$

  (D5): If $\phi$ is a classical tautology, then $D(\phi)$

  (D6): If $\phi$ is a theorem of PA, then $D(\phi)$

  (D7): If $\phi$ is an axiom of ADT, then $D(\phi)$.

- **Ascending Truth**

  (A1): $\phi \to A(\phi)$

  (A2): $\neg A(\phi) \to A(\neg \phi)$

  (A3): $[A(\phi) \lor A(\psi)] \to A(\phi \lor \psi)$

  (A4): $A(\phi \land \psi) \to [A(\phi) \land A(\psi)]$

  (A5): If $\phi$ maps to the falsum constant, then $\neg A(\phi)$

  (A6): If $\phi$ negates an axiom of PA, then $\neg A(\phi)$

- **Transformation Rules**
(M1): \( D(\phi) \iff \neg A \neg (\phi) \)
(M2): \( S(\phi) \iff D(\phi) \lor \neg A(\phi) \)
(M3): \( S(\phi) \land \phi \rightarrow D(\phi) \)
(M4): \( A(\phi) \land S(\phi) \rightarrow \phi \)

(E1): If \( s \) and \( t \) are terms; \( s = t \); and replacing \( s \) with \( t \) in a sentence \( p \) yields a sentence \( q \); then \( D(p) \iff D(q) \)
(E2): If \( s \) and \( t \) are terms; \( s = t \); and replacing \( s \) with \( t \) in a sentence \( p \) yields a sentence \( q \); then \( A(p) \iff A(q) \)
(E3): If \( s \) and \( t \) are terms; \( s = t \); and replacing \( s \) with \( t \) in a sentence \( p \) yields a sentence \( q \); then \( S(p) \iff S(q) \).

13.2.3 Semantics for ADT

Scharp advances a combination of relational semantics for a non-normal modal logic, as augmented by a neighborhood semantics. (A modal logic is normal if and only if it includes axiom K and the rule of Necessitation; respectively ‘\( \Box[\phi \rightarrow \psi] \rightarrow [\Box \phi \rightarrow \Box \psi] \)’ and ‘\( \vdash \phi \rightarrow \vdash \Box \phi \)’.) He refers to this as xeno semantics.

A model, \( M \), of ADT is a tuple, \( \langle D, W, R, I \rangle \), where \( D \) is a non-empty domain of entities constant across worlds, \( W \) denotes the space of worlds, \( R \) denotes a relation of accessibility on \( W \), and \( I \) is an interpretation-function mapping subsets of \( D \) to \( W \). The clauses for defining truth in a world in the model are familiar:

\[
\langle M, w \rangle \vdash \phi \iff w \in V(\phi)
\]
\[
\langle M, w \rangle \vdash \neg \phi \iff \text{it is not the case that } \langle M, w \rangle \vdash \phi
\]
\[
\langle M, w \rangle \vdash \phi \land \psi \iff \langle M, w \rangle \vdash \phi \text{ and } \langle M, w \rangle \vdash \psi
\]
\[
\langle M, w \rangle \vdash \phi \lor \psi \iff \langle M, w \rangle \vdash \phi \text{ or } \langle M, w \rangle \vdash \psi
\]
\[
\langle M, w \rangle \vdash \phi \rightarrow \psi \iff, \text{if } \langle M, w \rangle \vdash \phi, \text{ then } \langle M, w \rangle \vdash \psi
\]
\[
\langle M, w \rangle \vdash \phi \iff \psi \text{ iff } \langle M, w \rangle \vdash \phi \text{ iff } \langle M, w \rangle \vdash \psi
\]
\[
\langle M, w \rangle \vdash \Box \phi \iff \forall w' [\text{If } R(w, w'), \text{ then } \langle M, w' \rangle \vdash \phi]
\]
\[
\langle M, w \rangle \vdash \Diamond \phi \iff \exists w' [R(w, w') \text{ and } \langle M, w' \rangle \vdash \phi]
\]

Scharp augments his relational semantics with a neighborhood semantics. \( M = \langle D, W, R, I \rangle \) is thus enriched with a neighborhood function, \( N \), which maps sets of subsets of \( W \) to each world in \( W \).

Necessity takes then the revised clause:
\[
\langle M, w \rangle \vdash \Box \phi \iff \exists X \subseteq N(w) \forall w' [\langle M, w' \rangle \vdash \phi \iff w' \in X]
\]

Possibility takes the revised clause:
\[
\langle M, w \rangle \vdash \Diamond \phi \iff \neg [\exists X \subseteq N(w) \forall w' [\langle M, w' \rangle \vdash \neg \phi \iff w' \in X]]
\]
Let L be a language with Boolean connectives, and the operators □, ◦, and Σ. □ is the Descending Truth operator. ◦ is the Ascending Truth operator. Σ is the Safety operator. A xeno model $M = \langle F, R, N, V \rangle$ where F denotes a xeno frame, R is an accessibility relation on wff in L, N is a function from W to $2^W$, and V is an assignment-function from wff in L to the values $[0,1]$. Truth in a world is defined inductively as above. The operators take the following clauses:

**Descending Truth:**

$$\langle M, w \rangle \models □ \phi \iff \forall w' \in W[R \phi(w, w') \to \exists X \subseteq N(w')(\forall v \in W[\langle M, v \rangle \models \phi \iff v \in X])]$$

**Ascending Truth:**

$P(\phi)$ denotes the neighborhood structure – i.e., the set of subsets of worlds – at which $\phi$ is true.

$$\langle M, w \rangle \models ◦ \phi \iff \neg[\forall w' \in W[R \neg \phi(w, w') \to P(\neg \phi) \in N(w')]]$$

**Safety:**

$$\langle M, w \rangle \models Σ \phi \iff \forall w' \in W[R \phi(w, w') \to P(\phi) \in N(w')] \lor \exists w' \in W[R \neg \phi(w, w') \land P(\neg \phi) \in N(w')].$$

A reflexive and co-reflexive xeno frame is equivalent to a neighborhood frame:

- (Reflexivity) $\forall \phi \forall w \in W[R \phi(w, w)]$
- (Co-reflexivity) $\forall \phi \forall w \in W \forall w' \in W[R \phi(w, w') \to w = w']$

A sentential xeno frame is acceptable iff

- (i) $\forall w \in W \ N(w) \neq \emptyset$
- (ii) $\forall w \in W \forall X \subseteq N(w) \ X \neq \emptyset$
- (iii) $\forall w \in W \forall X \subseteq N(w) \ w \in X$
- (iv) $\forall \phi \in L \forall w \in W[R \phi(w, w)]$

(v) if $\phi$ and $\psi$ are of the same syntactic type, then $R \phi = R \psi$

A constant-domain xeno frame is a tuple, $F = \langle W, N, R_f, D \rangle$. A constant-domain xeno model adds an interpretation-function $I$ to $F$, s.t. $I$ maps pairs of $F$ and worlds $w$ to subsets of $D$, s.t. $M = \langle F, R_M, I \rangle$.

A substitution is a function from a set of variables to elements of $D$. A substitution $v'$ is $x$-variant of $v$, if $v(y) = v'(y)$ for all variables $y$.

