A Two-Dimensionalist Solution to the Access Problem

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In his (1973), Benacerraf inquires into how the semantics, by which he means the theory of truth, for mathematics, might interact with the theory of knowledge for mathematics. He raises the inquiry concerning how knowledge of acausal abstract objects such as those of mathematics (numbers, functions, and sets) is possible, assuming that the best theory of knowledge is that deployed in the empirical sciences and thus presupposes a condition of causal interaction. This is known in the literature in philosophy of mathematics as the access problem. Field (1989) generalizes Benacerraf’s problem by no longer presupposing the condition of causal interaction, and inquiring into what might explain the reliability of mathematical beliefs. Clarke-Doane (2016) has argued that the Benacerraf-Field problem might no longer thought to be pressing in light of mathematical beliefs satisfying conditions of safety and sensitivity. A belief is safe if it could not easily have been different. A belief is sensitive if, had the contents of the belief been false, we would not believe them. Mathematical beliefs are thus sensitive, because mathematical truth are metaphysically necessary, true at all worlds. Clarke-Doane quotes David Lewis, who writes: [I]f it is a necessary truth that so-and-so, then believing that so-and-so is an infallible method of being right. If what I believe is a necessary truth, then there is no possibility of being wrong. That is so whatever the subject matter [...] and no matter how it came to be believed’ (1986: 114-115). Mathematical beliefs are safe, because mathematical truths hold at all nearby worlds, indeed at all of them, and ‘there are reasons to think that we could not have easily had different mathematical beliefs. Our 'core' mathematical beliefs might be thought to be evolutionarily inevitable. Given that our mathematical theories best systematize those beliefs, there is a "bootstrapping" argument for the safety of our belief in those theories’ (24).

In this paper, I argue that the access problem can be solved by availing of the epistemic interpretation of a particular semantics, namely epistemic two-dimensional semantics. The semantics accounts for the truth-conditions of mathematical formulas, while also having an epistemic interpretation of the intensional, and, as I will argue, hyperintensional parameters relative to which those formulas receive their semantic values. Thus epistemic two-dimensional semantics can account for the convergence between the semantics and theory of knowledge for mathematics. Furthermore, however, epistemic two-dimensional semantics countenances two-dimensional intensions, which are functions from
epistemically possible worlds to metaphysically possible worlds to extensions. In epistemic two-dimensional semantics, the value of a formula or term relative to a first parameter ranging over epistemic scenarios determines the value of the formula or term relative to a second parameter ranging over metaphysically possible worlds. The dependence is recorded by 2D-intensions. Chalmers (2006: 102) provides a conditional analysis of 2D-intensions to characterize the dependence: 'Here, in effect, a term’s subjunctive intension depends on which epistemic possibility turns out to be actual. / This can be seen as a mapping from scenarios to subjunctive intensions, or equivalently as a mapping from (scenario, world) pairs to extensions. We can say: the two-dimensional intension of a statement S is true at (V, W) if V verifies the claim that W satisfies S. If [A]₁ and [A]₂ are canonical descriptions of V and W, we say that the two-dimensional intension is true at (V, W) if [A]₁ epistemically necessitates that [A]₂ subjunctively necessitates S. A good heuristic here is to ask ‘If [A]₁ is the case, then if [A]₂ had been the case, would S have been the case?’'. Formally, we can say that the two-dimensional intension is true at (V, W) iff ‘□₁([A]₁ → □₂([A]₂ → S))’ is true, where ‘□₁’ and ‘□₂’ express epistemic and subjunctive necessity respectively'.

Two-dimensional intensions thus provide a conduit from conceivable to metaphysical possibility, and can thus explain the connection between the conceivability of mathematical formulas and their metaphysical possibility. In previous work, I have availed of two-dimensional intensions to account for the interaction between the epistemic and objective or metaphysical profiles of abstraction principles, set-theoretic axioms (including large cardinal axioms), Orey i.e. undecidable propositions, indefinite extensibility, and rational intuition. However, by bridging the epistemic and metaphysical universes, the two-dimensional intensions of epistemic two-dimensional semantics can explain how our epistemic states about mathematical formulas can be a guide to their metaphysical profiles. In this way, epistemic two-dimensional semantics provides a solution to the access problem.