Thus,

$$\langle M, w \rangle \models_v F[(a_1), \ldots, F(a_m)],$$

where $a_i$ is either an individual constant or variable $f(a_1), \ldots, f(a_m) \in I(F,w)$, s.t. if $a_i$ is a variable $x_i$, then $f(a_i) = v(x_i)$, and if $a_i$ is an individual constant $c_i$, then $f(a_i) = I(c_i)$.
\[\langle M, w \rangle \models \neg \phi \text{ iff it is not the case that } \langle M, w \rangle \models \phi\]
\[\langle M, w \rangle \models \phi \land \psi \text{ iff } \langle M, w \rangle \models \phi \text{ and } \langle M, w \rangle \models \psi\]
\[\langle M, w \rangle \models \lor \phi \land \psi \text{ iff } \langle M, w \rangle \models \phi \text{ or } \langle M, w \rangle \models \psi\]
\[\langle M, w \rangle \models \phi \rightarrow \psi \text{ iff, if } \langle M, w \rangle \models \phi \text{ then } \langle M, w \rangle \models \psi\]
\[\langle M, w \rangle \models \forall x [\phi(x)] \text{ iff for all } x \text{-variant } v' \langle M, w \rangle \models \phi(x)\]
\[\langle M, w \rangle \models \exists x [\phi(x)] \text{ iff for some } x \text{-variant } v' \langle M, w \rangle \models \phi(x)\]
\[\langle M, w \rangle \models \Box \phi \text{ iff } \forall w' \in W[R(\phi)(w, w')] \rightarrow \exists x \in N(w') \forall v \in W[\langle M, w \rangle \models \phi \iff v \in X]\]
\[\langle M, w \rangle \models \Diamond \phi \text{ iff } \exists w' \in W[R(\neg \phi)(w, w')] \text{ and } P(\neg \phi) \text{ is not in } N(w').\]

13.3 New Extensions of ADT

In his discussion of Priest’s (2006) inclosure schema, Scharp disavows of a unified solution to the gamut of paradoxical phenomena (Scharp, 2013: 288). Despite the foregoing, I believe that there are at least four positive extensions of Scharp’s theory of Ascending and Descending Truth that he does not discuss, and yet that might merit examination.

13.3.1 First Extension: The Preface Paradox

The first extension of the theory of ADT might be to the preface paradox. A set of credence functions is Easwaran-Fitelson-coherent if and only if (i) the credences are governed by the Kolmogorov axioms; and it is not the case both (ii) that one’s credence is dominated by a distinct credence, s.t. the distinct credence is closer to the ideal, vindicated world, while (iii) one’s credence is assigned the same value as the remaining credences, s.t. they are tied for closeness (cf. Easwaran and Fitelson, 2015). Rather than eschew consistency in favor of a weaker epistemic norm such as Easwaran-Fitelson coherence, the ADT theorist might argue that consistency can be preserved, because the preface sentence, ‘All of the beliefs in my belief set are true, and one of them is false’ might be Ascending True rather than Descending True.

\[A \text{ credence function is here assumed to be a real variable, interpreted as a subjective probability density. The real variable is a function to the } [0, 1] \text{ interval, and is further governed by the Kolmogorov axioms: normality, } 'Cr(T) = 1'; \text{ non-negativity, } 'Cr(\phi) \geq 0'; \text{ finite additivity, } '\text{for disjoint } \phi \text{ and } \psi, Cr(\phi \cup \psi) = Cr(\phi) + Cr(\psi)'; \text{ and conditionalization, } 'Cr(\phi|\psi) = Cr(\phi \cap \psi) / Cr(\psi)'\]}
Because the models in Scharp’s replacement theory can preserve consistency in response to the Preface, ADT might, then, provide a compelling alternative to the Easwaran-Fitelson proposal.

13.3.2 Second Extension: Absolute Generality

A second extension of Scharp’s ADT theory might be to a central issue in the philosophy of mathematics; namely unrestricted quantification. A response to the latter might further enable the development of the property versions of AT and DT: i.e., being Ascending-True-of and being Descending-True-of. For example, Fine (2005) and Linnebo (2006) advance a distinction between sets and interpretations, where the latter are properties; and suggest that inconsistency might be avoided via a suitable restriction of the property comprehension scheme.² A proponent of Scharp’s ADT theory might be able: (i) to adopt the distinction between extensional and intensional groups (sets and properties, respectively); yet (ii) circumvent restriction of the property comprehension scheme, if they argue that the Russell property, R, is Ascending True-of yet not Descending True-of. The foregoing maneuver would parallel Scharp’s treatment of the derivation, within ADT, of the Ascending and Descending Liars and their revenge analogues (see Section 2.1 above).

13.3.3 Third Extension: Probabilistic Self-reference

A third extension of ADT might be to a self-referential paradox in the probabilistic setting. Caie (2013) outlines a puzzle, according to which:

(1) ‘*’ := ¬CrT(*) ≥ .5

that is, (*) says of itself that it is not the case that an agent has credence in the truth of (*) greater than or equal to .5. As an instance of the T-scheme, (1) yields: ‘T(*) ↔ ¬CrT(*) ≥ .5’. However, CrT(*) ought to map to the interval between .5 and 1. Then, ‘Cr(φ) + Cr(¬φ) ≠ 1’, violating the normality condition which states that one’s credences ought to sum to 1.

In ADT, the probabilist self-referential paradox might be blocked as follows. Axiom (A2) states that ¬A(φ) → A(¬φ); so if it is not an Ascending Truth that φ, then it is an Ascending Truth that not φ. However, (A2)

²See Field (2004; 2008) for a derivation of the Russell property, R, given the ‘naive comprehension scheme: ∀u₁ . . . uₙ∃y[Property(y) ∧ ∀x(x instantiates y ⇐⇒ Θ(x, u₁ . . . uₙ)])’ (2008: 294). R denotes ‘does not instantiate itself’, i.e. ∀x[x ∈ R ⇐⇒ ¬(x ∈ x)], s.t. R ∈ R ⇐⇒ ¬(R ∈ R) (2004: 78).
does not hold for Descending Truth. Thus, in the instance of the T-scheme which states that ‘T(*) \iff \neg\text{Cr}T(*) \geq .5’, the move from ‘\neg\text{Cr}T(*) \geq .5’ to ‘\text{Cr}(\neg T(*)) \geq .5’ is Ascending True, but not Descending True. So, if the move from ‘\neg\text{Cr}T(*)’ to ‘\text{Cr}(\neg T(*))’ is not Descending True, then the transition from ‘\text{Cr}(\phi) + \neg\text{Cr}(\phi) = 1’ to ‘\text{Cr}(\phi) + \text{Cr}(\neg\phi) = 1’ is not Descending True. Similarly, then, to the status of the Descending Liar in ADT, the derivation of probabilistic incoherence from the probabilist self-referential sentence, (1), is Ascending True, but not Descending True.

13.3.4 Fourth Extension: The Sorites Paradox

A fourth extension of ADT might, finally, be to the sorites paradox. Scharp’s xeno semantics is non-normal, such that the accessibility relation is governed by the axioms T (reflexivity) and 4 (transitivity), although not by axiom K. Suppose that there is a bounded, phenomenal continuum from orange to red, beginning with a color hue, c_i, and such that – by transitivity – if c_i is orange, then c_{i+1} is orange. The terminal color hue, in the continuum, would thereby be orange and not red. The transitivity of xeno semantics explains the generation of the sorites paradox. However, xeno semantics appears to be perfectly designed in order to block the paradox, as well: The neighborhood function in Scharp’s xeno semantics for ADT is such that one can construct a model according to which transitivity does not hold. Let M_k be a neighborhood model, s.t. W_k = \{a,b,c\}; N_k(a) = \{a,b\}; N_k(b) = \{a,b,c\}; N_k(c) = \{b,c\}; V_k(\phi) = \{a,b\}. Thus, \langle M_k,a \rangle \vdash \Box \phi; but not \langle M_k,b \rangle \vdash \Box \phi. So, it is not the case that \langle M_k,a \rangle \vdash \Box \Box \phi; so transitivity does not hold in the model. Scharp’s semantics for his ADT theory would thus appear to have the resources both to generate, and to solve, the sorites paradox.\footnote{Scharp suggests that the truth predicate is contextually invariant, although assessment-sensitive (9.4). A second means by which the proposal could be extended in order to account for vagueness is via its convergence with the interest-relative approach advanced by Fara (2000; 2008).}

In the remainder of the paper, I examine six issues for ADT, and then outline an epistemicist diagnosis of Curry’s paradox.
13.4 Issues for ADT

13.4.1 Issue 1: Revenge Paradoxes

- Descent
  \[ \delta \iff '\neg D(\delta) \lor \neg S(\delta)' \]
  (i) Suppose that \(D(\delta)\).
  Then, \(D[^{\neg D(\delta) \lor \neg S(\delta)}]\).
  So, \([\neg D(\delta) \lor \neg S(\delta)]\), contrary to the supposition.
  (ii) Suppose that \([\neg D(\delta) \lor \neg S(\delta)]\).
  So, \(D[^{\neg D(\delta) \lor \neg S(\delta)}]\). So, \(D(\delta)\), contrary to the supposition.

- Ascent
  \[ \alpha \iff '\neg A(\alpha) \lor \neg S(\alpha)' \]
  (i) Suppose that \(A(\alpha)\).
  Then, \(A[^{\neg A(\alpha) \lor \neg S(\alpha)}]\)
  So, \([\neg A(\alpha) \lor \neg S(\alpha)]\), contrary to the supposition.
  (ii) Suppose that \(A[^{\neg A(\alpha) \lor \neg S(\alpha)}]\)
  Then \(A(\alpha)\), contrary to the supposition.

Similarly to the response to the alethic paradoxes, Scharp avers that \(\alpha\) and \(\delta\) are unsafe, and so they are Ascending True although not Descending True.

However, Scharp concedes that ADT does not invalidate all unsafe sentences, because some theorems of ADT are not Descending True (cf. 154). Crucially, then, Scharp’s response, both to the alethic paradoxes and to the revenge sentences which are generated using only the resources of his own theory, fails to generalize. Because some theorems of ADT are not Descending True, some theorems of ADT are unsafe, and therefore Scharp’s proposed restriction to safe predicates in order to avoid paradox serves only, as it were, to temper the flames on one side of the room, while they flare throughout the remainder.

A second maneuver exploits the fact that some unsafe sentences are derivable in ADT.

A sentence, \(\gamma\), comprising the singleton \(U^+\) is positively unsafe iff it is derivable in ADT.

A sentence, \(\gamma\), comprising the singleton \(U^-\) is negatively unsafe iff it is not derivable in ADT.
\[ \gamma \iff \neg D(\gamma) \text{ and } \gamma \text{ is not } U^+ . \]

To show that \( \gamma \) is unsafe, suppose for reductio that \( D(\gamma) \). Then, \( D[\neg D(\gamma) \text{ and } \gamma \text{ is not in } U^+] \), so \( \neg D(\gamma) \) and \( \gamma \) is not in \( U^+ \). So, \( \neg(\gamma) \), contrary to the supposition. So, by reductio, \( \neg D(\gamma) \). So \( \gamma \) is unsafe.

Suppose for reductio that \( \neg A(\gamma) \). Then, \( \neg A[\neg D(\gamma) \text{ and } \gamma \text{ is not in } U^+] \). So, \( \neg \gamma \), i.e. either \( D(\gamma) \) (by the definition of \( \gamma \)) or \( D(\gamma) \), then \( A(\gamma) \) (from the definition of Safety). Thus, by reductio, \( A(\gamma) \). Thus, \( \gamma \) is unsafe.

Suppose for reductio that \( \gamma \) is \( U^+ \), s.t. it is an unsafe theorem of ADT. Some sentences of ADT are not Descending True, e.g. \( \beta \). So, assume that \( \beta \rightarrow \gamma \). So, (a) \( \neg D(\gamma) \) and (b) \( \gamma \) isn’t \( U^+ \). Thus, by reductio, \( \gamma \) is not \( U^+ \). Thus, \( \gamma \) is \( U^- \), i.e. unsafe and not derivable from ADT. (In order to make this proof work, Scharp needs to assume (c), i.e. that \( \beta \) is itself not in \( U^+ \). No argument is advanced for this. In some cases it could so be, and then the proof would be blocked.)

Scharp endeavors to minimize the crucial lacuna in his proposal to the effect that ADT validates sentences that are not Descending True. He argues:

- that a valid argument cannot take one from a \( D(\phi) \) to a \( \neg A(\phi) \);
- that – while \( D(\phi) \) can still entail \( \neg S(\phi) \) (by the construction of the paradoxes in ADT) – \( \neg S(\phi) \) entails \( A(\phi) \);
- that the Descending Liar is unsafe (caveat: the Descending Liar is provable in ADT);
- that the conjunction of the Ascending Liar and its negation is not Ascending True (caveat: the Ascending Liar is unsafe, and unsafe sentences are derivable from ADT); and
- that the axioms of ADT are at least Descending True.

### 13.4.2 Issue 2: Validity

Scharp mentions Field’s (2008) argument against identifying validity with necessary truth-preservation, although does not reconstruct the argument.

In order to argue against identifying validity with necessary truth-preservation, Field draws, inter alia, on Curry’s Paradox.

The argument from Curry’s Paradox is such that – by (T-In) and (T-Out) – one can derive the following. If \( \phi \) is a false sentence then,

1. \( \phi \iff [T(\phi) \rightarrow \bot] \)
2. \( T(\phi) \iff [T(\phi) \rightarrow \bot] \)
3. \( T(\phi) \rightarrow [T(\phi) \rightarrow \bot] \)
4. \([T(\phi) \land T(\phi)] \rightarrow \bot\) (by importation)
5. \(T(\phi) \rightarrow \bot\)
6. \([T(\phi) \rightarrow \bot] \rightarrow T(\phi)\)
7. \(T(\phi)\)
8. \(\bot\)

So, necessary truth-preservation entails contradiction.

However, the argument need not be valid, if one preserves (T-In) and (T-Out) yet weakens the logic. One can avail of the strong Kleene valuation scheme, such that \(|\phi|\) is ungrounded, i.e. maps to 1/2. One can then add a Determinacy operator, such that it is not determinately true that \(\phi\) and it is not determinately true that not \(\phi\); so, it is indeterminate whether \(\phi\).