I will define epistemic possibility via the notion of apriority, such that φ is epistemically possible iff φ is primary conceivable, where primary conceivability (⋄) is the dual of apriority (¬■, i.e. not apriori ruled out). So epistemic possibility is the dual of apriority i.e. epistemic necessity, i.e. not apriori ruled out. Chalmers (2002) distinguishes between primary and secondary conceivability. Secondary conceivability is subjunctive, so rejecting the metaphysical necessity of the identity between Hesperus and Phosphorus is not secondary conceivable. Primary conceivability targets epistemically possible (indicative) worlds rather than subjunctive worlds. Chalmers also distinguishes between positive and negative conceivability and prima facie and ideal conceivability. A scenario is positively conceivable when it can be imagined with perceptual detail. A scenario is negatively conceivable when nothing rules it out apriori, as above. A scenario is prima facie conceivable when it is conceivable ‘on first appearances’. E.g. a formula might be prima facie conceivable if it does not lead to contradiction after a finite amount of reasoning. A scenario is ideally conceivable if it is prima facie conceivable with a justification that cannot be
defeated by subsequent reasoning (op. cit.).

Chalmers distinguishes between deep and strict epistemic possibilities. He writes: "We might say that the notion of strict epistemic possibility – ways things might be, for all we know – is undergirded by a notion of deep epistemic possibility – ways things might be, prior to what anyone knows. Unlike strict epistemic possibility, deep epistemic possibility does not depend on a particular state of knowledge, and is not obviously relative to a subject" (62). About deep epistemic necessity, he writes: "For example, a sentence s is deeply epistemically possible when the thought that s expresses cannot be ruled out a priori / This idealized notion of apriority abstracts away from contingent limitations" (66). All references to epistemic possibility in this paper will be to Chalmers' notion of deep epistemic possibility.

Chalmers defines epistemic possibility as (i) not being apriori ruled out (2011: 63, 66), i.e. as the dual of epistemic necessity i.e. apriority (65), and as (ii) being true at an epistemic scenario i.e. epistemically possible world (62, 64). He also accepts a Plenitude principle according to which: 'A thought T is epistemically possible iff there exists a scenario S such that S verifies T' (64). Chalmers advances both epistemic and metaphysical constructions of epistemic scenarios. In the metaphysical construction of epistemic scenarios, epistemic scenarios are centered metaphysically possible worlds (69). In the epistemic construction of epistemic scenarios, they are sentence types comprising an infinitary ideal language, M, with vocabulary restricted to epistemically invariant expressions (75). He defines epistemically invariant expressions thus: "When s is epistemically invariant, then if some possible competent utterance of s is epistemically necessary, all possible competent utterances of s are epistemically necessary" (op. cit.). The sentence types in the infinitary language must also be epistemically complete. A sentence s is epistemically complete if s is epistemically possible and there is no distinct sentence t such that both s \land t and s \land \neg t are epistemically possible (76). The epistemic construction of epistemic scenarios transforms the Plenitude principle into an Epistemic Plenitude principle according to which: 'For all sentence tokens s, if s is epistemically possible, then some epistemically complete sentence of [M] implies s' (op. cit.).

The thesis of ‘weak modal rationalism’ states that conceivability can be a guide to 1-possibility, i.e. conceivability entails 1-possibility or truth at a centered metaphysically possible world (2002). Thus conceivability can be a guide to metaphysical possibility on the metaphysical construction of epistemically possible worlds.

However, in his (2002) and (2010), Chalmers argues that 1-possibility entails 2-possibility, in the case when the primary and secondary intensions for physics and consciousness coincide. Thus, there is no gap between the epistemic and

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1 One might also adopt a conception on which every proposition that is not logically contradictory is deeply epistemically possible, or on which every proposition that is not ruled out a priori is deeply epistemically possible. In this paper, I will mainly work with the latter understanding' (63).