Field argues, in virtue of the foregoing, that validity ought to be a primitive. In more recent work, Field (2015) argues that validity is primitive if and only if it is 'genuine', such that the notion cannot be identical with either its model-theoretic or proof-theoretic analyses. As an elucidation of the genuine concept, he writes that 'to regard an inference or argument as valid is to accept a constraint on belief […] s.t. (in the objective sense of 'shouldn't') we shouldn't fully believe the premises without fully believing the conclusion' (op. cit.). (The primitivist notion is intended to hold, as well, for partial belief.)

Scharp is persuaded by Field's argument, and endorses, in turn, a primitivist notion of validity, as a primitive canon of reasoning without necessary truth-preservation. Scharp takes this to be sufficient for the retention of Condition (iii), in Montague’s Lemma (151). Scharp does not provide any further account of the nature of validity in the book. In later sections of the book, he reiterates his sympathy with Field’s analysis, and also avails of Kreisel’s ‘squeezing’ argument (section 8.8), to the effect that the primitive notion of validity extensionally coincides with a formal notion of validity (i.e., derivation in a first-order axiomatizable quantified logic with identity).

However, one potential issue is that, in a subsequent passage, Scharp writes that: 'an argument whose premises are the members of the set \(G\) and whose conclusion is \(p\) is valid iff for every point of evaluation \(e\) [i.e., index], if all members of \(G\) are assigned tM-value [i.e., an AT- or DT-value of] 1 at \(e\), then \(p\) is assigned tM-value 1 at \(e'\) (240); and this would appear to be a definition of validity as necessary truth-preservation.

The primitivist approach to validity is the primary consideration that Scharp explicitly avails of, when arguing that closure ought to be rejected (Condition iv, in Montague’s Lemma), rather than rejecting logical tau-
ologies as candidates for the axioms of ADT (Condition iii, in Montague’s Lemma) (151). So, further remarks about the nature of validity would have been welcome. An objection to prescinding from more substantial remarks about the nature of validity might also be that Scharp exploits claims with regard to its uses. So, e.g., he writes that ‘a valid argument will never take one from descending truths to something not ascending true’ (177). However, that claim is itself neither a consequence of either Kreisel’s squeezing argument, nor the primitivist approach to validity.

13.4.3 Issue 3: Hybrid Principles and Compositional-ity

- ∧-T-Imb.
  \[ D(\phi) \land D(\psi) \rightarrow A(\phi \land \psi) \]
  v.
  \[ T(\phi) \land T(\psi) \rightarrow T(\phi \land \psi) \]

- ∨-T-Exc.
  \[ D(\phi \lor \psi) \rightarrow A(\phi) \lor A(\psi) \]
  v.
  \[ T(\phi \lor \psi) \rightarrow T(\phi) \lor T(\psi) \]
  (cf. 147, 171)

Feferman’s (1984) theory countenances a primitive truth predicate (Feferman-true, in what follows); a primitive falsity predicate; as well as a Determinacy operator (op. cit.). This is by salient contrast to Scharp’s approach, on which truth is replaced with DT and AT. Scharp argues that Feferman overemphasizes the significance of the compositionality of his Determinacy operator, at the cost of not having either logical truths or the axioms of his own theory satisfy Feferman-truth. By contrast, Scharp believes that he can avail of hybrid principles, such that it is not a requirement of ADT that Descending Truth and Ascending Truth obey compositionality (157).

One objection to this maneuver is that AT and DT are separated, in the hybrid principles, between the antecedent and consequent of the conditional.4 So, it is unclear whether Scharp’s hybrid principles are sufficient

\(^4\)Thanks here to Stephen Read.
to redress the failure of compositionality in ADT; i.e., there being truth-
conditions for sentences whose component semantic values are, alternatively,
DT and AT.
A further objection is that the foregoing might be in tension with Scharp’s
repeated mention of natural-language semantics, in order to argue against
competing proposals. If natural-language semantics were to vindicate prin-
ciples of compositionality, then this would provide a challenge to the empirical
adequacy of Scharp’s ADT theory, and thereby the viability of his replace-
ment concepts for the traditional alethic predicate.

13.4.4 Issue 4: ADT and Indeterminacy

This issue concerns whether ADT might generalize, in order to account for
other philosophical issues that concern indeterminacy. Whether ADT can
be so extended to other issues, such as vagueness and types of indetermi-
nacy, is not a necessary condition on the success of the theory. However, it
might be a theoretical virtue of other accounts – e.g., classical, paracomplete,
intuitionist, and supervaluational approaches – that they do so generalize;
and the extensions of logic and semantics to issues in metaphysics are both
familiar and legion.\(^5\)

E.g., McGee (1991) suggests replacing the truth predicate with (i) a vague
truth predicate, and (ii) super-truth. The replacement predicates are not
intended for deployment in inferences implicated in reasoning, such as con-
ditional proof and arguments by reductio (155). McGee introduces a Defi-
niteness operator, \(\mu\), in order to yield the notion of super-truth relative to
a set of precisifications. There is thus a truth predicate and a super-truth
predicate. Super-truth is governed by (T-In) and (T-Out). Vague truth is
governed by neither.

Thus:
\(\mu(p)\), then \(\mu\)\(’T(p)’\)
(If \(p\) is definitely true, then \(’p\) is true’ is definitely true)
\(\mu(\neg p)\), then \(\mu\)\(’\neg T(p)’\)
(If \(p\) is definitely not true, then \(’p\) is true’ is definitely not true)
\(\mu(p)\) is vague, then \(’T(p)’\) is vague
(vagueness here is secured by availing of the strong Kleene valuation

\(^5\)Cf. Williamson (2017), for an argument for the retention of classical logic despite the
semantic paradoxes, based on the abductive strength of its generalization.
scheme, such that p is ungrounded, i.e. maps neither to true nor false, and rather to .5)

McGee endeavors to avoid Revenge, by arguing that
‘u’ := ‘u is false or u is vague’
collapses to u is vague. So, u is not definitely true, and not definitely vague. Further, u is not derivable within McGee’s supervaluationist theory, nor within a separate, fixed-point theory that he also advances.

Scharp raises several issues for the supervaluational approach. One issue is that vague sentences cannot be precisified via supervaluation – i.e., rendered determinately true – on pain of Revenge (156).

Scharp argues that Descending Truth and Ascending Truth obey (T-Out) and (T-In), respectively, whereas – according to McGee – vague sentences do not. So, McGee’s replacement restricts expressivity, whereas Scharp argues that there are no expressive restrictions on his proposal. Scharp notes, as well, that some of the axioms for McGee’s theory are not definitely true – and are thus vague and not governed by T-Out or T-In – which would appear to be a considerable objection.

However, Field’s (2008) approach – K3 plus a Determinacy operator, with a multi-valued semantics for the conditional – would appear to remain a viable proposal. Extensions of Field’s proposal can be to an explanation of vagueness (Field, op. cit: ch. 5); to the logic of doxastic states (cf. Caie, 2012); and to the model-theory of metaphysics.

With regard, e.g., to the extension of Field’s treatment of the paradoxes to the logic of doxastic states, Caie demonstrates that – rather than rejecting the Liar sentence – it would no longer be the case, by K3 and indeterminacy, that one could believe the Liar, and it would no longer be the case that one ought not to believe the Liar. In addition to this proposal, I provide in section 5 an epistemicist approach to Curry’s paradox which is able to retain both classical logic and the normal truth rules.

With regard to the extension of Field’s treatment of the paradoxes to the logic and model-theory of metaphysics, consider the following. Given Curry’s paradox, the validity of an epistemic norm might depend, for its explanation, on one’s choice of logic. However, one’s choice of logic might depend for its explanation on considerations from metaphysics. Suppose, e.g., that one distinguishes between fundamental and derivative metaphysical states of affairs. The fundamental states of affairs might concern the entities located in 3n-dimensional spacetime, such as whatever is represented by the wavefunction. The derivative states of affairs might concern emergent entities
located in lower, 3-dimensional spacetime. In order to capture the priority of the fundamental to the derivative, fundamental states of affairs could take the classical values, $[0,1]$; by contrast, derivative states of affairs could take the value $0.5$ (in $K3+\text{indeterminacy}$), such that – while fundamental states are always either true or false – it is not determinate that a derivative state of affairs obtains, and it is not determinate that a derivative state of affairs does not obtain.

On the supervaluational treatment of the paradoxes, the approach can more generally be extended in order to account, e.g., for the metaphysical issues surrounding fission cases and indeterminate survival. Approaches which avail of a supervaluational response to fission scenarios, and similar issues at the intersection of nonclassical logic, metaphysical indeterminacy, and decision theory, can be found, e.g., in Williams (2014).

Thus, while it is not a necessary condition on the success of treatments to the alethic paradoxes that their proposals can generalize – in order, e.g., to aid in the resolution of other philosophical issues such as epistemic and metaphysical indeterminacy – there are viable proposals which can be so extended. The competing approaches thus satisfy a theoretical virtue that might ultimately elude ADT.

### 13.4.5 Issue 5: Descending Truth, Ascending Truth, and Objectivity

Scharp claims that considerations of space do not permit him to elaborate on the interaction between Descending Truth, Ascending Truth, and objectivity (Section 8.3). Suppose that – depending on the target domain of inquiry – the truth-conditions of sentences might be sensitive to the reality of the objects and properties that the sentences concern. So, e.g., second-order implicit definitions for the cardinals might be true only if the terms therein refer to abstract entities. By contrast, what is said in sentences about humor might be true, if and only if the sentence satisfies deflationary conditions such as the T-schema.

Another objection to the replacement strategy, and of Scharp’s candidate replacements in particular, is that it is unclear how – in principle – either Descending Truth or Ascending Truth can be deployed in order to capture the foregoing distinctions.
13.4.6 Issue 6: Paradox, Sense, and Signification

One final objection concerns the general methodology of the book. Scharp proceeds by endeavoring to summarize all of the extant approaches to the alethic paradoxes in the literature, and to marshall at least one issue adducing against their favor. However, there are two approaches to the paradoxes that Scharp overlooks. The first approach targets the notion of what is said by an utterance, i.e. the properties of sense and signification that a sentence might express. One such proposal is inspired by Bradwardine (c.1320/2010) and pursued by Read (2009). According to the proposal, if a sentence such as the Liar does not wholly signify that it is true, then one invalidates T-Introduction for the sentence. In a similar vein, Rumfitt (2014) argues that paradoxical sentences are a type of Scheignedanken, i.e. mock thoughts that might have a sense, although take no value; so, T-Introduction is similarly restricted.

Scharp takes it to be a virtue of his account that he can retain the disquotational principles, even though they get subsequently divided among his replacement concepts. He might then reply to the foregoing proposal by suggesting that they similarly induce expressive restrictions in a manner that his approach can circumvent.

However, there are other approaches which avail of what I shall refer to as the sense and signification strategy, and which eschew neither T-Elimination nor T-Introduction. Modulo a semantics for the conditional, K3 and indeterminacy at all orders ensures not only that hyper-determinacy – and therefore an assignment of classical values to the paradoxical sentences – is circumvented; but, furthermore, that revenge sentences cannot be derived either. Against this approach, Scharp reiterates his concern with regard to restrictions on expression. He writes, e.g., that 'Field avoids revenge only by an expressive limitation on his language' (107). However, a virtue of the approach is that, as in xeno semantics for ADT, T-Elimination and T-Introduction are preserved. Against ADT theory, K3+indeterminacy does not arbitrarily select the alethic principles that the semantic theory should satisfy. Crucially, moreover, the approach does not say more than one should like it to, as witnessed, e.g., by the derivability in ADT of both the DT and AT Liars and their revenge analogues. Rather, the language of paracompleteness and indeterminacy demonstrates that – without the loss of the foundational principles governing the alethic predicate – there are propositions which can satisfy the values in an abductively robust semantic theory.
13.5 Epistemicism and Alethic Paradox

Finally, the second approach that Scharp does not consider is one that has recently been developed by the present author. This approach provides an epistemicist solution to the alethic paradoxes, and is able to retain both classical logic and a univocal, non-replacement, alethic predicate which obeys both truth-elimination and truth-introduction. In the epistemic modal system at issue, the box-operator, $\Box_K$, is interpreted as 'the agent knows that' and the box-operator, $\Box_B$, is interpreted as 'the agent believes that'. Belief entails the dual of knowledge: $B\phi \rightarrow \neg K\neg \phi$. The epistemic modal system validates axiom T, which records the factivity of knowledge: $K\phi \rightarrow \phi$.

Curry’s paradox is, again, the following:\(^6\)

For any false sentence, $\phi$,
1. $\phi \iff [T(\phi) \rightarrow \bot]$
2. $T(\phi) \iff [T(\phi) \rightarrow \bot]$
3. $T(\phi) \rightarrow [T(\phi) \rightarrow \bot]$
4. $[T(\phi) \land T(\phi)] \rightarrow \bot$ (by importation)
5. $T(\phi) \rightarrow \bot$ (by contraction)
6. $[T(\phi) \rightarrow \bot] \rightarrow T(\phi)$
7. $T(\phi)$
8. $\bot$

Rather than weaken the logic in a type-free setting (Field, 2008); eschew contraction (Beall and Murzi, 2012); eschew T-introduction in virtue of defective properties of signification (Read, 2009); or retain classical logic for abductive reasons and disband of one of the truth rules (Williamson, 2017), the current approach argues that Curry’s paradox is classically sound; that the normal truth rules can yet be retained; and that the paradox is problematic because it exhibits an instance of epistemic indeterminacy.

The epistemic indeterminacy entrained by Curry’s paradox occurs because step 5 invalidates axiom K on its epistemic interpretation: $K(\phi \rightarrow \psi)$

---

\(^6\)Read (2010) notes that there are at least three forms of Curry’s paradox. One form, attributed to Albert of Saxony, is conjunctive: for any sentence $A$, $'\phi' := A \land F(\phi)$. A second form, attributed both to Albert of Saxony and Bradwardine, is disjunctive: for any sentence $A$, $'\phi' := A \lor F(\phi)$. A third form, owing to Löb (1955) is such that, for any sentence $A$: $'\phi' := T(\phi) \rightarrow A$. Read (op. cit.) argues that T-introduction is the culprit in the first three forms of the paradox. A fourth variation on Curry’s paradox can be found in Beall and Murzi (2012), who replace the truth predicate with a validity predicate. A fifth form is targeted by Field (2008), and is the version at issue in this talk.
→ (Kφ → Kψ)′. This provides a counter-instance to epistemic closure. To see this, let φ denote 'T(φ)' and ψ denote '⊥'. Then:

(*) K(Tφ → ⊥) → [K(Tφ) → K(⊥)].