2 We can say that s is deeply epistemically necessary when s is a priori: that is when s expresses actual or potential a priori knowledge’ (65).
metaphysical profiles for expressions involving physics or consciousness, and the conceivability about scenarios concerning them will entail the 1-possibility and the 2-, i.e. metaphysical, possibility of those scenarios.

Finally, in his (2012), Chalmers defines a notion which he refers to as super-rigidity: ‘When an expression is epistemically rigid and also metaphysically rigid (metaphysically rigid \textit{de jure} rather than \textit{de facto}, in the terminology of Kripke 1980), it is super-rigid’ (Chalmers, 2012: 239). He writes: ‘I accept Apriority/Necessity and Super-Rigid Scrutability. (Relatives of these theses play crucial roles in “The Two-Dimensional Argument against Materialism” (241). The Apriority/Necessity Thesis is defined as the ‘thesis that if a sentence S contains only super-rigid expressions, s is a priori iff S is necessary’ (468), and Super-Rigid Scrutability is defined as the ‘thesis that all truths are scrutable from super-rigid truths and indexical truths’ (474). This is thus a third way for conceivability to be a guide to metaphysical possibility. The epistemic necessity i.e. apriority of a sentence involving only super-rigid expressions is such that it converges with the metaphysical necessity of that sentence.

Hyperintensionality in Chalmers’ epistemic two-dimensional intensional semantics can be countenanced via structured intensions, i.e. intensions for each component expression of a sentence, rather than there being an intension for the sentence taken as a whole. However, there are other degrees of hyperintensionality which it would be ideal to capture. One dimension of hyperintensionality can concern sentences being true at parts of worlds rather than at whole worlds themselves. Thus, e.g., ‘snow is white or it is not the case that snow is white’ and ‘grass is green or it is not the case that grass is green’ are necessarily equivalent, but have different contents. In truthmaker semantics, this is owing to the two sentences being made true by different parts of worlds. These parts of worlds which verify and falsify sentences can thus be considered hyperintensional truthmakers and falsemakers (Fine, 2013, 2017.a-c).

Another dimension of hyperintensionality which it would be ideal to capture concerns subject matters. These subject matters are called topics in the literature, and capture the aboutness of atomic formulas. Thus, contents can be defined as pairs of intensions, i.e. functions from worlds to extensions, as well as topics which compose via mereological fusion (Berto, 2018, 2019; Canavotto et al, 2020; and Berto and Hawke, 2021).

In this paper, I will advance a version of epistemic two-dimensional semantics which is a truthmaker semantics and which is topic-sensitive.

According to truthmaker semantics for epistemic logic, a modalized state space model is a tuple \(\langle S, P, \leq, v \rangle\), where S is a non-empty set of states, P is the subspace of possible states where states s and t comprise a fusion when \(s \sqcap t \in P\), \(\leq\) is a partial order, and \(v: \text{Prop} \to (2^S \times 2^S)\) assigns a bilateral proposition \(\langle p^+, p^- \rangle\) to each atom \(p \in \text{Prop}\) with \(p^+\) and \(p^-\) incompatible (Fine 2017a,b; Hawke and Özgün, forthcoming: 10-11). Exact verification (\(\vdash\)) and exact falsification (\(\dashv\)) are recursively defined as follows (Fine, 2017a: 19; Hawke and Özgün, forthcoming: 11):

\[ s \vdash p \text{ if } s \in [p]^+ \]