Our epistemic modal system validates reflexivity, or axiom T: 'Kφ → φ'. Thus, in (*), K(⊥) is false, because – by reflexivity – only truths can be known. However, K(Tφ) can be known, because it is an instance of T-introduction. Thus, the conditional in the consequent of (*) has itself a true antecedent and false consequent, and is thus false. Finally, the conditional in the antecedent of (*), 'K(Tφ → ⊥)' is true, in virtue of the proof of Curry’s paradox. So the instance of K expressed in (*) has a true antecedent and false consequent, providing a counter-instance to K; and so K is not a valid axiom in our system of epistemic logic. Thus, what is problematic about Curry’s paradox is that – despite being classically sound and entailing contradiction via the normal truth rules – its derivation is epistemically indeterminate, by invalidating axiom K.

13.6 Concluding Remarks

In this essay, I have outlined Scharp’s theory of ADT and its semantics. I then proffered four novel extensions of the theory, and detailed five issues that the theory faces. I outlined, finally, a novel epistemicist diagnosis of Curry’s Paradox.

---

7Whether the system validates further axioms is a question which requires separate treatment. Axioms 4 [Kφ → KKφ]; G [BKφ → KBφ]; 4.4 [K(φ ∧ BKψ) → K(φ ∨ ψ)]; GL [K[K(φ) → φ] → K(φ)]; and Grz [KK(φ → Kφ) → φ → φ], have all been proposed as plausible axioms of epistemic logic. If one follows Priest (op. cit.) in arguing that the sorites paradox and the paradoxes of self-reference have the same form, and one believes that transitivity, i.e. axiom 4 on its epistemic interpretation, is the culprit in the sorites paradox, then perhaps eschewing of axiom 4 with regard to the paradoxes of self-reference might be a viable, similarly epistemicist approach. Cf. Williamson (1994; 2002), for the foundations of the latter proposal.
Chapter 14

Epistemic Modality, Intention, and Decision Theory

14.1 Introduction

Formal treatments of imperatival notions have been pursued both logically and semantically. In the logical setting, deontic claims have been interpreted as types of a modal operator, where a condition holding across the points of a space abbreviates the property of obligation, and its dual abbreviates the property of permissibility. In the twentieth century, research in deontic logic has examined the validity of the rule of necessitation ($\vdash \phi \rightarrow \vdash \Box \phi$) (von Wright, 1981); modal axiom 4 ($\Box \phi \rightarrow \Box \Box \phi$) (cf. Barcan, 1966); and modal axiom GL $[\Box (\Box \phi \rightarrow \phi) \rightarrow \Box \phi]$ (cf. Smiley, 1963). The semantic approach has been inspired by the works of Kratzer (1977, 2012), Stalnaker (1978), and Veltman (1996), arguing that there are modal operators on a set of points which are not straightforwardly truth-conditional, instead recording an update on that set which is taken to be pragmatic (cf. Yalcın, 2012). The types of obligation have proliferated, as variations on the 'ought'-operator – e.g., what one ought to do relative to a time and one’s states of information, by contrast to what one ought to do relative to the facts – have been codified by differences in the array of intensional parameters relative to which the operator receives a semantic value (cf. Yalcın, op. cit.; Cariani, 2013; Dowell, 2013; et al).

1 Deontic logic dates from at least as early as the fourteenth century, in the writings of Ockham, Holcot, and Rosetus. See Knuuttila (1981) for further discussion.
This essay aims to provide a theory of the structural content of the types of intention via a similar modal analysis; to explain the role of intention in practical reasoning; and to answer thereby what I will call the unification problem: i.e., the inquiry into how the various types of intention comprise a unified mental state.² The general significance of the present contribution is that it will provide some foundational structure to the topic, where the previous lack thereof has served only to exacerbate its intransigence.³ I will argue that – similarly to the case of deontic judgment – the foregoing types of intention can be countenanced as modal operators. The defining contours of the contents of the states may thus be targeted via their intensional-semantic profile. The types of intention on which I will focus include (i) the notion of 'intention-in-action', as evinced by cases in which agents act intentionally; (ii) the notion of 'intention-with-which', where an agent’s intentions figure as an explanation of their actions; and (iii) the notion of 'intention-for-the-future', as evinced by an agent’s plans to pursue a course of action at a future time.

I will argue that the unification problem has at least two, consistent solutions. The first manner in which the operations of intention are unified is that they are defined on a single space, whose points are states of information or epistemic possibilities. I argue, then, that the significance of examining how the state of intention interacts with practical reasoning is that it provides a second means by which to account for the unity of intention’s types. Although each type of intention has a unique formal clause codifying its structural content, the notions of 'intention-in-action', 'intention-with-which', and 'intention-for-the-future' are nevertheless unified, because each is directed toward the property of expected utility. Thus, acting intentionally, acting because of an intention, and intending to pursue a course of action

²The unification problem is first examined in Anscombe (1963), and has been pursued in contemporary research by, inter alia, Bratman (1984) and Setiya (2014).

³Compare the aims and methods pursued in the research projects of Fine (1981) and Williamson (2014b): 'The relevance of the undertaking [...] consists mainly in the general advantages that accrue from formalizing an intuitive theory. First of all, one thereby obtains a clearer view of its primitive notions and truths. This is no small thing in a subject [...] that is so conspicuously lacking in proper foundations' (Fine, op. cit.: ); 'The aim is to gain insight into a phenomenon by studying how it works under simplified, rigorously described conditions that enable us to apply mathematical or quasi-mathematical reasoning that we cannot apply directly to the phenomenon as it occurs in the wild, with all its intractable complexity. We can then cautiously transfer our insight about the idealized model back to the phenomenon in the wild’ (Williamson, op. cit.).
at a future time, are mental states whose unification consists in that each
type aims toward the satisfaction of the value of an outcome – the value of
which is the product of a partial belief conditional on one’s acts by the utility
thereof. The dissociation between an agent’s intention to pursue an action
and the causal relevance of the action’s outcome adds in favor of the
characterization of expected utility in the setting of evidential, rather than
causal, decision theory. The proposal that the content of intention is ex-
pected utility has, furthermore, the virtue of generalizing, in order to explain
the nature of the intentions of non-human organisms. The contents of non-
human organisms’ intentions can here be understood as the value intended
by their actions, as sensitive to both their prediction that the outcome will
occur and the utility of its occurrence. Finally, because the aim of intention
is expected utility, a precise account can be provided of how intention relates
to the notions of belief and desire, while yet retaining its status as a unique
mental state.

In Sections 2-3, I delineate the intensional-semantic profiles of the types
of intention, and provide a precise account of how the types of intention
are unified in virtue of both their operations in a single epistemic modal
space and their role in practical reasoning, i.e., evidential decision theory. I
endeavor to provide reasons adducing against the proposal that the types of
intention are reducible to the mental states of belief and desire, where the
former state is codified by subjective probability measures and the latter is
codified by a utility function. Section 4 provides concluding remarks.

14.2 The Modes of Intention

The epistemic modal space of an agent can be defined via a frame, comprised
of a set of points, and a relation of accessibility thereon (cf. Kripke, 1963;
Blackburn et al, 2001). The points in the frame are here interpreted as an
agent’s states of information, while the relation of accessibility can receive
various interpretations. A state of information is possible, just if there is at
least one point relative to which it is true, if and only if it is not necessary for
the formula to be false. One of the states of information is necessary, just if it
is true everywhere, i.e. relative to all the other points in the space, if and only
if it is impossible for it to be false. The distinctly epistemic interpretation
of possibility comes in at least two guises, defined as the dual of epistemic
necessity (’◊φ’ iff ’¬□¬φ’): The truth of a formula is epistemically possible,
just if the formula is believed by an agent, or is conceivable to the agent. The epistemic interpretation of necessity can itself come in at least two guises: The truth of a formula is epistemically invariant or necessary, just if the truth of the formula is known by an agent, or if it is inconceivable for the formula to be false (dually, epistemically necessary), and is thus in one sense apriori.

When an agent intends to φ, their intention may fall into three distinct types. One type of intention concerns the intentional pursuit, by the agent, of a course of action. A second type of intention can be witnessed, when the agent cites an intention as an explanation of her pursuit of a course of action. Finally, a third type of intention can be witnessed, when the agent intends to pursue a course of action at a future time.

14.2.1 Intention-in-Action

If the agent acts intentionally, then her intention can be understood as an operation relative to her states of information. The agent acts intentionally, just if there is a world and a unique array of intensional parameters relative to which her intention is realized and receives a positive semantic value. The array of intensional parameters is two-dimensional, because the value of intending to φ relative to one of the parameters will constrain the value of intending to φ relative to the subsequent parameters. Thus, we can say that an agent intends to φ, if and only if she acts intentionally, only if there is both a world and array of intensional parameters, relative to which her intention is realized, i.e. receives a positive value. The intensional parameters include a context comprised of a time and location, and a pair of indices on which spaces of the agent’s acts and of the outcomes of her actions are built. So, the agent’s intention-in-action receives a positive semantic value only if there is at least one world in her epistemic modal space at which – relative to the context of a particular time and location, which constrains the admissibility of the actions as defined at a first index, and which subsequently constrains the outcome thereof as defined at a second index – the intention is realized.

• \[\text{Intenton-in-Action}(\phi)\] = 1 only if \(\exists w' [\phi]_{w',c(=t,l),a,o} = 1\)

14.2.2 Intention-with-which

If the agent refers to an intention, in order to explain her pursuit of a course of action, then her intention can similarly be understood as an operation
relative to her states of information. In this case, the agent intends to $\phi$, just if there is a pair of formulas defined at points in her epistemic modal space, where one of the states is realized because it holds in virtue the other state being realized. Informally, the foregoing explanation can be referred to as the intention-with-which she acts. Thus, we can say that an agent intends to $\phi$, if and only if her intention is an explanation for her action, only if she acts in pursuit of $\psi$ because she intends to $\phi$. In order to capture the notion of one formula holding in virtue, or because, of a distinct formula, we define grounding operators on the agent’s epistemic modal space. Thus, the agent intends to $\phi$ because, there is an intention in virtue of which her action, $\psi$ so as to realize $\phi$, receives a positive value.

- $\llbracket$Intention-with-which($\phi$)$\rrbracket_w = 1$ only if $\exists w'$[$\llbracket\psi\rrbracket_{w'} = 1 \land [G(\phi,\psi)] = 1$],

where $G(x,y)$ is a grounding operator encoding the explanatory connection between $\phi$ and $\psi$. Following Fine (2012b,c), the grounding operator can have the following properties: The grounding operator is weak if and only if it induces reflexive grounding. The operator is strict if and only if it is not weak. The operator is full if and only if the intention to $\phi$ uniquely provides the explanatory ground for the action, $\psi$. The operator is part if and only if the intention to $\phi$ - along with other reasons for action - provide the explanatory ground for the action, $\psi$. Combinations of the foregoing explanatory operators may also obtain: $x < y$ iff $\phi$ is a strict full ground for $\psi$; $x \leq y$ iff $\phi$ is a weak full ground for $\psi$; $x \triangleleft y$ iff $\phi$ is a strict part ground for $\psi$; $x \preceq y$ iff $\phi$ is a weak part ground for $\psi$; $x \triangleleft^* y$ iff $\phi$ is a strict partial ground for $\psi$; $x \triangleleft^{*'} y$ iff $\phi$ is a partial strict ground for $\psi$; $x \triangleleft^{*'} z$ iff $[\phi \triangleleft^* \psi \land \psi \preceq \mu]$ iff $\phi$ is a part strict ground for some further action, $\mu$.

14.2.3 Intention-for-the-Future

Finally, an agent can intend to $\phi$, because she intends to pursue a course of action at a future time. In this case, the intensional-semantic profile which records the parameters relative to which her intention receives a positive semantic value converges with a future-directed modal operator to the effect that the agent will $\phi$. Thus, an agent realizes an intention-for-the-future only if there is a possible world and a future time, relative to which the possibility that $\phi$ is realized can be defined. Thus:
\[ \left[ \text{Intention-for-the-future(\(\phi\))} \right]_w = 1 \text{ only if } \exists w' \forall t \exists t' [t < t' \land \left[ \phi \right]^{w',t'} = 1]. \]

This section has endeavored to accomplish two aims. The first was to provide a precise delineation of the structural content of, and therefore the distinctions between, the types of intention. Intention was shown to be a modal mental state, whose operations have a unique intensional profile, and whose values are defined relative to an agent’s space of states of information. The second aim was to secure one of the means by which the unity of the distinct types of intention can be witnessed. Despite that each of the types of intention has a unique structural content, the contents of those types are each defined in a single, encompassing space; i.e, relative to the agent’s space of epistemic possibilities.

### 14.3 Intention in Decision Theory

In Section 2, I suggested that intention is a unified, modal mental state, the contents of which are defined relative to an agent’s states of information. This section examines the proposal that intentions have a dual profile (cf. Bratman, op. cit.), because intentions figure constitutively in practical reasoning. I argue that, because expected utility theories are the only axiomatized theories of practical reasoning, an account must be provided of the role that intention plays therein. The account will illuminate a precise relationship – which I argue is not identity – between the types of intention and the mental states of belief and desire. The account will further serve to provide a second explanation for the unity of intention’s types, given the uniform role that the types of intention play in decision theory.