\(s\) verifies \(p\), if \(s\) is a truthmaker for \(p\) i.e. if \(s\) is in \(p\)’s extension;
\(s \vdash p\) if \(s \in [p]^+\)
(s falsifies \(p\), if \(s\) is a falsifier for \(p\) i.e. if \(s\) is in \(p\)'s anti-extension);
\(s \vdash \neg p\) if \(s \vdash p\)
(s verifies not \(p\), if \(s\) falsifies \(p\));
\(s \vdash \neg p\) if \(s \vdash p\)
(s falsifies not \(p\), if \(s\) verifies \(p\));
\(s \vdash p \land q\) if \(\exists t, u, t \vdash p, u \vdash q,\) and \(s = t \cap u\)
(s verifies \(p\) and \(q\), if \(s\) is the fusion of states, \(t\) and \(u\), \(t\) verifies \(p\), and \(u\) verifies \(q\));
\(s \vdash p \lor q\) if \(s \vdash p\) or \(s \vdash q\)
(s verifies \(p\) or \(q\), if \(s\) verifies \(p\) or \(s\) verifies \(q\));
\(s \vdash \forall x \phi(x)\) if \(\exists s_1, \ldots, s_n,\) with \(s_1 \vdash \phi(a_1), \ldots, s_n \vdash \phi(a_n),\) and \(s = s_1 \cap \ldots \cap s_n\)
(s verifies \(\forall x \phi(x)\) "if it is the fusion of verifiers of its instances \(\phi(a_1), \ldots, \phi(a_n)\)" (Fine, 2017c));
\(s \vdash \exists x \phi(x)\) if \(s \vdash \phi(a)\) for some individual \(a\) in a domain of individuals (op. cit.)
(s falsifies \(\exists x \phi(x)\) "if it falsifies one of its instances" (op. cit.));
\(s \vdash \exists x \phi(x)\) if \(s \vdash \phi(a)\) for some individual \(a\) in a domain of individuals (op. cit.)
(s verifies \(\exists x \phi(x)\) "if it verifies one of its instances \(\phi(a_1), \ldots, \phi(a_n)\)" (op. cit.));
\(s \vdash \exists x \phi(x)\) if \(\exists s_1, \ldots, s_n,\) with \(s_1 \vdash \phi(a_1), \ldots, s_n \vdash \phi(a_n),\) and \(s = s_1 \cap \ldots \cap s_n\) (op. cit.)
(s falsifies \(\exists x \phi(x)\) "if it is the fusion of falsifiers of its instances" (op. cit.));
\(s\) exactly verifies \(p\) if and only if \(s \vdash p\) if \(s \in [p]^+\);
\(s\) inexactly verifies \(p\) if and only if \(s \triangleright p\) if \(\exists s' \leq S, s' \vdash p;\) and
\(s\) loosely verifies \(p\) if and only if, \(\forall t, s.t. s \sqcup t \vdash p,\) where \(\sqcup\) is the relation of compatibility (35-36):
\(s \vdash A \phi\) - i.e. \(A \phi\) - if and only if for all \(t \in P\) there is a \(t' \in P\) such that \(t' \cap t \in P\) and \(t' \vdash \phi,\) where \(A \phi\) denotes the apriority of \(\phi;\) and
\(s \vdash A \phi\) if and only if there is a \(t \in P\) such that for all \(u \in P\) either \(t \cap u \not\in P\) or \(u \vdash \phi.\)

In order to account for two-dimensional indexing, we augment the model, \(M\), with a second state space, \(S^*\), on which we define both a new parthood relation, \(\leq^*\), and partial function, \(V^*\), which serves to map propositions in a domain, \(D\), to pairs of subsets of \(S^*, \{1,0\}\), i.e. the verifier and falsifier of \(p\), such that \([P]^+ = 1\) and \([p]^− = 0\). Thus, \(M = \langle S, S^*, D, \leq, \leq^*, V, V^* \rangle\). The two-dimensional hyperintensional profile of propositions may then be recorded by defining the value of \(p\) relative to two parameters, \(c,i\): \(c\) ranges over subsets of \(S\), and \(i\) ranges over subsets of \(S^*\).
(134) (*): $M, s, s^* \in S, S^* \vdash p$ iff:
(i) $\exists c \in \llbracket p \rrbracket : c, c = 1$ if $s \in \llbracket p \rrbracket$ +
(ii) $\exists c \in \llbracket p \rrbracket : c, c = 1$ if $s^* \in \llbracket p \rrbracket$ +

(Distinct states, $s, s^*$, from distinct state spaces, $S, S^*$, provide a multi-
dimensional verification for a proposition, $p$, if the value of $p$ is provided a
truthmaker by $s$. The value of $p$ as verified by $s$ determines the value of $p$ as
verified by $s^*$).