A model of decision theory can be understood as a tuple \( \langle A, O, K, V \rangle \), where \( A \) is a set of acts; \( O \) is a set of outcomes; \( K \) encodes a set of counterfactual conditionals, where an act from \( A \) figures in the antecedent of the conditional and \( O \) figures in the conditional’s consequent; and \( V \) is a function assigning a real number to each outcome. The real number is a representation of the value of the outcome. The expected value of the outcome is calculated

---

\(^{4}\text{See Rao and Georgeff (1991), for the suggestion that operators in a multi-modal logic can model the notion of goal-oriented intention. The foregoing intensional semantics is consistent with the logic that they proffer.}\)
as the product of (i) the subjective probability – i.e., the agent’s partial belief or credence – that the outcome will occur, as conditional on her act, and (ii) the value or utility which she assigns to the outcome’s occurrence. The agent can prefer one assignment of values to the outcome’s occurrence over another. (Which preference axioms ought to be adopted is a contentious issue, and will not here be examined. Cf. von Neumann and Morgenstern, 1944; Savage, 1954; Jeffrey, 1983; and Joyce, 1999.) In evidential decision theory, the expected utility of an outcome is calculated as the product of the agent’s credence conditional on her action, by the utility of the outcome. In causal decision theory, the expected utility of an outcome is calculated as the product of the agent’s credence, conditional on both her action and the causal efficacy thereof, by the utility of the outcome. Expected utility can further be augmented by a risk-weighting function: If the agent’s expected utility diminishes with the order of the bets she might pursue – such that expected utility is sensitive to the agent’s propensity to take risks relative to the total ordering of the gambles – then she might have a preference for a sure-gain of .5 units of value, rather than prefer a bet with a 50 percent chance of winning either 0 or 1 units of value (cf. Buchak, 2014).

If intention plays a constitutive role in practical reasoning, and decision theories provide the most tractable models thereof, then what is the role of intention in decision theory? The parameters in the axiomatizations of decision theory encode variables for credences, actions, outcomes, assignments of utility, background states of information pertinent to the causal relevance of actions on outcomes, and the agent’s preferences. Expected utility is derived, as noted, by the interaction between an agent’s credences, actions, and utility assignments. Which, then, of these parameters do an agent’s intentions concern?

There are dissociations between intention and belief and between intention and desire. An agent can have a partial belief that the sun will rise, without intending to pursue any course of action. Conversely, an agent can intend to pursue a course of action, yet appreciate that there are, unfortunately, reasons for her to disbelieve that the act will obtain. An agent can desire that the sun rises, without the intention to entrain the sun’s rising as consequence. Conversely, a vegetarian can intend to consume meat, if it is the only available source of protein and they are in dire need thereof, while yet desire a distinct and orthogonal outcome.

There are dissociations between intention and preference. An agent can prefer the sun’s rising to the prevalence in her life of unprovoked antagonists,
without either acting intentionally, possessing an intention as an explanation for some course of action, or intending to pursue any particular course of action in the future. Conversely, whether an agent’s intention to pursue an action mandates a preference for the value of the outcome of that action will depend on one’s preference axioms. One such axiom might be maximin, according to which the best of the worst outcomes among a set of options should be preferred, while a distinct rule might be maximax, according to which one ought to prefer and pursue the maximally valuable outcome among a set of options. Thus, intending to $\phi$ is not sufficient for determining whether $\phi$ ought to be preferred.

There are, finally, dissociations between intention and action. One might intend to calculate the value of a formula, yet not be able so to act, because their attention might be allocated elsewhere.

Acting intentionally, intending to pursue a course of action in the future, and citing an intention as an explanation for one’s course of action are each, however, in some way related to the value of a course of action. When an agent acts intentionally, she acts in such a way as to obtain an outcome that she values. When an agent pursues a course of action, and refers to her intention so to act as the explanation for that action, the intention explains the value, for the agent, in which the action and its outcomes are supposed to consist. Finally, when an agent intends to pursue a course of action in the future, her intention is similarly guided by the value of the outcome that her action will hopefully entrain. The value of the outcome will not be her bare assessment of the utility of the outcome, because – in the setting of decision theory – utility functions codify desires, such that her intention would thereby be elided with her desire for the outcome.

Because the types of intention are all directed toward the value of an outcome of a course of action – while being irreducible to, because dissociable from, the states of belief and desire – the remaining and most suitable candidate for the role of the mental state of intention in decision theory is the aim of expected utility; i.e., the value of an outcome, as arising by the interaction between the agent’s partial belief or expectation that the outcome will occur as conditional on her act, and the utility that she associates with the outcome’s occurrence. Because of the dissociation between an agent’s intention to pursue a course of action in the future and the action’s occurrence – let alone the dissociation between the intention to act in future, and the causal efficacy of the action were it to obtain – the role of intention in practical reasoning appears to be more saliently witnessed in the setting of evidential
decision theory.

That the types of intention are each directed toward expected utility evinces how an agent’s intentions can be sensitive to her beliefs and desires, without being reducible to them. Crucially, moreover, that the types of intention are each directed toward expected utility provides a second explanation of the way that the types of intention comprise a unified mental state.

Theoretical advantages accruing to the foregoing proposal include that it targets a foundational role for intention in decision theory. The proposal might be foundational, because it targets a basic role for intention in practical reasoning, which is consistent with the possible augmentation of the proposal with other approaches which assume a more cognitively demanding role for intention’s aims. Such approaches include proposals to the effect (i) that the most fundamental type of intention is intention-with-which, such that intention’s role as an explanation can be elided with its causal efficacy (cf. Anscombe, op. cit.; Davidson, 1963); (ii) that the content of intention is the diachronic satisfaction of self-knowledge (cf. Velleman, 1989); and (iii) that the role of intention in practical reasoning ought to be understood as an evaluative constraint, as determined by the virtuous traits of an agent’s character (cf. Setiya, 2007).

14.4 Concluding Remarks

I have argued that the unification problem for the types of intention can be solved in two, consistent ways. The types of intention can be modeled as modal operators, where the unity of the operations consists, in the first instance, in that their values are defined relative to a single, encompassing, epistemic modal space. The second manner by which the unity of intention’s types can be witnessed is via intention’s unique role in practical reasoning. I argued that each of the types of intention – i.e., intention-in-action, intention-as-explanation, and intention-for-the-future – has as its aim the value of an outcome of the agent’s action, as derived by her partial beliefs and assignments of utility, and as codified by the value of expected utility in evidential decision theory. A precise account was thereby provided of the role of epistemic modality in the unification of the types of a unique, modal mental state, whose value figures constitutively in decision-making and practical reason.
Bibliography


van Benthem, J. 2010. *Modal Logic for Open Minds*. CSLI.


Boolos, G. 1984. To Be is to Be the Value of a Variable (or to Be some Values of some Variables). *Journal of Philosophy*, 81.


224


225


Chalmers, D. 2010b. Inferentialism and Analyticity.


Cotnoir, A. ms. Are Ordinary Objects Abstracta?


Dever, J. ms. Quantity without Quantities.


Fara, D.G. 2000. Shifting Sands. Philosophical Topics, 28. Originally published under the name 'Delia Graff'.


Friedman, H. 1975/ms. The Analysis of Mathematical Texts, and Their Calibration in Terms of Intrinsic Strength I.


237
Kamp, H. 1967. The Treatment of ‘Now’ as a 1-Place Sentential Operator. multilith, UCLA.


Leach-Krouse, G. ms. Ω-Consequence Interpretations of Modal Logic.


Pettigrew, R. ms. An Introduction to Toposes.


250


Venema, Y. 2012. Lectures on the Modal $\mu$-Calculus.


Venema, Y. 2020. Lectures on the Modal $\mu$-Calculus.


Watzl, S. Forthcoming. Can Intentionalism Explain how Attention Affects

Waxman, D. ms. Imagining the Infinite.


254


Woodin, W.H. ms. The Ω Conjecture.