We say that $p$ is hyper-rigid iff:

(134) (**): $M, s, s^* \in S, S^* \vdash p$ iff:
(i) $\forall c' \in \llbracket p \rrbracket : c, c' = 1$ if $s \in \llbracket p \rrbracket$ +
(ii) $\forall i \in \llbracket p \rrbracket : c, c' = 1$ if $s^* \in \llbracket p \rrbracket$ +

The foregoing provides a two-dimensional hyperintensional semantic frame-
work within which to interpret the values of a proposition:

$s$ is a two-dimensional exact truthmaker of $p$ if and only if (*)
$s$ is a two-dimensional inexact truthmaker of $p$ if and only if $\exists s' \in S, s \rightarrow s'$,
$s' \vdash p$ and such that
$\exists c \in \llbracket p \rrbracket : c, c = 1$ if $s' \in \llbracket p \rrbracket$ +
$\exists i \in \llbracket p \rrbracket : c, c = 1$ if $s^* \in \llbracket p \rrbracket$ +

$s$ is a two-dimensional loose truthmaker of $p$ if and only if, $\exists t, s.t. s \sqcup t, s$
$\sqcup t \vdash p$:
$\exists c \in \llbracket p \rrbracket : c, c = 1$ if $s' \in \llbracket p \rrbracket$ +
$\exists i \in \llbracket p \rrbracket : c, c = 1$ if $s^* \in \llbracket p \rrbracket$ +

Epistemic (primary), subjunctive (secondary), and 2D hyperintensions can
be defined as follows, where hyperintensions are functions from states to exten-
sions, and intensions are functions from worlds to extensions:

- Epistemic Hyperintension:
  $\text{pri}(x) = \lambda s. \llbracket x \rrbracket^{s, s},$ with $s$ a state in an epistemic state space;

- Subjunctive Hyperintension:
  $\text{sec}_v(x) = \lambda i. \llbracket x \rrbracket^{v, i},$ with $i$ a state in metaphysical state space $I$;

- 2D-Hyperintension:
  $\text{2D}(x) = \lambda s \lambda i. \llbracket x \rrbracket^{s, i} = 1.$

Following the presentation of topic models in Berto (op. cit.), atomic topics
comprising a set of topics, $T$, record the hyperintensional intentional content
of atomic formulas, i.e. what the atomic formulas are about at a hyperintensional
level. Topic fusion is a binary operation, such that for all $x, y, z \in T$, the following
properties are satisfied: idempotence $(x \odot x = x)$, commutativity $(x \odot y = y$
$\odot x)$, and associativity $[(x \odot y) \odot z = x \odot (y \odot z)]$ (Berto, 2018: 5). Topic
parthood is a partial order, $\leq$, defined as $\forall x, y \in T (x \leq y) \iff x \oplus y = y$ (op.
cit.: 5-6). Atomic topics are defined as follows: $\text{Atom}(x) \iff \neg \exists y < x$, with
A strict order. Topic parthood is thus a partial ordering such that, for all \( x, y, z \in T \), the following properties are satisfied: reflexivity (\( x \leq x \)), antisymmetry (\( x \leq y \land y \leq x \rightarrow x = y \)), and transitivity (\( x \leq y \land y \leq z \rightarrow x \leq z \)) (6). A topic frame can then be defined as \( \{ W, T, \oplus, t \} \), with \( t \) a function assigning atomic topics to atomic formulas. For formulas, \( \phi, \psi \), atomic formulas, \( p, q, r \) (\( p_1, p_2, \ldots \)), and a set of atomic topics, \( \{ p_1, \ldots, p_n \} \), the topic of \( \phi \), \( t(\phi) = \oplus t(p_1) \oplus \cdots \oplus t(p_n) \) (op. cit.). Topics are hyperintensional, though not as fine-grained as syntax. Thus \( t(\phi) = t(\neg \neg \phi) \), \( t(\phi) = t(\neg \phi) \), \( t(\phi \land \psi) = t(\phi) \oplus t(\psi) = t(\phi \lor \psi) \) (op. cit.).

The diamond and box operators can then be defined relative to topics:

\[
\langle M, w \rangle \models \Diamond^t \phi \iff \langle R_{w,t} \rangle(\phi)
\]

\[
\langle M, w \rangle \models \Box^t \phi \iff [R_{w,t}](\phi) \subseteq \phi
\]

\[
[R_{w,t}](\phi) := \{ w' \in W : w' \in T \land R_{w,t}[w', t'] \cap \phi \neq \emptyset \land t'(\phi) \leq t(\phi) \}
\]

We can then combine topics with truthmakers rather than worlds, thus countenancing a multi-hyperintensional semantics, i.e. topic-sensitive epistemic two-dimensional truthmaker semantics:

- Epistemic Hyperintension:

\[
\text{pri}_u(x) = \lambda s \lambda t. \llbracket x \rrbracket_{s \cap t}^{s \cap t}, \text{ with } s \text{ a truthmaker from an epistemic state space.}
\]

- Subjunctive Hyperintension:

\[
\text{sec}_v \cap t(x) = \lambda w \lambda t. \llbracket x \rrbracket_{w \cap t}^{w \cap t}, \text{ with } w \text{ a truthmaker from a metaphysical state space.}
\]

- 2D-Hyperintension:

\[
2D(x) = \lambda s \lambda w \lambda t. \llbracket x \rrbracket_{s \cap t}^{s \cap t} = 1.
\]

It is easy to see that mathematical sentences - whether arbitrary formulas, axioms, or Orey sentences - can be evaluated two-dimensionally, such that their epistemic profile can be a guide to their metaphysical profile. The two-dimensional hyperintensions of mathematical sentences captures the interaction between the epistemic and objective profiles of the foregoing sentences.

The foregoing proposal also differs from full-blooded platonism in the following respects. According to full-blooded platonism, if a mathematical formula is consistent and thus logically possible, as well as for whatever objects are logically possible, then those formulas are true and those objects exist (Balaguer, 1998). Formulas such as the Continuum Hypothesis (CH) which states that all infinite sets of reals have the cardinality of either the natural numbers or the real numbers, as well as the negation of CH are both logically possible and thus are actually satisfied in different universes (Balaguer, 2001: 97).

Epistemic two-dimensional semantics differs from full-blooded platonism in concerning epistemic possibilities rather than logical possibilities, as well as metaphysical possibility rather than existence. Thus, primary intensions are functions from epistemically possible worlds considered as actual to extensions.
So only epistemically possible worlds considered as actual can be a guide to metaphysical possibility, by contrast to the case of full-blooded platonism according to which any logically possible object or formula actually exists or is true.

A second point of departure from full-blooded platonism is that epistemic two-dimensionalism is consistent with monism about the universe of sets, i.e. there being a cumulative hierarchy of sets comprising a single universe. This contrasts to the set-theoretic pluralism entailed by the unsettled yet logically possible status of both CH and ¬CH as in full-blooded platonism.

Finally, two-dimensional intensions can be availed of as a bridge between what Cantor (1883/1996: §8) refers to as ‘immanent’ mathematical reality and ‘transient’ mathematical reality. Immanent reality concerns what exists relative to the ‘understanding’, whereas transient reality concerns what exists relative to the ‘external world’ (op. cit.). Immanent reality is constrained by conditions of coherence and consistency. Cantor famously argues that mathematics is free to stipulate the existence of any objects or concepts which satisfies those conditions and that they are thus possessed of immanent reality. He leaves it as an open question what conditions need to be satisfied in order for immanent reality to be connected to metaphysics or transient reality, although he appeals to the ‘unity of the all to which we ourselves belong’ (op. cit.) in order to account for their convergence. Two-dimensional intensions are natural candidates for bridging the divide between conceivability and metaphysics, and thus provide a more satisfying explanation of how immanent and transient reality might converge than Cantor’s own.

References


